

Maker–Breaker domination number for Cartesian products of path graphs P_2 and P_n

Jovana Forcan^{1,2}

Jiayue Qi^{3*}

¹ Department of Mathematics and Informatics, Faculty of Sciences, University of Novi Sad, Serbia

² Department of Computer Sciences and Systems, Faculty of Philosophy, University of East Sarajevo, Bosnia and Herzegovina

³ Research Institute for Symbolic Computation, Johannes Kepler University, Linz, Austria

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We study the Maker–Breaker domination game played by Dominator and Staller on the vertex set of a given graph. Dominator wins when the vertices he has claimed form a dominating set of the graph. Staller wins if she makes it impossible for Dominator to win, or equivalently, she is able to claim some vertex and all its neighbours. Maker–Breaker domination number $\gamma_{MB}(G)$ ($\gamma'_{MB}(G)$) of a graph G is defined to be the minimum number of moves for Dominator to guarantee his winning when he plays first (second). We investigate these two invariants for the Cartesian product of any two graphs. We obtain upper bounds for the Maker–Breaker domination number of the Cartesian product of two arbitrary graphs. Also, we give upper bounds for the Maker–Breaker domination number of the Cartesian product of the complete graph with two vertices and an arbitrary graph. Most importantly, we prove that $\gamma'_{MB}(P_2 \square P_n) = n$ for $n \geq 1$, $\gamma_{MB}(P_2 \square P_n)$ equals $n, n - 1, n - 2$, for $1 \leq n \leq 4, 5 \leq n \leq 12$, and $n \geq 13$, respectively. For the disjoint union of $P_2 \square P_n$ s, we show that $\gamma'_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n$ ($n \geq 1$), and that $\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i)$ equals $k \cdot n, k \cdot n - 1, k \cdot n - 2$ for $1 \leq n \leq 4, 5 \leq n \leq 12$, and $n \geq 13$, respectively.

Keywords: Positional game, Maker–Breaker domination game, Maker–Breaker domination number for grids, winning strategy on grids

1 Introduction

1.1 Background

In this paper, we study the Maker–Breaker domination game which was first introduced in the literature by Duchêne, Gledel, Parreau and Renault in [Duchene et al., 2020]. This game combines two following research directions. In the original domination game introduced by Brešar, Klavžar, and Rall in [Brešar et al., 2010], two players, *Dominator* and *Staller*, alternately take a turn in claiming a vertex of the finite graph G which were not yet chosen in the course of the game — note that the chosen vertex of either Dominator or Staller must enlarge the set of dominated vertices. Dominator has the goal to dominate the graph in as few moves as possible while Staller tries to prolong the game as much as possible. This

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indicates that both players play optimally in such a game. This is why later on, in our paper, we sometimes use subjunctive mood in the sentences, when we analyze the cases that are evidently not optimal for the player. However, in general, we keep the grammar in present tense, since many times it is not clear yet whether the focused moves are optimal for the player or not.

The Maker–Breaker games, introduced by Erdős and Selfridge in [Erdős and Selfridge, 1973], are played on a finite hypergraph (X, \mathcal{F}) with the vertex set X and a set $\mathcal{F} \subseteq 2^X$ of hyper-edges. The set X is called the *board* of the game, while the set \mathcal{F} is called the *family of winning sets*. Two players, *Maker* and *Breaker*, take turns in claiming previously unclaimed elements of X . Maker wins the game if, by the end of the game, he has claimed all elements of some $F \in \mathcal{F}$. Otherwise, Breaker wins. For a deeper and more comprehensive analysis of Maker–Breaker games, see the book of Beck [Beck, 2008], and the monograph of Hefetz, Krivelevich, Stojaković and Szabó [Hefetz et al., 2014].

The Maker–Breaker domination game (MBD for short) is played on a graph $G = (V, E)$ by two players, *Dominator* and *Staller*. The aim of Dominator is to build a dominating set of the graph, which is a set T such that every vertex not in T has a neighbour in T . The aim of Staller is to claim a vertex from the graph G and all its neighbours, so that Dominator cannot dominate this vertex, which leads to Dominator failing in dominating all vertices of G . As concluded in [Duchene et al., 2020], this game is equivalent to a Maker–Breaker game played on the vertex set of a given graph as the board, with winning sets being all closed neighbourhoods of vertices.

