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Computing the number of $h$-edge spanning forests in complete bipartite graphs

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Let $f_{m,n,h}$ be the number of spanning forests with $h$ edges in the complete bipartite graph $K_{m,n}$. Kirchhoff’s Matrix Tree Theorem implies $f_{m,n,m} + n - 1 = m^n - 1 - n^m - 1$ when $m \geq 1$ and $n \geq 1$, since $f_{m,n,m} + n - 1$ is the number of spanning trees in $K_{m,n}$. In this paper, we give an algorithm for computing $f_{m,n,h}$ for general $m, n, h$. We implement this algorithm and use it to compute all non-zero $f_{m,n,h}$ when $m \leq 50$ and $n \leq 50$ in under 2 days.

MSC 2010: 68R10, 05C85, 05C30, 05C38, 05C05

Keywords: complete bipartite graph, spanning forest

1 Introduction

A consequence of Kirchhoff’s Matrix Tree Theorem is that the number of spanning trees in the complete bipartite graph $K_{m,n}$ is $m^n - 1 - n^m - 1$ when $m \geq 1$ and $n \geq 1$. In this paper, we consider a generalisation of this counting problem: the number $f_{m,n,h}$ of $h$-edge spanning forests in $K_{m,n}$. The number of spanning trees in $K_{m,n}$ is $f_{m,n,m}$ when $m \geq 1$ and $n \geq 1$, since spanning trees of $K_{m,n}$ have $m + n - 1$ edges.

The numbers $f_{m,n,h}$ arise in relation to the Tutte polynomial of $K_{m,n}$. For a general graph $G = (V, E)$, the Tutte polynomial is

$$T_G(x,y) = \sum_{A \subseteq E} (x-1)^{k(A)}(y-1)^{k(A)+|A|-|V|}$$

where $k(A)$ is the number of connected components in the graph $(V, A)$. Importantly, for the Tutte polynomial, we take $0^0 = 1$. Thus, the number of spanning forests of $G$ is given by

$$T_G(2,1) = \sum_{A \subseteq E} 0^{k(A)+|A|-|V|}$$

$$= \# \{ A \subseteq E : k(A) + |A| = |V| \}$$

$$= \text{number of spanning forests of } G$$

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since \( k(A) + |A| = |V| \) if and only if \( A \) is a spanning forest.

In the present paper, we are interested in the case of when \( G = K_{m,n} \), and will look for \( h \)-edge spanning forests. In this situation,

\[
T_{K_{m,n}}(2, 1) = \sum_{h \geq 0} f_{m,n,h}.
\]

For general graphs \( G = (V, E) \), Myrvold [10] gave an algorithm for computing the number of \( k \)-component spanning forests of \( G \). Björklund et al. [4] described an algorithm that could compute \( T_{G}(2, 1) \) in time \( 2^{|V|}|V|^{O(1)} \) by using Kirchhoff’s Matrix Tree Theorem for each of the \( 2^{|V|} \) subsets of \( V \), then combining the results together with Inclusion-Exclusion. Other relevant work for the enumeration of spanning forests includes [9]. Porter [11] gave an algorithm for generating the spanning trees of \( K_{m,n} \).

Bounds for the number of spanning forests in graphs were given by Teranishi [13], from which we can deduce

\[
T_{K_{m,n}}(2, 1) \geq \sum_{k \geq 0} \left( \min(m,n) \right)^{m+n-k} \binom{m+n}{k}
\]

for all \( m \geq 0 \) and \( n \geq 0 \). Jin and Liu [8] gave a simple formula for the number of rooted spanning forests in \( K_{m,n} \) (see also [7] and [12]).

This paper instead heads in a different direction to previous work: we derive algorithms for enumerating \( h \)-edge spanning forests in \( K_{m,n} \).

2 Basic results

To be clear, all forests in this paper will be labelled. Let \( \mathbb{N} = \{0, 1, 2, \ldots\} \). The following lemma gives the boundary conditions on \( f_{m,n,h} \).

**Lemma 1** Suppose \( m, n, h \in \mathbb{N} \). Then

- \( f_{m,n,0} = 1 \),
- if \( h \geq 1 \), then \( f_{0,n,h} = f_{m,0,h} = 0 \),
- if \( m \geq 1 \) and \( n \geq 1 \), then \( f_{m,n,h} > 0 \) if and only if \( 0 \leq h \leq m + n - 1 \),
- if \( m \geq 1 \) and \( n \geq 1 \), then \( f_{m,n,m+n-1} = m^{n-1}n^{m-1} \).

**Proof:** If \( h = 0 \), then there is exactly one subgraph of \( K_{m,n} \) with no edges (and it is a spanning forest), so \( f_{m,n,0} = 1 \). If \( m = 0 \) or \( n = 0 \), then \( K_{m,n} \) has no edges, and thus we obtain the second bulleted item.

