Tight upper bound on the maximum anti-forcing numbers of graphs *

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Let G be a simple graph with a perfect matching. Deng and Zhang showed that the maximum anti-forcing number of G is no more than the cyclomatic number. In this paper, we get a novel upper bound on the maximum anti-forcing number of G and investigate the extremal graphs. If G has a perfect matching M whose anti-forcing number attains this upper bound, then we say G is an extremal graph and M is a *nice* perfect matching. We obtain an equivalent condition for the nice perfect matchings of G and establish a one-to-one correspondence between the nice perfect matchings and the edge-involutions of G, which are the automorphisms α of order two such that v and $\alpha(v)$ are adjacent for every vertex v. We demonstrate that all extremal graphs can be constructed from K_2 by implementing two expansion operations, and G is extremal if and only if one factor in a Cartesian decomposition of G is extremal. As examples, we have that all perfect matchings of the complete graph K_{2n} and the complete bipartite graph $K_{n,n}$ are nice. Also we show that the hypercube Q_n , the folded hypercube FQ_n ($n \ge 4$) and the enhanced hypercube $Q_{n,k}$ ($0 \le k \le n-4$) have exactly n, n + 1 and n + 1 nice perfect matchings respectively.

Keywords: Maximum anti-forcing number, Perfect matching, Edge-involution, Cartesian product, Hypercube, Folded hypercube

1 Introduction

Let G be a finite and simple graph with vertex set V(G) and edge set E(G). We denote the number of vertices of G by v(G), and the number of edges by e(G). For $S \subseteq E(G)$, G - S denotes the subgraph of G with vertex set V(G) and edge set $E(G) \setminus S$. A *perfect matching* of G is a set M of edges of G such that each vertex is incident with exactly one edge of M. A perfect matching of a graph coincides with a Kekulé structure in organic chemistry.

The *innate degree of freedom* of a Kekulé structure was firstly proposed by Klein and Randić (1987) in the study of resonance structure of a given molecule in chemistry. In general, Harary et al. (1991) called the innate degree of freedom as the forcing number of a perfect matching of a graph. The *forcing number* of a perfect matching M of a graph G is the smallest cardinality of subsets of M not contained in other perfect matchings of G. The *minimum forcing number* and *maximum forcing number* of G are the minimum and maximum values of forcing numbers over all perfect matchings of G, respectively. Computing the minimum forcing number of a bipartite graph with the maximum degree three is an NP-complete problem, see Afshani et al. (2004). As we know, the forcing numbers of perfect matchings have been studied for many specific graphs, see Adams et al. (2004); Che and Cheng (2011); Jiang and Zhang (2011, 2016); Lam and Pachter (2003); Pachter and Kim (1998); Shi and Zhang (2016); Zhang and Deng (2015); Zhang et al. (2010, 2015); Zhao and Zhang (2016).

Vukičević and Trinajstić (2007) defined the anti-forcing number of a graph as the smallest number of edges whose removal results in a subgraph with a unique perfect matching. Recently Lei et al. (2016) introduced the anti-forcing

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number of a single perfect matching M of a graph G as follows. A subset $S \subseteq E(G) \setminus M$ is called an *anti-forcing* set of M if G - S has a unique perfect matching M. The *anti-forcing number* of a perfect matching M is the smallest cardinality of anti-forcing sets of M, denoted by af(G, M). Obviously, the anti-forcing number of G is the minimum value of the anti-forcing numbers over all perfect matchings of G. The *maximum anti-forcing number* of Gis the maximum value of the anti-forcing numbers over all perfect matchings of G, denoted by Af(G). It is an NPcomplete problem to determine the anti-forcing number of a perfect matching of a bipartite graph with the maximum degree four, see Deng and Zhang (2017a). For some progress on this topic, see refs. Vukičević and Trinajstić (2008); Che and Cheng (2011); Deng (2007, 2008); Deng and Zhang (2017a,b,c); Lei et al. (2016); Li (1997); Shi and Zhang (2016); Yang et al. (2015b); Zhang et al. (2011).

For a bipartite graph G, Riddle (2002) proposed the trailing vertex method to get a lower bound on the forcing numbers of perfect matchings of G. Applying this lower bound, the minimum forcing number of Q_n is 2^{n-2} if n is even. However, for odd n, determining the minimum forcing number of Q_n is still an open problem. For the maximum forcing number of Q_n , Alon proved that for sufficiently large n this number is near to the total number of edges in a perfect matching of Q_n (see Riddle (2002)), but its specific value is still unknown. Afterwards, Adams et al. (2004) generalized Alon's result to a k-regular bipartite graph and for a hexagonal system, a polyomino graph or a (4, 6)-fullerene, Xu et al. (2013); Zhang and Zhou (2016); Shi et al. (2017) showed that its maximum forcing number of a perfect matching of G, and showed that the maximum forcing number of G is no more than Af(G). Particularly, for a hexagonal system H, Lei et al. (2016) showed that Af(H) equals the Fries number (see Fries (1927)) of H. Recently, see Shi et al. (2017), we also showed that for a (4, 6)-fullerene graph G, Af(G) equals the Fries number of G.

The cyclomatic number of a connected graph G is defined as r(G) = e(G) - v(G) + 1. Deng and Zhang (2017c) recently obtained that the maximum anti-forcing number of a graph is no more than the cyclomatic number.

Theorem 1.1 (Deng and Zhang (2017c)). For a connected graph G with a perfect matching, $Af(G) \leq r(G)$.

Deng and Zhang (2017c) further showed that the connected graphs with the maximum anti-forcing number attaining the cyclomatic number are a class of plane bipartite graphs. In this paper, we obtain a novel upper bound on the maximum anti-forcing numbers of a graph G as follows.

Theorem 1.2. Let G be any simple graph with a perfect matching. Then for any perfect matching M of G,

$$af(G,M) \le Af(G) \le \frac{2e(G) - v(G)}{4}.$$
(1)

In fact, this upper bound is also tight. By a simple comparison we immediately get that the upper bound is better than the previous upper bound r(G) when 3v(G) < 2e(G) + 4. In next sections we shall see that many non-planar graphs can attain this upper bound, such as complete graphs K_{2n} , complete bipartite graphs $K_{n,n}$, hypercubes Q_n , etc.

We say that a graph G is *extremal* if the maximum anti-forcing number Af(G) attains the upper bound in Theorem 1.2, that is, G has a perfect matching M such that both equalities in (1) hold. Such M is said to be a *nice* perfect matching of G. In Section 2, we give a proof to Theorem 1.2, obtain an equivalent condition for the nice perfect matchings of G, and establish a one-to-one correspondence between the nice perfect matchings of G and the edge-involutions of G. In Section 3, we provide a construction of all extremal graphs, which can be obtained from K_2 by implementing two expansion operations, and show that such a graph is an elementary graph (each edge belongs to some perfect matchings). In Section 4, we investigate Cartesian decompositions of an extremal graph. Let $\Phi^*(G)$ denote the number of nice perfect matchings of a graph G. For a Cartesian decomposition $G = G_1 \Box \cdots \Box G_k$, we obtain $\Phi^*(G) = \sum_{i=1}^k \Phi^*(G_i)$. This implies that a graph G is extremal if and only if in a Cartesian decomposition of G one factor is an extremal graph. As applications we show that three cube-like graphs, the hypercubes Q_n , the

folded hypercubes FQ_n and the enhanced hypercubes $Q_{n,k}$ are extremal. In particular, in the final section we prove that Q_n has exactly n nice perfect matchings and $Af(Q_n) = (n-1)2^{n-2}$, FQ_n $(n \ge 4)$ has exactly n+1 nice perfect matchings and $Af(FQ_n) = n2^{n-2}$, and for $0 \le k \le n-4$, $Q_{n,k}$ has n+1 nice perfect matchings and $Af(Q_{n,k}) = n2^{n-2}$. We also show that FQ_n is a prime graph under the Cartesian decomposition.