When it is not hard to determine the identity of the winner in some Maker–Breaker game, then a more interesting question to ask is how fast the player with the winning strategy can win. Fast winning strategies for Maker in the Maker–Breaker games have received a lot of attention in the recent years (see for example [Clemens et al., 2012, 2015; Hefetz et al., 2009]). Specifically, for the Maker–Breaker domination game the smallest number of moves for Dominator is studied in [Gledel et al., 2019], where Gledel, Iršič, and Klavžar introduced the Maker–Breaker domination number $\gamma_{MB}(G)$ of a graph G , as the minimum number of moves for Dominator to win in the game on G where he is the first player. If Dominator is the second player, then the corresponding invariant is denoted by $\gamma'_{MB}(G)$ in their paper.

1.2 Preliminaries

Assume that the MBD game is in progress. As in [Gledel et al., 2019], we say that a game is the *D-game* if Dominator is the first to play, i.e. one *round* consists of a move by Dominator followed by a move of Staller. In the *S-game*, one round consists of a move by Staller followed by a move of Dominator. We denote by D_1, D_2, \dots (or D'_1, D'_2, \dots) the sequence of vertices chosen by Dominator and by S_1, S_2, \dots (or S'_1, S'_2, \dots) the sequence of vertices chosen by Staller, in a D-game (or in an S-game). We say that the vertex v is *isolated* by Staller if v and all its neighbours are claimed by Staller. For a given graph G , by $V(G)$ and $E(G)$ we denote its vertex set and edge set, respectively. If G is a graph and $S \subset V(G)$, then let $G|S$ denote the graph G in which the vertices from S are declared to be already dominated. As a preparation, we still need to introduce the *Continuation Principle* and the *No-Skip Lemma*.

Remark 1.1. (*The Continuation Principle, [Gledel et al., 2019, Remark 2.4]*) Let G be a graph with $A, B \subset V(G)$. If $B \subset A$, then $\gamma_{MB}(G|A) \leq \gamma_{MB}(G|B)$ and $\gamma'_{MB}(G|A) \leq \gamma'_{MB}(G|B)$.

Lemma 1.2. (*No-Skip Lemma, [Gledel et al., 2019, Lemma 2.3]*) In an optimal strategy of Dominator to achieve $\gamma_{MB}(G)$ or $\gamma'_{MB}(G)$, it is never an advantage for him to skip a move. Moreover, if Staller skips a move it can never disadvantage Dominator.

1.3 Main results

In [Gledel et al., 2019], the authors proposed finding the minimum number of moves for Dominator in the MBD game on the Cartesian product of two graphs. Motivated by the given problem, we give upper bounds for Maker–Breaker domination number for the Cartesian product of K_2 and an arbitrary graph, and that of two general graphs as well, in Section 2. The corresponding results is Theorem 1.3. Most importantly, we focus on determining how fast can Dominator win on the graphs $P_2 \square P_n$, for $n \geq 1$. From the results (Theorem 1.5 and Theorem 1.4) on the grids, we get some results on the disjoint union of $P_2 \square P_n$ s as stated in Theorem 1.6 and Theorem 1.7. We denote by $\dot{\cup}$ the disjoint union; and by $\dot{\cup}_{i=1}^k (G)_i$ we denote the disjoint union of k copies of graph G .

The paper is organized as follows. In Section 2, we prove Theorem 1.3. We prove Theorem 1.5 and Theorem 1.4 in Section 3. The corresponding proofs of Theorem 1.6 and Theorem 1.7 are given in Section 3.5. We list the main results of this paper as the following.

Theorem 1.3. *Let G and H be two arbitrary graphs on n and m vertices, respectively. Suppose that Dominator has winning strategies on G and H , for both D -games and S -games. Then*

$$\gamma_{MB}(G \square H) \leq \min\{\gamma_{MB}(G) + (m - 1) \cdot \gamma'_{MB}(G), \gamma_{MB}(H) + (n - 1) \cdot \gamma'_{MB}(H)\}.$$

Suppose that Dominator has a winning strategy for the S -games on both G and H , then

$$\gamma'_{MB}(G \square H) \leq \min\{m \cdot \gamma'_{MB}(G), n \cdot \gamma'_{MB}(H)\}.$$

Theorem 1.4. *For every positive integer n , it holds that $\gamma'_{MB}(P_2 \square P_n) = n$. If Dominator skipped any moves during this game, he would lose the game.*

Theorem 1.5. *Consider the D -games on $P_2 \square P_n$. Then*

1. *If $1 \leq n \leq 4$, then $\gamma_{MB}(P_2 \square P_n) = n$.*
2. *If $5 \leq n \leq 12$, then $\gamma_{MB}(P_2 \square P_n) = n - 1$.*
3. *If $n \geq 13$, then $\gamma_{MB}(P_2 \square P_n) = n - 2$.*

Theorem 1.6. *For every two positive integers n and k , it holds that*

$$\gamma'_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n.$$

Theorem 1.7. *Let k be any positive integer. Then*

1. *If $1 \leq n \leq 4$, then $\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n$.*
2. *If $5 \leq n \leq 12$, then $\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n - 1$.*
3. *If $n \geq 13$, then $\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n - 2$.*

2 MBD games on the Cartesian product of two graphs

In this section, we prove the result on the Cartesian products of general graphs.