Now assume \( m \geq 1 \) and \( n \geq 1 \). Since \( K_{m,n} \) is connected, it has a spanning tree, which must have exactly \( m + n - 1 \) edges. By deleting edges from this spanning tree, we find \( h \)-edge forests in \( K_{m,n} \) for all \( 0 \leq h \leq m + n - 1 \). Hence \( f_{m,n,h} > 0 \) when \( 0 \leq h \leq m + n - 1 \).

Now assume \( A \) is an \( h \)-subset of the edges in \( K_{m,n} \), where \( m \geq 1 \) and \( n \geq 1 \) and \( h \geq m + n \). To be a spanning forest, we need \( k(A) + h = m + n \) as per [1], but if \( h \geq m + n \), then \( k(A) \leq 0 \), giving a contradiction.

\(^{(i)}\) \#P is the set of counting problems which ask for the number of “yes” instances for decision problems in NP. A \#P problem is in \#P-complete whenever any other problem in \#P can be reduced to it by a polynomial-time counting reduction.
The number of $h$-edge spanning forests in $K_{m,n}$

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Fig. 1: Small non-zero values of $f_{m,n,h}$.

The fourth bulleted item is from Kirchhoff’s Matrix Tree Theorem (already mentioned).

The following lemma is an important (but basic) consequence of the symmetry of the problem.

**Lemma 2** For all $m, n, h \in \mathbb{N}$, we have $f_{m,n,h} = f_{n,m,h}$.

In Figure 1 we list the non-zero values of $f_{m,n,h}$ when $0 \leq m \leq n \leq 6$. These values were generated using a straightforward backtracking algorithm.

3 Combinatorial equivalence

In this section, we will describe some combinatorial objects that are equivalent to $h$-edge spanning forests of $K_{m,n}$.

Let $M$ be an $m \times n$ $(0,1)$-matrix. We define a cycle in $M$ to be a set $\{e_1, e_2, \ldots, e_t\}$ of entries of $M$, such that:

- each $e_i$ contains the symbol 1, and

- for $i \in \{1, 2, \ldots, t\}$, we have that (a) $e_{2i-1}$ belongs to the same row as $e_{2i}$ and (b) $e_{2i}$ belongs to the same column as $e_{2i+1}$ (where we take $e_{2t+1} = e_1$).

For example, the matrix

$\begin{bmatrix}
0 & 1 & 0 & 1 & 1 \\
0 & 0 & 1 & 0 & 1 \\
1 & 0 & 0 & 0 & 1 \\
1 & 1 & 0 & 0 & 0 \\
\end{bmatrix}$
Rebecca J. Stones

has a cycle (consisting of the circled \(1\)'s), whereas the matrices

\[
\begin{bmatrix}
0 & 1 & 1 \\
0 & 0 & 0
\end{bmatrix} \quad \text{and} \quad \begin{bmatrix}
0 & 1 & 1 & 0 & 1 \\
1 & 0 & 0 & 1 & 0 \\
1 & 0 & 0 & 0 & 0
\end{bmatrix}
\]
do not have cycles.

Let \(A_{m,n,h}\) be the set of \((0, 1)\)-matrices with exactly \(h\) elements equal to 1. Let \(p_{m,n,h}\) be the probability that an element of \(A_{m,n,h}\) chosen uniformly at random contains a cycle. An anonymous user of math.stackexchange.com asked for a formula for \(p_{m,n,h}\). In fact, this was the original motivation for the author to study this problem.

**Lemma 3** For all \(m, n, h \in \mathbb{N}\), we have \(f_{m,n,h}\) is the number of \(m \times n\) \((0, 1)\)-matrices with \(h\) 1’s that do not contain a cycle, and hence

\[
f_{m,n,h} = (1 - p_{m,n,h}) \binom{mn}{h}.
\]

**Proof:** A matrix \(M \in A_{m,n,h}\) can be interpreted as the biadjacency matrix of an \(h\)-edge subgraph \(G = G(M)\) of \(K_{m,n}\). A cycle in \(M \in A_{m,n,h}\) corresponds to a cycle in \(G\). Hence

\[
f_{m,n,h} = (1 - p_{m,n,h}) |A_{m,n,h}| = (1 - p_{m,n,h}) \binom{mn}{h}.
\]

\(\square\)

A problem related to the \((0, 1)\)-matrix problem above comes from the study of \((0, 1)\)-matrices that do not have a submatrix which is the incidence matrix of any cycle of length at least 3; these are called “totally balanced matrices” (see e.g. [2]).

Another interpretation of matrices in \(A_{m,n,h}\) is as induced subgraphs of \(K_m \square K_n\), where \(\square\) represents the Cartesian product of graphs. The graph \(K_m \square K_n\) is sometimes called the “rook’s graph”, since the edges represent the legal moves of a rook on an \(m \times n\) chess board. The vertices in \(K_m \square K_n\) are \(\{(i, j) : 1 \leq i \leq m \text{ and } 1 \leq j \leq n\}\) and there is an edge between distinct vertices \((i, j)\) and \((i', j')\) whenever \(i = i'\) or \(j = j'\). The graph \(K_m \square K_n\) is also the line graph of \(K_{m,n}\).