2 Upper bound and nice perfect matchings

2.1 The proof of Theorem 1.2

Let G be a graph with a perfect matching M. A cycle of G is called an M-alternating cycle if its edges appear alternately in M and $E(G) \setminus M$. If G has not M-alternating cycles, then M is a unique perfect matching since the symmetric difference of two distinct perfect matchings is the union of some M-alternating cycles. So M is a unique perfect matching of G if and only if G has no M-alternating cycles. Lei et al. obtained the following characterization for an anti-forcing set of a perfect matching.

Lemma 2.1 (Lei et al. (2016)). A set $S \subseteq E(G) \setminus M$ is an anti-forcing set of M if and only if S contains at least one edge of every M-alternating cycle of G.

A compatible *M*-alternating set of *G* is a set of *M*-alternating cycles such that any two members are either disjoint or intersect only at edges in *M*. Let c'(M) denote the maximum cardinality of compatible *M*-alternating sets of *G*. By Lemma 2.1, the authors obtained the following theorem.

Theorem 2.2 (Lei et al. (2016)). For any perfect matching M of G, we have $af(G, M) \ge c'(M)$.



Fig. 1. A perfect matching M of Q_3 (thick edges) and an anti-forcing set S of M (" \times ").

In general, for any anti-forcing set S of a perfect matching M of G, the edge set $E(G) \setminus (M \cup S)$ may not be an anti-forcing set of M (see Fig. 1). However, for any minimal anti-forcing set in a bipartite graph, we have Lemma 2.3. Here an anti-forcing set is *minimal* if its any proper subset is not an anti-forcing set. Recall that for an edge subset E of a graph G, G[E] is an edge induced subgraph of G with vertex set being the vertices incident with some edge of E and edge set being E.

Lemma 2.3. Let G be a simple bipartite graph with a perfect matching M, and S a minimal anti-forcing set of M. Then $S^* := E(G) \setminus (M \cup S)$ is an anti-forcing set of M.

Proof: Clearly, M is a perfect matching of $G[M \cup S]$. It is sufficient to show that $G[M \cup S]$ has no M-alternating cycle by Lemma 2.1. By the contrary, we suppose that C is an M-alternating cycle of $G[M \cup S]$. Then the edges of C appear alternately in M and S. Let $E(C) \cap S = \{e_1, e_2, \ldots, e_k\}$ (see Fig. 2 for k = 3). Since S is a minimal anti-forcing set of M in G, the subgraph $G - (S \setminus \{e_i\})$ has an M-alternating cycle C_i such that $E(C_i) \cap S = \{e_i\}$, $i = 1, 2, \ldots, k$. Then G - S has a closed M-alternating walk $W = G[\bigcup_{i=1}^{k} (E(C_i) \setminus \{e_i\})$ as depicted in Fig. 2. Since G is a bipartite graph, W contains an M-alternating cycle C'. So G - S has an M-alternating cycle C'. This implies that S is not an anti-forcing set of M, a contradiction. So S^* is an anti-forcing set of M.

Let X and Y be two vertex subsets of a graph G. We denote by E(X, Y) the set of edges of G with one end in X and the other end in Y. The subgraph induced by E(X, Y), for convenience, is denoted by G(X, Y). For a vertex

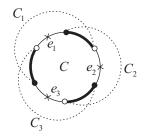


Fig. 2. Example of k = 3.

subset X of G, G[X] is a vertex induced subgraph of G with vertex set X and any two vertices are adjacent if and only if they are adjacent in G. The edge set of G[X] is denoted by E(X).

Proof of Theorem 1.2. For any perfect matching M of G, let A be a vertex subset of G consisting of one end vertex for each edge of M, $\overline{A} := V(G) \setminus A$. Then $G' := G(A, \overline{A})$ is a bipartite graph and M is a perfect matching of G'. Let S be a minimum anti-forcing set of M in G'. By Lemma 2.3, $S^* := E(G') \setminus (M \cup S)$ is an anti-forcing set of M in G'. So both $S \cup E(A)$ and $S^* \cup E(\overline{A})$ are anti-forcing sets of M in G. Hence

$$2af(G,M) \le |S \cup E(A)| + |S^* \cup E(\bar{A})| = e(G) - |M| = e(G) - \frac{v(G)}{2}.$$

Then $af(G, M) \leq \frac{2e(G)-v(G)}{4}$. By the arbitrariness of M, $Af(G) \leq \frac{2e(G)-v(G)}{4}$.

For any perfect matching M of a complete bipartite graph $K_{m,m}$ $(m \ge 2)$, any two edges of M belong to an M-alternating 4-cycle. Since any two distinct M-alternating 4-cycles are compatible, $c'(M) \ge {m \choose 2} = \frac{m^2 - m}{2}$. By Theorems 2.2 and 1.2, we obtain $af(K_{m,m}, M) = \frac{m^2 - m}{2} = Af(K_{m,m})$. Let M' be any perfect matching of a complete graph K_{2n} . For any two edges e_1 and e_2 of M', there are two distinct M'-alternating 4-cycles each of which simultaneously contains edges e_1 and e_2 . So $af(K_{2n}, M') \ge c'(M') \ge {n \choose 2} \times 2 = n^2 - n$. By Theorem 1.2, we know that $af(K_{2n}, M') = Af(K_{2n}) = n^2 - n$. Hence every perfect matching of $K_{m,m}$ and K_{2n} is nice.

Recall that the *n*-dimensional hypercube Q_n is the graph with vertex set being the set of all 0-1 sequences of length n and two vertices are adjacent if and only if they differ in exactly one position. For $x \in \{0, 1\}$, set $\bar{x} := 1 - x$. The edge connecting the two vertices $x_1 \cdots x_{i-1} x_i x_{i+1} \cdots x_n$ and $x_1 \cdots x_{i-1} \bar{x}_i x_{i+1} \cdots x_n$ of Q_n is called an *i*-edge of Q_n . We denote by E_i the set of all the *i*-edges of Q_n , i = 1, 2, ..., n. In fact, E_i is a Θ_{Q_n} -class of Q_n . We can show the following result for Q_n .

Lemma 2.4. Θ_{Q_n} -class E_i of Q_n is a nice perfect matching, that is, $af(Q_n, E_i) = Af(Q_n) = (n-1)2^{n-2}$.

Proof: It is sufficient to discuss E_1 . Clearly, E_1 is a perfect matching of Q_n . For vertices $x = x_1 x_2 \cdots x_n$ and $y = \bar{x}_1 x_2 \cdots x_n$, the edge $xy \in E_1$ belongs to n-1 E_1 -alternating 4-cycles. Over all edges of E_1 , since each E_1 -alternating 4-cycle is counted twice, there are $\frac{(n-1)2^{n-1}}{2} = (n-1)2^{n-2}$ distinct E_1 -alternating 4-cycles in Q_n . Since any two distinct E_1 -alternating 4-cycles are compatible, $c'(E_1) \ge (n-1)2^{n-2}$. So $af(Q_n, E_1) \ge c'(E_1) \ge (n-1)2^{n-2}$ by Theorem 2.2. Since $Af(Q_n) \le (n-1)2^{n-2}$ by Theorem 1.2, $af(Q_n, E_1) = Af(Q_n) = (n-1)2^{n-2}$.