Proof of Theorem 1.3: Consider, first, the D-game on $G \square H$. By $G^{(1)}, G^{(2)}, \dots, G^{(m)}$ denote the G -layers of graph $G \square H$. By S_D and S'_D denote Dominator's winning strategy on G in the D-game and in the S-game, respectively. Dominator will play his first move on one G -layer according to his winning strategy S_D . In every other round $i \geq 2$, he looks at the $(i-1)^{\text{th}}$ move of Staller. If Staller in his $(i-1)^{\text{th}}$ move claims a vertex from $V(G^{(j)})$, let Dominator respond by claiming a vertex from the same set $V(G^{(j)})$ according to the corresponding winning strategy S_D or S'_D on graph $G^{(j)}$ — the choice of S_D or S'_D depends on whether Dominator or Staller started the first move on this G -layer. Since Staller can be the first player on at most $m-1$ G -layers of graph $G \square H$, we know that Dominator can win within

$$\gamma_{MB}(G) + (m-1) \cdot \gamma'_{MB}(G)$$

many moves.

Since Dominator also has a winning strategy on H both as the first and as the second player, we obtain analogously that he can win within $\gamma_{MB}(H) + (n-1) \cdot \gamma'_{MB}(H)$ many moves. Therefore we get

$$\gamma_{MB}(G \square H) \leq \min\{\gamma_{MB}(G) + (m-1) \cdot \gamma'_{MB}(G), \gamma_{MB}(H) + (n-1) \cdot \gamma'_{MB}(H)\}.$$

Now, we consider the S-game on $G \square H$. Staller starts the game, we let Dominator respond on the same G -layer, using the winning strategy he has on G . Hence he can win within $m \cdot \gamma'_{MB}(G)$ many steps. Also, we can focus on the n H -layers, in this Cartesian graph. With the analogous analysis, we know that Dominator can win within $n \cdot \gamma'_{MB}(H)$ many steps. Hence we get

$$\gamma'_{MB}(G \square H) \leq \min\{m \cdot \gamma'_{MB}(G), n \cdot \gamma'_{MB}(H)\}.$$

Note that there can be a situation in any of the above cases, where Dominator cannot respond to the suggested strategy since he has already dominated every vertex in the corresponding G -layer. In this case, either he skip the planned move or not does not influence the situation of the game. The number we obtained above is in the situation where (we imagine that) he would not skip any move. So, we know that, in reality, he needs even less moves to reach the winning status — this is due to the “No-Skip Lemma” (Lemma 1.2). Therefore, the above inequalities still hold, under these situations. \square

We have the following corollary from Theorem 1.3.

Corollary 2.1. *Let G be a graph on n vertices. If Dominator has a winning strategy both as the first and as the second player in the game on G , then*

$$\gamma_{MB}(G \square K_2) \leq \min\{\gamma_{MB}(G) + \gamma'_{MB}(G), n\},$$

and

$$\gamma'_{MB}(G \square K_2) \leq \min\{2 \cdot \gamma'_{MB}(G), n\}.$$

3 MBD game on $P_2 \square P_n$

In order to prove our main results on MBD games on grids, namely Theorem 1.4 and Theorem 1.5, we need to introduce several “MBD graphs” and also two types of traps that Staller can make so as to win the game, as preparations.

3.1 MBD graphs

An *MBD graph* is a pair (G, \mathcal{I}) , where $G = (V, E)$ is a graph and function

$$\mathcal{I}: V \rightarrow \{s, d, n\} \times \{0, 1\}$$

assigns to each vertex a pair from $\{s, d, n\} \times \{0, 1\}$, describing the current situation of the vertex. Note that many MBD graphs defined in this section are standard partially dominated graphs. However, some are not; also, we intend to have a consistent notation for all the graphs that show up in this section, and in precise mathematical language as well. Hence, we introduce this notion of “MBD graphs”.