There is a bijection between induced subgraphs \(H = H(M)\) of \(K_m \square K_n\) and \((0, 1)\)-matrices \(M = (m_{ij}) \in A_{m,n,h}\): we include the vertex \((i, j)\) in \(H\) if and only if \(m_{ij} = 1\). For example, if we ignore the vertices marked 0 in

\[
\begin{array}{ccc}
0 & 1 & 1 \\
1 & 0 & 0 \\
1 & 0 & 0
\end{array}
\]

(ii) Full URL: [http://math.stackexchange.com/q/24800](http://math.stackexchange.com/q/24800)

(iii) The biadjacency matrix of a bipartite graph \(G\) on the vertex set \(\{u_i\}_{1 \leq i \leq m} \cup \{v_j\}_{1 \leq j \leq n}\) is the \(m \times n\) \((0, 1)\)-matrix \(M = (m_{ij})\) with \(m_{ij} = 1\) if and only if \(u_i v_j\) is an edge in \(G\).
The number of $h$-edge spanning forests in $K_{m,n}$

we obtain an induced subgraph $H(M)$ of $K_3 \Box K_5$, corresponding to the matrix

$$M = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \end{pmatrix}$$

in $A_{3,5,6}$. We see that cycles in $M$ map from cycles $C_k$ in $H$ for some $k \in \{4, 6, 8, \ldots\}$ (an example of a 6-cycle is highlighted). There are other cycles in $H$ (e.g. $C_3$ in the above example; in fact, if $M$ above had a row of 1’s, then $H$ would have a $K_5$ subgraph). However, induced cycles $C_k$ in $H(M)$ of length $k \in \{4, 6, 8, \ldots\}$ are in one-to-one correspondence with cycles in $M$ via the above bijection.

**Lemma 4** For all $m, n, h \in \mathbb{N}$, we have that $f_{m,n,h}$ is the number of $h$-vertex induced subgraphs of $K_{m,n}$ that do not contain an induced $C_k$ for any $k \in \{4, 6, 8, \ldots\}$.

### 4 Simplifying the equation

#### 4.1 A formula for $f_{m,n,h}$

We will look at two related ways of simplifying $f_{m,n,h}$. Let $K_{m,n}$ have the vertex bipartition $M \cup N$, where $M = \{u_1, u_2, \ldots, u_m\}$ and $N = \{v_1, v_2, \ldots, v_n\}$. Let $B_{m,n,h}$ be the set of $h$-edge spanning forests of $K_{m,n}$. Hence $f_{m,n,h} = |B_{m,n,h}|$. Define

$$W_{m,n,h} = \{G \in B_{m,n,h} : G \text{ has no isolated vertices}\}$$

and $w_{m,n,h} = |W_{m,n,h}|$. We observe the following:

- Given any spanning forest in $B_{m,n,h}$ we can delete the isolated vertices to obtain a spanning forest in $W_{i,j,h}$ for some $i \leq m$ and $j \leq n$.
- Conversely, given any spanning forest in $W_{i,j,h}$, we can add isolated vertices to it in $\binom{m}{i-j} \binom{n}{j}$ ways to obtain a spanning forest in $B_{m,n,h}$.

In both of the operations above, we need to relabel the vertices, but we preserve the order of the indices of the non-isolated vertices within $M$ and $N$.

Hence

$$f_{m,n,h} = \sum_{i=0}^{m} \sum_{j=0}^{n} \binom{m}{i} \binom{n}{j} w_{i,j,h}. \quad (2)$$

This formula eliminates the need to further account for isolated vertices.

We can decompose any $G \in W_{m,n,h}$ into disjoint components, each of which belong to $W_{i,j,k}$ for some $i \leq m$ and $j \leq n$ and $k \leq h$. Hence any $G \in W_{m,n,h}$ decomposes into the following.

- A partition $P$ of $M$ where $x, y \in M$ belong to the same part if there is a path from $x$ to $y$ in $G$.
- A partition $Q$ of $N$ where $x', y' \in N$ belong to the same part if there is a path from $x'$ to $y'$ in $G$.
- A bijection $\alpha : P \rightarrow Q$ such that every edge $ab$ in $G$ has $a \in p$ and $b \in \alpha(p)$ for some $p \in P$. Hence we have $|P| = |Q|$.
For each \( p \in P \), a subgraph in \( W_{|p|,|\alpha(p)|,|p|+|\alpha(p)|-1} \) induced by the vertices in \( p \cup \alpha(p) \).

Regarding the final bulleted item, we observe that the subgraph induced by \( p \cup \alpha(p) \) for some \( p \in P \) must be a spanning tree on \( K_{|p|,|\alpha(p)|} \). Any more edges would cause a cycle, while fewer edges would cause disjoint components (in which case \( G \) could be further decomposed).

We give an example of a graph \( G \) in \( W_{3,5,6} \) below.