The above three examples show that the upper bound in Theorem 1.2 is tight.

2.2 Nice perfect matchings

In the following, we will characterize the nice perfect matchings of a graph. The set of neighbors of a vertex v in G is denoted by $N_G(v)$. The degree of a vertex v is the cardinality of $N_G(v)$, denoted by $d_G(v)$.

Theorem 2.5. For any perfect matching M of a simple graph G, M is a nice perfect matching of G if and only if for any two edges $e_1 = xy$ and $e_2 = uv$ of M, $xu \in E(G)$ if and only if $yv \in E(G)$, and $xv \in E(G)$ if and only if $yu \in E(G)$.

Proof: Here we only need to consider simple connected graphs. To show the sufficiency, we firstly estimate the value of c'(M) for such perfect matching M of G. Let $c'_{wz}(M)$ be the number of M-alternating 4-cycles that contain edge wz. Since for any two edges $e_1 = xy$ and $e_2 = uv$ of M, $xu \in E(G)$ if and only if $yv \in E(G)$, and $xv \in E(G)$ if and only if $yu \in E(G)$, $c'_{wz}(M) = d_G(w) - 1 = d_G(z) - 1$ for any edge wz of M. Obviously, any two distinct M-alternating 4-cycles are compatible. Then

$$c'(M) \ge \frac{\sum_{wz \in M} c'_{wz}(M)}{2}$$

$$= \frac{\sum_{wz \in M} \frac{1}{2} [(d_G(w) - 1) + (d_G(z) - 1)]}{2}$$

$$= \frac{\sum_{w \in V(G)} \frac{1}{2} (d_G(w) - 1)}{2}$$

$$= \frac{e(G) - \frac{v(G)}{2}}{2}.$$
(2)

By Theorems 1.2 and 2.2, $c'(M) \le af(G, M) \le Af(G) \le \frac{2e(G)-v(G)}{4}$. So $af(G, M) = \frac{2e(G)-v(G)}{4}$, that is, M is a nice perfect matching of G.

Conversely, suppose that M is a nice perfect matching of G. Let A be a vertex subset of G consisting of one end vertex for each edge of M and $\overline{A} := V(G) \setminus A$. Then (A, \overline{A}) is a partition of V(G). Given any bijection $\omega : M \to \{1, \ldots, |M|\}$, we extend weight function ω on M to the vertices of G: if $v \in V(G)$ is incident with $e \in M$, then $\omega(v) := \omega(e)$. This weight function ω gives a natural ordering of the vertices in $A(\overline{A})$. Clearly, if $e = xy \in M$, then $\omega(x) = \omega(y)$, otherwise, $\omega(x) \neq \omega(y)$. Set

$$\begin{split} E^{\omega}_A &:= \{ xy \in E(G) : \omega(x) > \omega(y), x \in A \text{ and } y \in \bar{A} \}, \\ E^{\omega}_{\bar{A}} &:= \{ xy \in E(G) : \omega(x) < \omega(y), x \in A \text{ and } y \in \bar{A} \}. \end{split}$$

Since $G - E_A^{\omega} \cup E(A)$ has a unique perfect matching M, $E_A^{\omega} \cup E(A)$ is an anti-forcing set of M in G. Similarly, $E_{\bar{A}}^{\omega} \cup E(\bar{A})$ is also an anti-forcing set of M in G. Since M is a nice perfect matching of G, $af(G, M) = \frac{2e(G) - v(G)}{4}$. So $|E_A^{\omega} \cup E(A)| \ge \frac{2e(G) - v(G)}{4}$, $|E_{\bar{A}}^{\omega} \cup E(\bar{A})| \ge \frac{2e(G) - v(G)}{4}$. Since $|E_A^{\omega} \cup E(A)| + |E_{\bar{A}}^{\omega} \cup E(\bar{A})| = e(G) - |M| = e(G) - \frac{v(G)}{2}$, $|E_A^{\omega} \cup E(A)| = |E_{\bar{A}}^{\omega} \cup E(\bar{A})| = \frac{2e(G) - v(G)}{4}$. Hence $E_A^{\omega} \cup E(A)$ is a minimum anti-forcing set of M in G.

Now we show that for any two edges $e_1 = xy$ and $e_2 = uv$ of M, $xu \in E(G)$ if and only if $yv \in E(G)$, and $xv \in E(G)$ if and only if $yu \in E(G)$. It is sufficient to show that $xv \in E(G)$ implies $yu \in E(G)$. Given two bijections $\omega_1 : M \to \{1, \ldots, |M|\}$ and $\omega_2 : M \to \{1, \ldots, |M|\}$ with $\omega_1(e_1) = 1$, $\omega_1(e_2) = 2$, $\omega_2(e_1) = 2$, $\omega_2(e_2) = 1$ and $\omega_2|_{M \setminus \{e_1, e_2\}} = \omega_1|_{M \setminus \{e_1, e_2\}}$. As the above extension of ω , we extend the weight functions ω_1 and ω_2 on M to the vertices of G.

We first consider the case that $x, u \in A$. Suppose to the contrary that $xv \in E(G)$ but $yu \notin E(G)$. Set $A' := A \setminus \{x, u\}, \bar{A}' := \bar{A} \setminus \{y, v\}, E'_1 := \{wz \in E(G) : \omega_1(w) > \omega_1(z), w \in A' \text{ and } z \in \bar{A}'\}$. Then

$$E_A^{\omega_2} \cup E(A) = \{xv\} \cup E(\{y,v\}, A') \cup E'_1 \cup E(A) = \{xv\} \cup E_A^{\omega_1} \cup E(A).$$
(3)

By the above proof we know that both $E_A^{\omega_1} \cup E(A)$ and $E_A^{\omega_2} \cup E(A)$ are minimum anti-forcing sets of M in G, it contradicts to the equation (3). Thus $yu \in E(G)$.

For the case that $x \in A$ and $u \in \overline{A}$, set $U := (A \setminus \{v\}) \cup \{u\}, \overline{U} := (\overline{A} \setminus \{u\}) \cup \{v\}$. Then each edge in M is incident with exactly one vertex in U. Substituting the partition (A, \overline{A}) of V(G) with the partition (U, \overline{U}) , by a similar argument as the above case, we can also show that $xv \in E(G)$ implies $yu \in E(G)$.

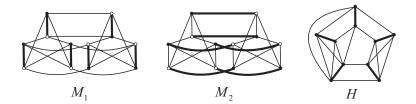


Fig. 3. Two nice perfect matchings M_1 and M_2 of G' and a nice perfect matching of H.

By Theorem 2.5, we can easily check whether a perfect matching of a graph is nice. For example, in Fig. 3, the two perfect matchings M_1 and M_2 of the bipartite graph G' are nice, and the perfect matching of the non-bipartite graph H is also nice.

Proposition 2.6. Let M be a nice perfect matching of G and S a subset of V(G). Then $M \cap E(S)$ is a nice perfect matching of G[S] if $M \cap E(S)$ is a perfect matching of G[S].

Proof: By Theorem 2.5, it holds.

In the proof of Theorem 2.5, we notice that $d_G(u) = d_G(v)$ for every edge uv of a nice perfect matching of G. So we have the following necessary but not sufficiency condition for the upper bound in Theorem 1.2 to be attained.

Proposition 2.7. Let G be a graph with a perfect matching. Then $Af(G) < \frac{2e(G)-v(G)}{4}$ if there are an odd number of vertices of the same degree in G.

Proposition 2.7 is not sufficient. For example, for a hexagonal system with a perfect matching, it does not have a nice perfect matching by Theorem 2.5, that is, its maximum anti-forcing number can not be the upper bound in Theorem 1.2, but it has an even number of vertices of degree 3 and an even number of vertices of degree 2.

Abay-Asmerom et al. (2010) introduced a *reversing involution* of a connected bipartite graph G with partite sets X and Y as an automorphism α of G of order two such that $\alpha(X) = Y$ and $\alpha(Y) = X$. Here we give the following definition of a general graph.

Definition 2.8. Suppose that G is a simple connected graph. An edge-involution of G is an automorphism α of G of order two such that v and $\alpha(v)$ are adjacent for any vertex v in G.

Hence an edge-involution of a bipartite graph is also a reversing involution, but a reversing involution of a bipartite graph may not be an edge-involution. In the following, we establish a relationship between a nice perfect matching and an edge-involution of G.

Theorem 2.9. Let G be a simple connected graph. Then there is a one-to-one correspondence between the nice perfect matchings of G and the edge-involutions of G.

Proof: For a nice perfect matching M of G, we define a bijection α_M of order 2 on V(G) as follows: for any vertex v of G, there is exactly one edge e in M such that v is incident with e, let $\alpha_M(v)$ be the other end-vertex of e. Let x and y be any two distinct vertices of G. If $xy \in M$, then $\alpha_M(x) = y$, $\alpha_M(y) = x$ and $\alpha_M(x)\alpha_M(y) = yx \in E(G)$. If $xy \notin M$ (x may not be adjacent to y), then both $x\alpha_M(x)$ and $y\alpha_M(y)$ belong to M. Since M is a nice perfect matching, $xy \in E(G)$ if and only if $\alpha_M(x)\alpha_M(y) \in E(G)$ by Theorem 2.5. This implies that α_M is an automorphism of G. Thus α_M is an edge-involution of G.

Conversely, let α be an edge-involution of G. Then for any vertex y of G, $y\alpha(y) \in E(G)$. Since α is a bijection of order 2 on V(G), $M' := \{y\alpha(y) : y \in V(G)\}$ is a perfect matching of G. For any two distinct edges $y_1\alpha(y_1)$ and $y_2\alpha(y_2)$ of M', $y_1y_2 \in E(G)$ if and only if $\alpha(y_1)\alpha(y_2) \in E(G)$, and $y_1\alpha(y_2) \in E(G)$ if and only if $\alpha(y_1)y_2 \in E(G)$ since α is an automorphism of order 2 of G. So M' is a nice perfect matching of G by Theorem 2.5. We can also see that $\alpha_{M'} = \alpha$. This establishes a one-to-one correspondence between the nice perfect matchings of G and the edge-involutions of G.

3 Construction of the extremal graphs

In the following, we will show that every extremal graph can be constructed from a complete graph K_2 by implementing two expansion operations.

Definition 3.1. Let G_i be a simple graph with a nice perfect matching M_i , i = 1, 2 (note that $V(G_1) \cap V(G_2) = \emptyset$). We define two expansion operations as follows:

(i) $G := G_i + e + e'$, where $e, e' \notin E(G_i)$ and there are edges $e_1, e_2 \in M_i$ such that the four edges e, e', e_1, e_2 form a 4-cycle.

(ii) For $M'_1 \subseteq M_1$ and $M'_2 \subseteq M_2$ with $|M'_1| = |M'_2|$, given a bijection ϕ from $V(M'_1)$ to $V(M'_2)$ with $uv \in M'_1$ if and only if $\phi(u)\phi(v) \in M'_2$. G_1 joins G_2 over matchings M'_1 and M'_2 about bijection ϕ , denoted by $G_1 \circledast G_2$, is a graph with vertex set $V(G_1) \cup V(G_2)$ and edge set $E(G_1) \cup E(G_2) \cup E'$, where $E' := \{u\phi(u) : u \in V(M'_1)\}$.

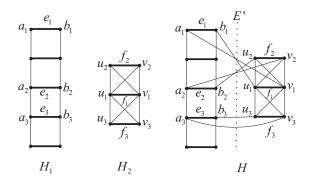


Fig. 4. $H = H_1 \circledast H_2$ over matchings M'_1 and M'_2 about bijection ϕ' .

For example, in Fig. 4 graph H is $H_1 \otimes H_2$ over matchings M'_1 of H_1 and M'_2 of H_2 about bijection ϕ' , where $M'_1 = \{e_1, e_2, e_3\}, M'_2 = \{f_1, f_2, f_3\}, \phi'(a_i) = v_i, \phi'(b_i) = u_i, i = 1, 2, 3$. H has a nice perfect matching which is marked by thick edges in Fig. 4. Recall that nK_2 is the disjoint union of n copies of K_2 .

Theorem 3.2. A simple graph G is an extremal graph if and only if it can be constructed from K_2 by implementing operations (i) or (ii) in Definition 3.1 (regardless of the orders).

Proof: Let \mathcal{P}' be the set of all the graphs that can be constructed from K_2 by implementing operations (*i*) or (*ii*). For any graph $G \in \mathcal{P}'$, G is a simple graph with a nice perfect matching by the definition of the two operations.

Conversely, we suppose that G is an extremal graph, that is, G has a nice perfect matching $M = \{e_1, e_2, \ldots, e_n\}$. If n = 1, or 2, then G must be isomorphic to K_2 , $2K_2$, C_4 or K_4 . So $G \in \mathcal{P}'$. Next, we suppose that $n \ge 3$ and it holds for n - 1. Let $G' := G[\bigcup_{i=1}^{n-1} V(e_i)]$. Then $\{e_1, \ldots, e_{n-1}\}$ is a nice perfect matching of G' by Proposition 2.6. So $G' \in \mathcal{P}'$ by the induction. If e_n is an isolated edge in G, then $G = G' \cup \{e_n\} \in \mathcal{P}'$. Otherwise, $e_n = u_n v_n$ has adjacent edges $u_n v_i$ and $v_n u_i$, or $u_n u_i$ and $v_n v_i$ for some $i \in \{1, \ldots, n-1\}$, where $u_i v_i = e_i \in M$. Let $G'' = G' \circledast K_2$ over matchings $\{e_i\}$ and $\{e_n\}$ about bijection $\phi : \{u_i, v_i\} \to \{u_n, v_n\}$. So $G'' \in \mathcal{P}'$. Then G can be constructed from G'' by implementing several times operations (i). So $G \in \mathcal{P}'$.

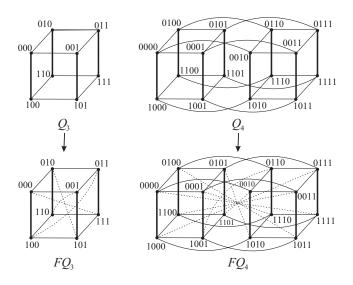


Fig. 5. The nice perfect matchings E_1 of Q_3 and Q_4 are depicted by thick edges; the dashed edges are the complementary edges.