Here, the partitions \( P \) and \( Q \) are given by \( P = \{\{u_1, u_2\}, \{u_3\}\} \) and \( Q = \{\{v_1, v_2\}, \{v_3, v_4, v_5\}\} \) and we have the bijection \( \alpha \) such that

\[
\{u_1, u_2\} \mapsto \{v_1, v_2\}, \\
\{u_3\} \mapsto \{v_3, v_4, v_5\}.
\]

We see that \( G \) decomposes into a graph in \( W_{2,2,3} \) and a graph in \( W_{1,3,3} \).

In general, for any \( p \in P \), we must have exactly \( |p| + |\alpha(p)| - 1 \) edges in \( p \cup \alpha(p) \). Summing this over all \( p \in P \) gives \( h = m + n - |P| \). We conclude that

\[
|P| = m + n - h
\]

(and thus \( |Q| = m + n - h \)).

We will now describe how to construct any graph in \( W_{m,n,h} \) via its decomposition. Let \( \mathcal{P} \) be the set of partitions of \( M \) and let \( \mathcal{Q} \) be the set of partitions of \( N \). Given (a) a partition \( P \in \mathcal{P} \) of size \( |P| = m + n - h \), (b) a partition \( Q \in \mathcal{Q} \) and (c) a bijection \( \alpha : P \rightarrow Q \), we can construct \( \prod_{p \in P} w_{|p|,|\alpha(p)|,|p|+|\alpha(p)|-1} \) distinct graphs \( G \in W_{m,n,h} \) by the following steps:

1. Start with \( G \) as the graph with vertex set \( M \cup N \) and no edges.

2. For each \( p \in P \) add one of the subgraphs in \( W_{|p|,|\alpha(p)|,|p|+|\alpha(p)|-1} \) on the vertices \( p \cup \alpha(p) \).

Since all graphs in \( W_{m,n,h} \) can be constructed uniquely by the above steps we have the following theorem.

**Theorem 1** For \( m, n, h \in \mathbb{N} \),

\[
f_{m,n,h} = \sum_{i=0}^{m} \sum_{j=0}^{n} \binom{m}{i} \binom{n}{j} w_{i,j,h}
\]
The number of $h$-edge spanning forests in $K_{m,n}$

where

$$w_{m,n,h} = \sum_{P \in P} \sum_{Q \in Q} \sum_{\alpha: P \to Q} \prod_{p \in P} w_{|p|, |\alpha(p)|, |p|+|\alpha(p)|-1}$$

where

$$w_{|p|, |\alpha(p)|, |p|+|\alpha(p)|-1} = |p|^{\alpha(p)}|p|-1.$$  

4.2 An improved formula for $f_{m,n,h}$

We will derive two further equations for $f_{m,n,h}$ (in Theorems 2 and 3), which are related to Theorem 1. They result in a slightly more complicated algorithm, but will allow us to compute $f_{m,n,h}$ faster than (2).

For all $m$, $n$, $h \in \mathbb{N}$, define

$$W'_{m,n,h} = \{G \in W_{m,n,h} : \text{the vertices in } N \text{ have degree } \geq 2\}.$$  

Given any $G \in B_{m,n,h}$ we can delete the isolated vertices and the leaves (vertices of degree 1) in $N$ to obtain a spanning forest in $W'_{i,j,h-k}$ for some $i \leq m$ and $j \leq n$ and $k \in \mathbb{N}$. Conversely, given any $G \in W'_{i,j,h-k}$, we can add:

- $m - i$ isolated vertices to $M$ in \binom{m}{m-i}$ ways, so as to increase $|M|$ to $m$,
- $n - j$ isolated vertices to $N$ in \binom{n}{n-j}$ ways, so as to increase $|N|$ to $n$,
- $k$ edges to $G$ in \binom{n-j}{k}$ ways, so as to increase the number of leaves in $N$ by $k$,

thereby obtaining a spanning forest in $B_{m,n,h}$. Similar to the derivation of (2), we need to relabel the vertices so as to preserve the order of the indices of the non-isolated vertices within $M$ and the non-isolated non-leaf vertices within $N$. Hence

$$f_{m,n,h} = \sum_{i=0}^{m} \sum_{j=0}^{n} \sum_{k=0}^{\min(n-j,h)} \binom{m}{i} \binom{n}{j,k,n-j-k} m^k w'_{i,j,h-k}.$$  

All of the steps involved for finding the formula (3) for $w_{m,n,h}$ are still valid for $w'_{m,n,h}$. Hence (3) remains true if we replace $w$ with $w'$. However, we can no longer make use of the simple formula for $w_{|p|, |\alpha(p)|, |p|+|\alpha(p)|-1}$. Nevertheless, we will be able to find a formula for $w'_{|p|, |\alpha(p)|, |p|+|\alpha(p)|-1}$ via Inclusion-Exclusion.