As a variant of the *n*-dimensional hypercube Q_n , the *n*-dimensional folded hypercube FQ_n , proposed first by El-Amawy and Latifi (1991), is a graph with $V(FQ_n) = V(Q_n)$ and $E(FQ_n) = E(Q_n) \cup \overline{E}$, where $\overline{E} := \{x\overline{x} : x = x_1x_2\cdots x_n, \overline{x} = \overline{x}_1\overline{x}_2\cdots\overline{x}_n\}, \overline{x}_i := 1 - x_i$. Each edge in \overline{E} is called a *complementary edge*. The graphs shown in Fig. 5 are FQ_3 and FQ_4 , respectively.

Corollary 3.3. FQ_n is an extremal graph and $Af(FQ_n) = n2^{n-2}$.

Proof: By Lemma 2.4, E_1 is a nice perfect matching of Q_n . FQ_n is constructed from Q_n by applying the operation (*i*) over the nice perfect matching E_1 of Q_n (see Fig. 5 for n = 3, 4). So E_1 is also a nice perfect matching of the folded hypercube FQ_n .

For any positive integer n, a connected graph G with at least 2n + 2 vertices is said to be *n*-extendable if every matching of size n is contained in a perfect matching of G.

Proposition 3.4. Any connected extremal graph G other than K_2 is 1-extendable.

Proof: Since G is an extremal graph, it has a nice perfect matching M. For any edge uv of $E(G) \setminus M$, there are edges ux and vy of M. By Theorem 2.5, $xy \in E(G)$. So uv belongs to an M-alternating 4-cycle C := uxyvu. Then $M \triangle E(C) := (M \cup E(C)) \setminus (M \cap E(C))$ is a perfect matching of G that contains edge uv. So G is 1-extendable. \Box

By Proposition 3.4, any connected extremal graph except for K_2 is 2-connected.

4 Cartesian decomposition

The Cartesian product $G \Box H$ of two graphs G and H is a graph with vertex set $V(G) \times V(H) = \{(x, u) : x \in V(G), u \in V(H)\}$ and two vertices (x, u) and (y, v) are adjacent if and only if $xy \in E(G)$ and u = v or x = y and $uv \in E(H)$. For a vertex (x_i, v_j) of $G \Box H$, the subgraphs of $G \Box H$ induced by the vertex set $\{(x, v_j) : x \in V(G)\}$ and the vertex set $\{(x_i, v) : v \in V(H)\}$ are called a G-layer and an H-layer of $G \Box H$, and denoted by G^{v_j} and H^{x_i} , respectively.

For any graph H, let E' be the set of edges of all K_2 -layers of $H \Box K_2$. Clearly, E' is a perfect matching of $H \Box K_2$. Define a bijection α on $V(H \Box K_2)$ as follows: for every edge $uv \in E'$, $\alpha(u) := v$ and $\alpha(v) := u$. Then α is an edge-involution of $H \Box K_2$. So $H \Box K_2$ is an extremal graph by Theorem 2.9. This fact inspires us to consider the Cartesian product decomposition of an extremal graph. Let $\Phi^*(G)$ be the number of all the nice perfect matchings of a graph G. We have Theorem 4.1. Recall that for an edge e = uv of G and an isomorphism φ from G to H, $\varphi(e) := \varphi(u)\varphi(v)$.

Theorem 4.1. Let G_1 and G_2 be two simple connected graphs. Then

$$\Phi^*(G_1 \square G_2) = \Phi^*(G_1) + \Phi^*(G_2).$$

Proof: Let $V(G_1) = \{x_1, x_2, \dots, x_{n_1}\}$ and $V(G_2) = \{v_1, v_2, \dots, v_{n_2}\}$. Since $\Phi^*(K_1) = 0$, we suppose $n_1 \ge 2$ and $n_2 \ge 2$.

We define an isomorphism ρ_{v_i} from G_1 to $G_1^{v_i}$ and an isomorphism σ_{x_j} from G_2 to $G_2^{x_j}$: $\rho_{v_i}(x) := (x, v_i)$ for any vertex x of G_1 and $\sigma_{x_i}(v) := (x_j, v)$ for any vertex v of G_2 . For any nice perfect matching M_i of G_i , i = 1, 2, let

$$\rho(M_1) := \bigcup_{v_i \in V(G_2)} \rho_{v_i}(M_1), \quad \sigma(M_2) := \bigcup_{x_j \in V(G_1)} \sigma_{x_j}(M_2), \tag{4}$$

By Theorem 2.5, $\rho_{v_i}(M_1)$ is a nice perfect matching of $G_1^{v_i}$ and $\rho(M_1)$ is a nice perfect matching of $G_1 \square G_2$. Similarly, $\sigma(M_2)$ is also a nice perfect matching of $G_1 \square G_2$.

Conversely, since $E(G_1^{v_1}), \ldots, E(G_1^{v_{n_2}}), E(G_2^{x_1}), \ldots, E(G_2^{x_{n_1}})$ is a partition of $E(G_1 \square G_2)$, for any nice perfect matching M of $G_1 \square G_2$ there is some x_i or v_j such that $M \cap E(G_2^{x_i}) \neq \emptyset$ or $M \cap E(G_1^{v_j}) \neq \emptyset$. If $M \cap E(G_2^{x_i}) \neq \emptyset$ for some x_i , then we have the following Claim.

Claim: $M \cap E(G_2^{x_j})$ is a nice perfect matching of $G_2^{x_j}$ for each $x_j \in V(G_1)$, and $\sigma_{x_j}^{-1}(M \cap E(G_2^{x_j})) = \sigma_{x_i}^{-1}(M \cap E(G_2^{x_j}))$. $E(G_2^{x_i})$). So $M \cap E(G_1^v) = \emptyset$ for each $v \in V(G_2)$.

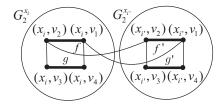


Fig. 6. Illustration for the proof of the Claim in Theorem 4.1.

Take an edge $f = (x_i, v_1)(x_i, v_2) \in M \cap E(G_2^{x_i})$. Then $v_1v_2 \in E(G_2)$. If $n_2 = 2$, then $M \cap E(G_2^{x_i}) = \{f\}$ is a nice perfect matching of $G_2^{x_i}$. For $n_2 \geq 3$, since G_2 is connected, without loss of generality we may assume that $d_{G_2}(v_2) \geq 2$. Let v_3 be a neighbor of v_2 that is different from v_1 . So $(x_i, v_2)(x_i, v_3) \in E(G_2^{x_i})$. Let g be an edge of M with an end-vertex (x_i, v_3) . Since M is a nice perfect matching of G, the other end-vertex of g must be adjacent to (x_i, v_1) by Theorem 2.5. So the other end-vertex of g belongs to $V(G_2^{x_i})$ (see Fig. 6), that is, $g \in E(G_2^{x_i})$. Since $G_2^{x_i} \cong G_2$ is a connected graph, we can obtain that $M \cap E(G_2^{x_i})$ is a perfect matching of $G_2^{x_i}$ in the above way. So $M \cap E(G_2^{x_i})$ is a nice perfect matching of $G_2^{x_i}$ by Proposition 2.6.

Since G_1 is connected and $n_1 \ge 2$, there is some vertex $x_{i'}$ of G_1 such that x_i and $x_{i'}$ are adjacent in G_1 . So vertex $(x_{i'}, v_1) \notin G_2^{x_i}$ is adjacent to (x_i, v_1) in $G_1 \square G_2$ (see Fig. 6). Let f' be an edge of M that is incident with $(x_{i'}, v_1)$. Since M is a nice perfect matching of $G_1 \square G_2$, the other end-vertex of f' must be adjacent to (x_i, v_2) by Theorem 2.5. So $f' = (x_{i'}, v_1)(x_{i'}, v_2) \in M \cap E(G_2^{x_{i'}})$. As the above proof, we can similarly show that $M \cap E(G_2^{x_{i'}})$ is a nice perfect matching of $G_2^{x_{i'}}$. Since G_1 is connected, in an inductive way we can show that $M \cap E(G_2^{x_j})$ is a nice perfect matching of $G_2^{x_j}$ for any $x_j \in V(G_1)$.

Notice that $\sigma_{x_i}^{-1}(f) = v_1 v_2 = \sigma_{x'_i}^{-1}(f')$. Let g' be the edge of M that is incident with $(x_{i'}, v_3)$. Since $(x_{i'}, v_3)$ is adjacent to (x_i, v_3) , the other end vertex of g' must be adjacent to the other end vertex (x_i, v_4) of g by Theorem 2.5. So

 $g' = (x_{i'}, v_3)(x_{i'}, v_4) \text{ since } g' \in E(G_2^{x_{i'}}). \text{ This implies that } \sigma_{x_i'}^{-1}(g') = v_3v_4 = \sigma_{x_i}^{-1}(g). \text{ In an inductive way, we can show that } \sigma_{x_{i'}}^{-1}(M \cap E(G_2^{x_{i'}})) = \sigma_{x_i}^{-1}(M \cap E(G_2^{x_i})). \text{ Similarly, we also have } \sigma_{x_j}^{-1}(M \cap E(G_2^{x_j})) = \sigma_{x_i}^{-1}(M \cap E(G_2^{x_i})) \text{ for any } x_j \in V(G_1).$

By this Claim, $M_2 := \sigma_{x_i}^{-1}(M \cap E(G_2^{x_i}))$ is a nice perfect matching of G_2 with $M = \sigma(M_2)$. If $M \cap E(G_1^{v_j}) \neq \emptyset$, then we can similarly show that G_1 has a nice perfect matching M_1 with $M = \rho(M_1)$. So $\Phi^*(G_1 \square G_2) = \Phi^*(G_1) + \Phi^*(G_2)$.

In fact, we can get the following corollary.

Corollary 4.2. Let G be a simple connected graph. Then we have $\Phi^*(G) = \sum_{i=1}^k \Phi^*(G_i)$ for any decomposition $G_1 \Box \cdots \Box G_k$ of G.

Now, it is easy to get the following proposition.

Proposition 4.3. A simple connected graph G is an extremal graph if and only if one of its Cartesian product factors is an extremal graph.

The *n*-dimensional enhanced hypercube $Q_{n,k}$, see Tzeng and Wei (1991), is the graph with vertex set $V(Q_{n,k}) = V(Q_n)$ and edge set $E(Q_{n,k}) = E(Q_n) \cup \{(x_1x_2 \cdots x_{n-1}x_n, \bar{x}_1\bar{x}_2 \cdots \bar{x}_{n-k-1}\bar{x}_{n-k}x_{n-k+1}x_{n-k+2} \cdots x_n : x_1x_2 \cdots x_n \in V(Q_{n,k})\}$, where $0 \le k \le n-1$. Clearly, $Q_n \cong Q_{n,n-1}$ and $FQ_n \cong Q_{n,0}$, i.e., the hypercube and the folded hypercube are regarded as two special cases of the enhanced hypercube. By Yang et al. (2015a), we have $Q_{n,k} \cong FQ_{n-k} \Box Q_k$, for $0 \le k \le n-1$. Hence we obtain the following result by the Proposition 4.3.

Corollary 4.4. $Q_{n,k}$ is an extremal graph and $Af(Q_{n,k}) = n2^{n-2}$.

According to the above discussion, for any graph G, we know that $K_{m,m} \Box G$, $K_{2n} \Box G$, $Q_n \Box G$, $FQ_n \Box G$ and $Q_{n,k} \Box G$ are extremal graphs. Moreover, we can produce an infinite number of extremal graphs from an extremal graph by the Cartesian product operation.

5 Further applications

From examples we already know that $K_{m,m}$, K_{2n} , Q_n , FQ_n and $Q_{n,k}$ are extremal graphs. Two perfect matchings M_1 and M_2 of a graph G are called *equivalent* if there is an automorphism φ of G such that $\varphi(M_1) = M_2$. So we know that all the perfect matchings of $K_{m,m}$ (or K_{2n}) are nice and equivalent. Further in this section we will count nice perfect matchings of the three cube-like graphs.

Theorem 5.1. Q_n has exactly *n* nice perfect matchings E_1, E_2, \ldots, E_n , all of which are equivalent.

Proof: By Lemma 2.4, E_1, E_2, \ldots, E_n are *n* distinct nice perfect matchings of Q_n . Since Q_n is the Cartesian product of n K_2 's, Q_n has exactly *n* nice perfect matchings by Corollary 4.2. So the first part is done. Now, it remains to show that E_i and E_j are equivalent for any $1 \le i < j \le n$. Let the automorphism f_{ij} of Q_n be defined as $f_{ij}(x_1 \cdots x_{i-1}x_ix_{i+1} \dots x_{j-1}x_jx_{j+1} \cdots x_n) = x_1 \cdots x_{i-1}x_jx_{i+1} \dots x_{j-1}x_ix_{j+1} \cdots x_n$ for each vertex $x_1x_2 \cdots x_n$ of Q_n . Then $f_{ij}(E_i) = E_j$.

The theorem can be obtained by applying the reversing-involutions of bipartite graphs, see Abay-Asmerom et al. (2010), but the computation is tedious.

Since $FQ_2 \cong K_4$ and $FQ_3 \cong K_{4,4}$, we have $\Phi^*(FQ_2) = 3$ and $\Phi^*(FQ_3) = 24$. For $n \ge 4$, we have a general result as follows.

Theorem 5.2. FQ_n has exactly n + 1 nice perfect matchings for $n \ge 4$.

Proof: By Lemma 2.4, E_i is a perfect matching of Q_n . Then E_i is also a perfect matching of FQ_n . We can easily check that E_i is a nice perfect matching of FQ_n by Theorem 2.5.

Let E_{n+1} be the set of all the complementary edges of FQ_n . Then E_{n+1} is a perfect matching of FQ_n . Let $u\bar{u}$ and $v\bar{v}$ be two distinct edges in E_{n+1} . Since any two distinct complementary edges are independent, the edge linked u to v or \bar{v} (if exist) does not belong to E_{n+1} . We can easily show that $uv \in E_j$ if and only if $\bar{u}\bar{v} \in E_j$ for some j = 1, 2, ..., n, and $u\bar{v} \in E_s$ if and only if $\bar{u}v \in E_s$ for some s = 1, 2, ..., n. So E_{n+1} is also a nice perfect matching of FQ_n .