For $I \subseteq N$, define

$$A_I = \{G \in W_{m,n,m+n-1} : \text{vertices in } I \text{ are leaves}\}.$$  

Note that spanning forests in $A_I$ might also have leaves outside of $I$. Hence $|A_0| = |W_{m,n,m+n-1}| = m^{n-1}n^{m-1}$. By symmetry, if $I, J \subseteq N$ and $|I| = |J|$ then $|A_I| = |A_J|$. So we will assume $I = \{n - i + 1, n - i + 2, \ldots, n\}$. We can construct the graphs in $A_I$ by adding $i$ leaves to the graphs in $W_{m,n-|I|,m+n-|I|-1}$ (i.e., the set of spanning trees of $K_{m,n-|I|}$), and these leaves can be added in $m^i$ ways. Hence

$$|A_I| = m^i \cdot \text{number of spanning trees of } K_{m,n-i}$$

$$= m^i m^{n-i}(n-i)^{m-1}$$

$$= m^{n-1}(n-i)^{m-1}.$$
when \( m \geq 1 \) and \( n \geq 1 \). Here we take \( 0^0 = 1 \), since we want to account for \( K_{1,0} \in W_{1,0,0} \). Hence, by Inclusion-Exclusion, we find

\[
\begin{align*}
\wpr_{m,n,m+n-1} &= |A_0| - \left| \bigcup_{I \subseteq N} A_I \right| \\
&= m^{n-1}n^{m-1} - \sum_{i=1}^{n} (-1)^{i-1} \binom{n}{i} m^{n-1}(n-i)^{m-1} \\
&= m^{n-1} \sum_{i=0}^{n} (-1)^{i} \binom{n}{i} (n-i)^{m-1}
\end{align*}
\]

when \( m \geq 1 \) and \( n \geq 1 \). By definition, we also have \( \wpr_{m,n,m+n-1} = 0 \) when either \( m = 0 \) or \( n = 0 \) (and \( \wpr_{m,n,m+n-1} \) is not defined when \( m = 0 \) and \( n = 0 \)).

Before identifying formulae for \( \wp_{m,n,h} \) and \( \wpr_{m,n,h} \) in the next section, we make the following observation.

**Lemma 5** For all \( m, n, h \in \mathbb{N} \), we have \( n! \) divides \( \wpr_{m,n,h} \).

**Proof:** The symmetric group on \( N \) acts on \( W'_{m,n,h} \) by permuting the vertices in \( N \). Suppose a graph \( G \in W'_{m,n,h} \) has a non-trivial stabiliser subgroup under this action. Then there exists a non-identity permutation \( \alpha \) such that \( \alpha G = G \), and distinct vertices \( v, v' \in N \) for which \( \alpha(v) = v' \). Since \( G \in W'_{m,n,h} \), the degree of \( v \) is 2 or more, so assume \( v \) has distinct neighbours \( a, b \in M \). Then, since \( \alpha G = G \), we find that \( v' = \alpha(v) \) has the neighbours \( a \) and \( b \) too. Thus, \( \{v, v', a, b\} \) induces a 4-cycle, giving a contradiction. Hence, all stabiliser subgroups are trivial, and by the Orbit-Stabiliser Theorem, all orbits have size \( n! \). Hence \( n! \) divides \( \wpr_{m,n,h} \). \( \square \)

### 4.3 Formulae for \( \wp_{m,n,h} \) and \( \wpr_{m,n,h} \)

Now we will simplify (3), noting that our simplifications remain valid when we replace \( w \) with \( w' \).

For \( a \geq 1 \) and \( t \geq 1 \), let \( S_{a,t} \) be the set of (number) partitions of \( a \) into \( t \) non-zero parts. We will interpret elements of \( S_{a,t} \) as multisets; for example \( \{4, 3, 3, 1, 1\} \in S_{12,5} \). We will require \( S_{a,0} = \emptyset \) when \( t < 0 \), and \( S_{0,0} = \{\emptyset\} \) and \( S_{0,t} = \emptyset \) when \( t > 0 \). For \( z \in S_{a,t} \), let \( T[z] \) be the set of partitions of \( \{1, 2, \ldots, a\} \) whose part sizes induce the number partition \( z \). For \( z \in S_{a,t} \), let \( \hat{z} \) denote an arbitrary element of \( T[z] \). If \( P \in T[x] \) and \( Q \in T[z] \), then any bijection \( \alpha : P \to Q \) induces a bijection \( \beta : x \to z \).
The number of \( h \)-edge spanning forests in \( K_{m,n} \)

Hence

\[
\begin{align*}
w_{m,n,h} &= \sum_{x,z \in S_{m,n,m+n-h}} \sum_{P \in \mathcal{T}[x]} \sum_{Q \in \mathcal{T}[z]} \alpha: P \to Q \sum_{\alpha: P \to Q} \prod_{p \in P} w_{\alpha[p],|\alpha(p)|,|p|+|\alpha(p)|-1} \\
&= \sum_{x,z \in S_{m,n,m+n-h}} \sum_{P \in \mathcal{T}[x]} \sum_{Q \in \mathcal{T}[z]} \alpha: z \to \hat{z} \sum_{\alpha: z \to \hat{z}} \prod_{p \in \hat{z}} w_{\alpha[p],|\alpha(p)|,|p|+|\alpha(p)|-1} \\
&= \sum_{x,z \in S_{m,n,m+n-h}} |T[x]| \cdot |T[z]| \sum_{\beta: x \to \hat{x}} \prod_{\beta: x \to \hat{x}} w_{r,\beta(r),r+\beta(r)-1}.
\end{align*}
\]

If \( z \in S_{a,t} \), then

\[
|T[z]| = \frac{a!}{\prod_{i \geq 1} i^{s_i(z)} s_i(z)!}
\]

where \( s_i(z) \) denotes the number of parts \( i \) in \( z \) (see e.g. [1] Theorem 13.2]). Hence we have the following theorems.

**Theorem 2** For \( m, n, h \in \mathbb{N} \),

\[
f_{m,n,h} = \sum_{i=0}^{m} \sum_{j=0}^{n} \left( \begin{array}{c} m \\ i \end{array} \right) \left( \begin{array}{c} n \\ j \end{array} \right) w_{i,j,h}
\]

where

\[
w_{m,n,h} = \sum_{x,z \in S_{m,n,m+n-h}} \sum_{\beta: x \to \hat{x}} \frac{m!}{\prod_{i \geq 1} i^{s_i(x)} s_i(x)!} \frac{n!}{\prod_{i \geq 1} i^{s_i(z)} s_i(z)!} \prod_{r \in x} r^{\beta(r)-1} \beta(r)^{r-1}. \tag{4}
\]

**Theorem 3** For \( m, n, h \in \mathbb{N} \),

\[
f_{m,n,h} = \sum_{i=0}^{m} \sum_{j=0}^{n} \sum_{k=0}^{\min(n-j,h)} \left( \begin{array}{c} m \\ i \end{array} \right) \left( \begin{array}{c} n \\ j \end{array} \right) m^k w_{i,j,h-k}.
\]

where

\[
w'_{m,n,h} = \sum_{x,z \in S_{m,n,m+n-h}} \sum_{\beta: x \to \hat{x}} \frac{m!}{\prod_{i \geq 1} i^{s_i(x)} s_i(x)!} \frac{n!}{\prod_{i \geq 1} i^{s_i(z)} s_i(z)!} \prod_{r \in x} w_{r,\beta(r),r+\beta(r)-1}
\]

where

\[
w'_{m,n,m+n-1} = m^{n-1} \sum_{i=0}^{n} (-1)^i \binom{n}{i} (n-i)^{m-1}.
\]

The main advantage of Theorems 2 and 3 over Theorem 1 is the summation is over number partitions (rather than set partitions).
5 Implementation

The author has implemented the two formulae for \( \tilde{f}_{m,n,h} \) described in the preceding section, along with a simple backtracking algorithm. The results between all three implementations concur, giving confidence in the accuracy of the code. The C source code is available as supplementary material to this document. The GMP library was used for arbitrary precision arithmetic \([6]\).

The formula involving \( w'_{m,n,h} \) instead of \( w_{m,n,h} \) was (unsurprisingly) much faster, largely because many values of \( w_{m,n,h} \) equal 0, such as when \( m \leq n \), which can be used to drastically reduce the search tree.

5.1 Symmetry breaking

The run-time of the program was also improved through the use of symmetry breaking, which we will now describe. For \( x \in S_{m,m+n-h} \) and \( z \in S_{m,m+n-h} \), there are often many bijections \( \beta : x \rightarrow z \) which map the same elements to the same elements (since both \( x \) and \( z \) are multisets). In these instances, a na"ive algorithm would repeat the same computation unnecessarily. Let \( x \) and \( z \) be the multisets \( x = \{x_1, x_2, \ldots, x_{m+n-h}\} \) and \( z = \{z_1, z_2, \ldots, z_{m+n-h}\} \). We define the condition:

**Symmetry breaking condition 1:** We say \( \beta : x \rightarrow z \) is half-canonical if \( \beta^{-1}(z_{i-1}) < \beta^{-1}(z_i) \) whenever \( z_{i-1} = z_i \).

We need to add a multiplicative factor to adjust for the restriction to half-canonical bijections. Hence

\[
\begin{align*}
 w_{m,n,h} &= \sum_{x,z \in S_{m,m+n-h}} \sum_{\beta : x \rightarrow z, \beta \text{ is a bijection}} \left( \prod_{i \geq 1} l^{s_i(z)} \right) \frac{m!}{\prod_{i \geq 1} l^{s_i(x)} s_i(x)!} \frac{n!}{\prod_{i \geq 1} l^{s_i(z)} s_i(z)!} \prod_{r \in x} w_{r, \beta(r), r+\beta(r)-1} \\
\end{align*}
\]

and similarly with \( w' \) in place of \( w \). Using this assumption, we reduce the number of bijections by a factor of \( \prod_{i \geq 1} l^{s_i(z)} \), which results in a substantial time saving.

Instead of the single symmetry breaking assumption, it is possible to utilise symmetry breaking using an additional condition:

**Symmetry breaking condition 2:** We say \( \beta : x \rightarrow z \) is canonical if it is half-canonical and \( \beta(x_{i-1}) < \beta(x_i) \) whenever \( x_i = x_{i-1} \).