Now, we have found n + 1 nice perfect matchings of FQ_n . Next, we will show that FQ_n has no other nice perfect matchings. By the contrary, we suppose that M is a nice perfect matching of FQ_n that is different from any $E_i, i = 1, 2, \ldots, n+1$. Since E_1, \ldots, E_{n+1} is a partition of the edge set $E(FQ_n)$, there is E_k with $k \neq n+1$ such that $M \cap E_k \neq \emptyset$ and $E_k \neq M$. Clearly, $FQ_n - (E_{n+1} \cup E_k)$ has exactly two components both of which are isomorphic to Q_{n-1} . We notice that the k-th coordinate of each vertex in one component is 0, and 1 in the other component. We denote the two components by Q_n^0 and Q_n^1 , respectively. In fact, $V(Q_n^i) = \{x_1 \cdots x_{k-1} i x_{k+1} \cdots x_n :$ $x_j = 0$ or $1, j = 1, \dots, k - 1, k + 1, \dots, n$, i = 0, 1. Since $M \cap E_k \neq \emptyset$, there is some edge $vv' \in M \cap E_k$ with $v \in V(Q_n^0)$ and $v' \in V(Q_n^1)$. For any vertex w of Q_n^0 with w and v being adjacent, we consider the edge g of M that is incident with w. By Theorem 2.5, the other end-vertex of g is adjacent to v'. If $g = w\bar{w}$ is a complementary edge of FQ_n , then there are exactly two same bits in the strings of \bar{w} and v'. So the edge $\bar{w}v' \in E(FQ_n)$ is not a complementary edge of FQ_n . Since \bar{w} and v' are adjacent, there is exactly one different bit in the strings of \bar{w} and v'. So n = 3, a contradiction. If $g = wz \in E(Q_n^0)$, then there are exactly three different bits in the strings of z and v'. Since z and v' are adjacent in FQ_n , the edge zv' is a complementary edge of FQ_n . So n = 3, a contradiction. Hence $g \in E_k$. Since Q_n^0 is connected, using the above method repeatedly, we can show that $M = E_k$, a contradiction. So FQ_n has exactly n + 1 nice perfect matchings.

Proposition 5.3. All the nice perfect matchings of FQ_n $(n \ge 2)$ are equivalent.

Proof: We notice that $FQ_2 \cong K_4$ and $FQ_3 \cong K_{4,4}$. So all the nice perfect matchings of FQ_n are equivalent for $2 \leq n \leq 3$. Suppose that $n \geq 4$. From the proof of Theorem 5.2 we know that $E_1, E_2, \ldots, E_{n+1}$ are all the nice perfect matchings of FQ_n . f_{ij} defined in the proof of Theorem 5.1 is also an automorphism of FQ_n such that $\varphi(E_i) = E_j$ for $1 \leq i < j \leq n$. We will show that E_1 and E_{n+1} are equivalent. Clearly, $FQ_n - (E_1 \cup E_{n+1})$ has exactly two components each isomorphic to Q_{n-1} , denoted by Q_n^0 and Q_n^1 . Set $V(Q_n^i) = \{ix_2x_3 \cdots x_n : x_j = 0 \text{ or } 1, j = 2, \ldots, n\}$, i = 0, 1. We define a bijection f on $V(FQ_n)$ as follows:

$$f(x_1x_2\cdots x_n) = \begin{cases} \bar{x}_1x_2\cdots x_n, & \text{if } x_1x_2\cdots x_n \in V(Q_n^0), \\ \bar{x}_1\bar{x}_2\cdots \bar{x}_n, & \text{if } x_1x_2\cdots x_n \in V(Q_n^1). \end{cases}$$

It is easy to check that f is an automorphism of FQ_n . In addition, $f(E_1) = E_{n+1}$. Hence all the nice perfect matchings of FQ_n are equivalent.

By Corollary 4.2 and Theorems 5.1 and 5.2, we can obtain the following conclusion.

Corollary 5.4. $\Phi^*(Q_{n,n-1}) = n$, $\Phi^*(Q_{n,n-2}) = n + 1$, $\Phi^*(Q_{n,n-3}) = n + 21$ and $\Phi^*(Q_{n,k}) = n + 1$ for any $0 \le k \le n - 4$.

Proposition 5.5. For 0 < k < n - 1, $Q_{n,k}$ has exactly two nice perfect matchings up to the equivalent.

Proof: Since $Q_{n,k} = FQ_{n-k} \Box Q_k$, by adapting the notations in Eq. (4) and by the proof of Theorem 4.1 we know that all the nice perfect matchings of $Q_{n,k}$ are divided into two classes \mathcal{M}' and \mathcal{M}'' , where $\mathcal{M}' = \{\rho(M) : M \text{ is a nice perfect matching of } FQ_{n-k}\}$ and $\mathcal{M}'' = \{\sigma(M) : M \text{ is a nice perfect matching of } Q_k\}$.

For $M'_1, M'_2 \in \mathcal{M}'$, there are two nice perfect matchings M_1 and M_2 of FQ_{n-k} such that $M'_i = \rho(M_i), i = 1, 2$. By Proposition 5.3, there exists an automorphism φ of FQ_{n-k} such that $\varphi(M_1) = M_2$. Let $\varphi'(x, u) := (\varphi(x), u)$ for each vertex (x, u) of $FQ_{n-k} \Box Q_k$. It is easy to check that φ' is an automorphism of $Q_{n,k}$ and $\varphi'(M'_1) = M'_2$. By the arbitrariness of M'_1 and M'_2 , we know that all the nice perfect matchings in \mathcal{M}' are equivalent. Similarly, we can show that all the nice perfect matchings in \mathcal{M}'' are equivalent.

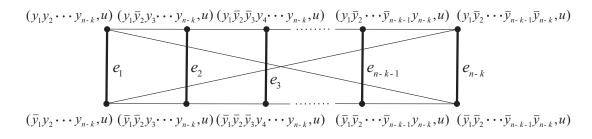


Fig. 7. The graph H.

Let F_1 and E_1 be the sets of all the 1-edges of FQ_{n-k} and Q_k respectively. Then F_1 is a nice perfect matching of FQ_{n-k} and E_1 is a nice perfect matching of Q_k . So $\rho(F_1) \in \mathcal{M}'$ and $\sigma(E_1) \in \mathcal{M}''$. See Fig. 7, we choose a subset $S := \{e_1, \ldots, e_{n-k}\}$ of $\rho(F_1)$. Then all the vertices incident with S induce a subgraph H as depicted in Fig. 7. For any subset $R \subseteq \sigma(E_1)$ of size n - k, let G be the subgraph of $Q_{n,k}$ induced by all the vertices incident with R. We note that $Q_{n,k} - \sigma(E_1)$ has exactly two components A and B each of which is isomorphic to $FQ_{n-k} \Box Q_{k-1}$, and $\sigma(E_1) = E(A, B)$. So G - R has at least two components. Clearly H - S is connected. So for any automorphism ψ of $Q_{n,k}, \psi(S) \neq R$. By the arbitrariness of R we know that $\rho(F_1)$ and $\sigma(E_1)$ are not equivalent. Then we are done. \Box

From Corollary 4.2 it is helpful to give a Cartesian decomposition of an extremal graph. It is known that $Q_n \cong K_2 \Box \cdots \Box K_2$ and $Q_{n,k} \cong FQ_{n-k} \Box Q_k$. However we shall see surprisedly that FQ_n is undecomposable.

A nontrivial graph G is said to be *prime* with respect to the Cartesian product if whenever $G \cong H \square R$, one factor is isomorphic to the complete graph K_1 and the other is isomorphic to G. Clearly, for $m \ge 3$ and $n \ge 2$, $K_{m,m}$ and K_{2n} are prime extremal graphs. In the sequel, we show that FQ_n is a prime extremal graph, too.