Again, we find

\[
\begin{align*}
 w_{m,n,h} &= \sum_{x,z \in S_{m,m+n-h}} \sum_{\beta : x \rightarrow z, \beta \text{ is a bijection}} \Gamma_{x,z,\beta} \frac{m!}{\prod_{i \geq 1} l^{s_i(x)} s_i(x)!} \frac{n!}{\prod_{i \geq 1} l^{s_i(z)} s_i(z)!} \prod_{r \in x} w_{r, \beta(r), r+\beta(r)-1} \\
\end{align*}
\]

and similarly with \( w' \) in place of \( w \), where \( \Gamma_{x,z,\beta} \) is the number of bijections between the multisets \( x \) and \( z \) which map the same elements to the same elements (as \( \beta \)). A formula for \( \Gamma_{x,z,\beta} \) was given by e.g. \([3]\), namely

\[
\Gamma_{x,z,\beta} = \frac{\prod_{i \geq 1} s_i(x)! \prod_{i \geq 1} s_i(z)!}{\prod_{i \geq 1} \prod_{j \geq 1} s_{i,j}(x, \beta)!}
\]

where
The number of $h$-edge spanning forests in $K_{m,n}$

- $s_i(z)$ denotes the number of parts $i$ in $z$ (as before), and
- $s_{i,j}(x,\beta)$ is the number of elements $(i, j)$ in the multiset $\{(r, \beta(r)) : r \in x\}$.

The author has implemented both of these symmetry breaking schemes in order to compare their performance (see Section 5.3).

5.2 Pseudo-code

Algorithm 1 gives a pseudo-code version of the C code used to implement the algorithm described by Theorem 3 using the half-canonical symmetry breaking condition. The partitions of $m$ and $n$ into $k$ parts were computed whenever needed and stored in memory. Iterating through the half-canonical bijections was performed “on the fly” using a backtracking algorithm.

While Theorems 2 and 3 are valid for all $m, n, h \in \mathbb{N}$, we need to set $f_{m,n,0} = 1$ separately in the C code.

5.3 Performance

The C code was run on a $2 \times 2.66$ GHz processor (although the code itself is not parallelised). The following table gives the run-times (in seconds) for the two algorithms under the two symmetry breaking schemes when computing all non-zero $f_{m,n,h}$ with $m, n \leq 19$.

<table>
<thead>
<tr>
<th>$h$</th>
<th>half-canonical</th>
<th>canonical</th>
</tr>
</thead>
<tbody>
<tr>
<td>Theorem 2</td>
<td>27.1</td>
<td>6.8</td>
</tr>
<tr>
<td>Theorem 3</td>
<td>4.1</td>
<td>4.5</td>
</tr>
</tbody>
</table>

The fastest version is Theorem 3 using half-canonical symmetry breaking. Under these conditions, the code had the following run-times to find all non-zero values of $f_{m,n,h}$ with $m, n \leq t$:

<table>
<thead>
<tr>
<th>$t$</th>
<th>time (sec)</th>
</tr>
</thead>
<tbody>
<tr>
<td>25</td>
<td>20.5</td>
</tr>
<tr>
<td>26</td>
<td>27.7</td>
</tr>
<tr>
<td>27</td>
<td>35.4</td>
</tr>
<tr>
<td>28</td>
<td>46.0</td>
</tr>
<tr>
<td>29</td>
<td>62.2</td>
</tr>
<tr>
<td>30</td>
<td>83.0</td>
</tr>
</tbody>
</table>

This table indicates the scalability of the program, which is not overwhelming for these values of $t$. The author also ran the program to compute $f_{m,n,h}$ with $m, n \leq t$ where $t = 50$, which took under 2 days (to be precise, it took 1 day 15 hours and 17 minutes). The largest number encountered was $f_{50,50,101}$, which has 167 digits, and is equal to the number of spanning trees of $K_{50,50}$, which is $50^{50-1} \cdot 50^{50-1} = 50^{98}$.

5.4 Complexity

In the worst case, $x = z = \{1, 2, 3, \ldots, t\}$, where $\min(m,n) = 1 + 2 + \cdots + t = \frac{1}{2}t(t + 1)$ (so $t = O(\sqrt{m})$), in which case the program must iterate through all $t!$ bijections from $x$ to $z$. So, for a worst case analysis, we assume $m = n$ and $m + n - h = \lceil \sqrt{m} \rceil$. 