Recall that the length of a shortest path between two vertices x and y of G is called the distance between x and y, denoted by $d_G(x, y)$. Let G be a connected graph. Two edges e = xy and f = uv are in the relation Θ_G if $d_G(x, u) + d_G(y, v) \neq d_G(x, v) + d_G(y, u)$. Notice that Θ_G is reflexive and symmetric, but need not to be transitive. We denote its transitive closure by Θ_G^* . For an even cycle C_{2n} , $\Theta_{C_{2n}}$ consists of all pairs of antipodal edges. Hence, $\Theta_{C_{2n}}^*$ has n equivalence classes and $\Theta_{C_{2n}} = \Theta_{C_{2n}}^*$. For an odd cycle C, any edge of C is in relation Θ with its two antipodal edges. So all edges of C belong to an equivalence class with respect to Θ_C^* . By the Cartesian product decomposition Algorithm depicted in Imrich and Klavzar (2000), we have the following lemma.

Lemma 5.6. If all the edges of a graph G belong to an equivalence class with respect to Θ_G^* , then G is a prime graph under the Cartesian product.

The *Hamming distance* between two vertices x and y in Q_n is the number of different bits in the strings of both vertices, denoted by $H_{Q_n}(x, y)$.

Theorem 5.7 (Xu and Ma (2006)). For a folded hypercube FQ_n , we have

(1) FQ_n is a bipartite graph if and only if n is odd.

(2) The length of any cycle in FQ_n that contains exactly one complementary edge is at least n + 1. If n is even, then the length of a shortest odd cycle in FQ_n is n + 1.

(3) Let u and v be two vertices in FQ_n . If $H_{Q_n}(u,v) \leq \lfloor \frac{n}{2} \rfloor$, then any shortest uv-path in FQ_n contains no complementary edges. If $H_{Q_n}(u,v) > \lceil \frac{n}{2} \rceil$, then any shortest uv-path in FQ_n contains exactly one complementary edge.

Here we list some known properties of Q_n that will be used in the sequel. For any two vertices x and y in Q_n , $d_{Q_n}(x, y) = H_{Q_n}(x, y)$. For any shortest path P from $x_1 x_2 \cdots x_n$ to $\bar{x}_1 \bar{x}_2 \cdots \bar{x}_n$ in Q_n , $|E(P) \cap E_i| = 1$ for each

i = 1, 2, ..., n. For any integer j $(1 \le j \le n)$, there is a shortest path P from $x_1 x_2 \cdots x_n$ to $\bar{x}_1 \bar{x}_2 \cdots \bar{x}_n$ in Q_n such that the edge in $E(P) \cap E_i$ is the *j*th edge when traverse P from $x_1 x_2 \cdots x_n$ to $\bar{x}_1 \bar{x}_2 \cdots \bar{x}_n$.

For every subgraph F of a graph G, the inequality $d_F(u, v) \ge d_G(u, v)$ obviously holds. If $d_F(u, v) = d_G(u, v)$ for all $u, v \in V(F)$, we say F is an *isometric subgraph* of G.

Proposition 5.8 (Hammack et al. (2011)). Let C be a shortest cycle of G. Then C is isometric in G.

Theorem 5.9. FQ_n is a prime graph under the Cartesian product.

Proof: Clearly, FQ_2 and FQ_3 are prime. So we suppose that $n \ge 4$. We recall that E_i is the set of all the *i*-edges of Q_n , i = 1, 2, ..., n. Let E_{n+1} be the set of all the complementary edges of FQ_n . Then $E_1, E_2, ..., E_{n+1}$ is a partition of $E(FQ_n)$. Since the girth of FQ_n is 4 for $n \ge 4$, any two opposite edges of a 4-cycle are in relation Θ_{FQ_n} . So E_i is contained in an equivalence class with respect to $\Theta_{FQ_n}^*$, i = 1, 2, ..., n + 1. For any vertex $x_1x_2 \cdots x_n$, it is linked to $\bar{x}_1\bar{x}_2\cdots\bar{x}_n$ by a complementary edge e in FQ_n . Let P be any shortest path from $x_1x_2\cdots x_n$ to $\bar{x}_1\bar{x}_2\cdots\bar{x}_n$ in Q_n . Then the length of P is n and $|P \cap E_i| = 1$ for any i = 1, 2, ..., n. Set $C := P \cup \{e\}$. Then C is a cycle of length n + 1.

If n is even, then the length of any shortest odd cycle in FQ_n is n + 1 by Theorem 5.7 (2). So C is a shortest odd cycle in FQ_n . By Proposition 5.8, C is an isometric odd cycle in FQ_n . So all edges of C belong to an equivalence class with respect to $\Theta_{FQ_n}^*$. Since $E(C) \cap E_i \neq \emptyset$ for any i = 1, 2, ..., n + 1, all edges of $E(FQ_n) = \bigcup_{i=1}^{n+1} E_i$ belong to an equivalence class with respect to $\Theta_{FQ_n}^*$, that is, FQ_n is a prime graph under the Cartesian product by Lemma 5.6.

For *n* being odd, we first show that *C* is an isometric cycle in FQ_n . It is sufficient to show that $d_C(u, v) = d_{FQ_n}(u, v)$ for any two distinct vertices *u* and *v* of *C*. By Theorem 5.7 (3), there are two cases for the shortest *uv*-path in FQ_n . If $H_{Q_n}(u, v) \leq \lfloor \frac{n}{2} \rfloor$, then any shortest *uv*-path in FQ_n contains no complementary edges. So $d_{FQ_n}(u, v) = d_{Q_n}(u, v) = H_{Q_n}(u, v) = d_C(u, v)$. If $H_{Q_n}(u, v) > \lceil \frac{n}{2} \rceil$, then any shortest *uv*-path in FQ_n contains exactly one complementary edge. Let P_1 be the *uv*-path on *C* that contains the unique complementary edge *e*. Since $H_{Q_n}(u, v) > \lceil \frac{n}{2} \rceil$ and the length of *C* is n + 1, $d_C(u, v) = |P_1| = n + 1 - H_{Q_n}(u, v) < \lceil \frac{n}{2} \rceil$. Clearly $d_{FQ_n}(u, v) \leq d_C(u, v)$, that is, P_1 is not a shortest *uv*-path in FQ_n . Let P_2 be a shortest *uv*-path in FQ_n . Then P_2 contains exactly one complementary edge by Theorem 5.7 (3). Set $P' := C - (V(P_1) \setminus \{u, v\})$. Then $P' \cup P_2$ is a walk in FQ_n that has exactly one complementary edge. So there is a cycle $C' \subseteq P' \cup P_2$ that contains exactly one complementary edge. We can deduce a contradiction by Theorem 5.7 (2) as follows:

$$n+1 \le |C'| \le |P'| + |P_2| < |P'| + |P_1| = |C| = n+1.$$

So $d_{FQ_n}(u, v) = d_C(u, v)$.

For any $i \in \{1, 2, ..., n\}$, let P^i be a shortest path from $x_1 x_2 \cdots x_n$ to $\bar{x}_1 \bar{x}_2 \cdots \bar{x}_n$ in Q_n such that the unique edge in $P^i \cap E_i$ is the antipodal edge of e on $C^i := P^i \cup \{e\}$. Since C^i is an isometric even cycle by the above proof, the unique complementary edge e on C^i and its antipodal edge $P^i \cap E_i$ are in relation Θ_{FQ_n} . So E_i and E_{n+1} are contained in an equivalence class with respect to $\Theta^*_{FQ_n}$, i = 1, 2, ..., n. Hence FQ_n is a prime graph under the Cartesian product by Lemma 5.6.

Now we know that for $m \ge 3$ and $n \ge 2$, $K_{m,m}$, K_{2n} and FQ_n are prime extremal graphs. From Proposition 4.3, it is interesting to characterize all the prime extremal graphs.

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