Algorithm 1 Implementation of Theorem 3 using half-canonical symmetry breaking

\begin{align*}
&f_{m,n,m+n-1} := 0; \text{ for all } m, n, h \\
&w'_{m,n,m+n-1} := 0; \text{ for all } m, n, h \\
&w'_{0,0,0} := 1; \\
&\text{for } m = 0, 1, \ldots, \text{MAX}(m) \text{ do} \\
&\hspace{1em} \text{for } n = 0, 1, \ldots, \text{MAX}(n) \text{ do} \\
&\hspace{2em} \text{if } m > n \text{ and } m > 0 \text{ and } n > 0 \text{ then} // \text{ Other values of } w'_{m,n,h} \text{ are 0} \\
&\hspace{3em} \text{for } i = 0, 1, \ldots, n \text{ do} \\
&\hspace{4em} // \text{Using the Inclusion-Exclusion formula} \\
&\hspace{5em} w'_{m,n,m+n-1} := w'_{m,n,m+n-1} + (-1)^i \binom{n}{i} (n-i)^{m-1}; \\
&\hspace{2em} \text{end for} \\
&\hspace{1em} \text{end for} \\
&\text{end if} \\
&\text{f}_{m,n,0} = 1; \\
&\text{for } h = 1, 2, \ldots, m + n - 1 \text{ do} \\
&\text{for } i = 0, 1, \ldots, m \text{ do} \\
&\text{for } j = 0, 1, \ldots, n \text{ do} \\
&\text{for } k = 0, 1, \ldots, \min(n-j,h) \text{ do} \\
&\hspace{1em} f_{m,n,h} := f_{m,n,h} + \binom{m}{i} \binom{n}{j,k,n-j-k} m^k w'_{i,j,h-k}; \\
&\hspace{1em} \text{end for} \\
&\hspace{1em} \text{end for} \\
&\hspace{1em} \text{end for} \\
&\text{end for} \\
&\text{for } h = 0, 1, \ldots, m + n - 1 \text{ do} \\
&\text{Print } f_{m,n,h}; \\
&\text{end for} \\
&\text{end for} \\
&\text{end for} \\
&\text{end for}
\end{align*}
The number of \( h \)-edge spanning forests in \( K_{m,n} \)

Hence, given \( x, z \in S_{m,m+n-h} \), there can be \( O(\sqrt{m}!) \) canonical bijections from \( x \) to \( z \). Hardy and Ramanujan’s asymptotic formula for the number of partitions of \( m \), namely

\[
\frac{1}{4m\sqrt{3}} e^{\pi \sqrt{2m/3}},
\]
gives a crude asymptotic upper bound on \( |S_{m,\lfloor\sqrt{m}\rfloor}|^2 \), specifically

\[
|S_{m,\lfloor\sqrt{m}\rfloor}|^2 = O\left( \frac{\text{const} \cdot m}{m^2} \right)
\]
as the number of ways of choosing \( x \) and \( z \). Hence, \( (4) \) has

\[
O\left( \text{const} \cdot \sqrt{m} \cdot \sqrt{m/2 + \text{const}} \right)
\]
terms, by Stirling’s Approximation. In contrast, a backtracking algorithm would need to generate and check around \( (\frac{m^2}{\lfloor\sqrt{m}\rfloor}) \leq \frac{1}{\lfloor\sqrt{m}\rfloor} m^2 \sqrt{m} \) graphs, which, when \( m = n \) and \( h = 2m - \sqrt{m} \), is

\[
O(\sqrt{m} \cdot m^{1.5 \sqrt{m + \text{const}}})
\]
itations, by Stirling’s Approximation. Of course, when implementing these algorithms, we use pruning whenever possible to reduce the search space, which makes a drastic difference not accounted for in these approximations.

5.4.1 When \( h \) is fixed

We conclude this paper with the observation that, when \( h \) is fixed, computing \( f_{m,n,h} \) is asymptotically “easy”. The underlying reason is that, for sufficiently large \( m \) or \( n \), we must have isolated vertices in graphs in \( B_{m,n,h} \). Thus, \( (2) \) only contains a finite number of non-zero terms.

**Theorem 4** For fixed \( h \), computing \( f_{m,n,h} \) can be performed in time \( O(\log(mn)) \).

**Proof:** If \( i > h \) or \( j > h \), then \( w_{i,j,h} = 0 \) (since any graph in \( B_{i,j,h} \) must have an isolated vertex). Hence, \( (2) \) is equivalent to

\[
f_{m,n,h} = \sum_{i=0}^{h} \sum_{j=0}^{h} \binom{m}{i} \binom{n}{j} w_{i,j,h}.
\]

For fixed \( h \), there is a finite number of terms in this sum. Thus, for fixed \( h \), we could write a program in which:

- we store a list of the pairs \( (i, j) \) for which \( w_{i,j,h} \) is non-zero, along with the value of \( w_{i,j,h} \),
- we iterate through this list, computing \( \binom{m}{i} \binom{n}{j} \) \( w_{i,j,h} \), and add it to a running total.

We can compute \( \binom{m}{i} = \frac{1}{i} m(m-1) \cdots (m-i+1) \) using \( O(h) \) multiplications (since \( i \leq h \)), each of which takes time \( O(\log m) \), and one division. Hence, \( \binom{m}{i} \) can be computed in time \( O(\log(m)) \) time (since \( h \) is fixed). Similarly \( \binom{n}{j} \) can be computed in \( O(\log(n)) \) time. We conclude that the whole summation can be performed in time \( O(\log(mn)) \). \( \square \)
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References