To say it more intuitively, s means the vertex is already claimed by Staller, d means the vertex is already claimed by Dominator. And n refers to “null” which means that the vertex is still free — not yet claimed by any player in the current game. While 1 means that the vertex is already dominated, i.e., it has a neighbouring vertex with the assigned value d , and 0 means that the vertex is not yet dominated, i.e., none of its neighbours have been assigned to with the value d yet.

An *MBD subgraph* S of a graph G is a subgraph of G which is an MBD graph, meaning that it has the value function $\mathcal{I}: V(S) \rightarrow \{s, d, n\} \times \{0, 1\}$ attached to it. In the sequel, we define several MBD subgraphs of the Cartesian graph $P_2 \square P_m$. They will be helpful in our later proofs. Let

$$V_m = \{u_1, \dots, u_m, v_1, \dots, v_m\}$$

and

$$E_m = \{\{u_i, u_{i+1}\} \mid i = 1, 2, \dots, m-1\} \cup \{\{v_i, v_{i+1}\} : i = 1, 2, \dots, m-1\} \\ \cup \{\{u_i, v_i\} : i = 1, 2, \dots, m\}.$$

It is easy to check that $P_2 \square P_m = (V_m, E_m)$.

1. By \mathfrak{X}_m ($m \geq 1$) denote the MBD graph (G, \mathcal{I}) with $G = P_2 \square P_m$, \mathcal{I} assigns to vertex u_1 with $(n, 1)$ and to all other vertices with the value $(n, 0)$. That is to say, vertex u_1 is already dominated in the ongoing game. Therefore, in the remaining game, Dominator does not need to consider dominating u_1 . In the standard notation (see [Gledel et al., 2019]), \mathfrak{X}_m is a partially dominated graph which is denoted by $G|\{u_1\}$. (see Figure 1(a))
2. By \mathfrak{Y}_m ($m \geq 3$) denote the MBD graph (G, \mathcal{I}) with $G = P_2 \square P_m$, \mathcal{I} assigns to vertex u_1 value $(n, 1)$, to vertex v_2 value $(s, 0)$, to vertices u_m, v_m the value $(n, 1)$ and to all other vertices the value $(n, 0)$. When considering the D-game on \mathfrak{Y}_m , we set $S_0 = v_2$ — consider the MBD game played on this graph, keeping in mind that Staller has claimed the vertex v_2 and vertices u_m and v_m are already dominated (but not claimed by any players) before the game starts. (see Figure 1(b))
3. By \mathfrak{Z}_m ($m \geq 1$) denote the MBD graph (G, \mathcal{I}) with $G = P_2 \square P_m$, \mathcal{I} assigns to vertices u_1 and v_1 the value $(n, 1)$, to all other vertices value $(n, 0)$. That is to say, vertices u_1 and v_1 are already dominated in the ongoing game. Therefore, in the remaining game, Dominator does not need to consider dominating u_1 or v_1 . In the standard notation (see [Gledel et al., 2019]), \mathfrak{Z}_m is a partially dominated graph which is denoted by $G|\{u_1, v_1\}$. (see Figure 1(c))

4. By \mathfrak{W}_m ($m \geq 1$) denote the MBD graph (G, \mathcal{I}) , where $G = (V, E)$, and

$$V = V_m \cup \{v_0\}, E = E_m \cup \{\{v_0, v_1\}\};$$

$\mathcal{I}(v_0) = \mathcal{I}(u_1) = (n, 1)$, and \mathcal{I} assigns to all other vertices value $(n, 0)$. That is to say, v_0 and u_1 are already dominated before the game starts. In the standard notation (see [Gledel et al., 2019]), \mathfrak{W}_m is a partially dominated graph which is denoted by $G|\{v_0, u_1\}$. (see Figure 1(d))

5. By \mathfrak{R}_m ($m \geq 0$) denote the MBD graph (G, \mathcal{I}) with $G = P_2 \square P_m$ and

$$\mathcal{I}(u_1) = (n, 1), \mathcal{I}(v_2) = (s, 0).$$

The function \mathcal{I} assigns to all other vertices value $(n, 0)$. When considering the D-game on \mathfrak{R}_m , we set $S_0 = v_2$ — consider the MBD game played on this graph, but Staller has claimed the vertex v_2 and u_1 is already dominated by Dominator before the game starts. Note that we use S_0 to denote the “move” of Staller before the game starts. Note that here we allow $m = 0$, where \mathfrak{R}_0 simply refers to the null graph; this allowance aims to make the later-on proof more smooth. When $m = 1$, it just refers to the complete graph with two vertices. (Figure 1(e)).

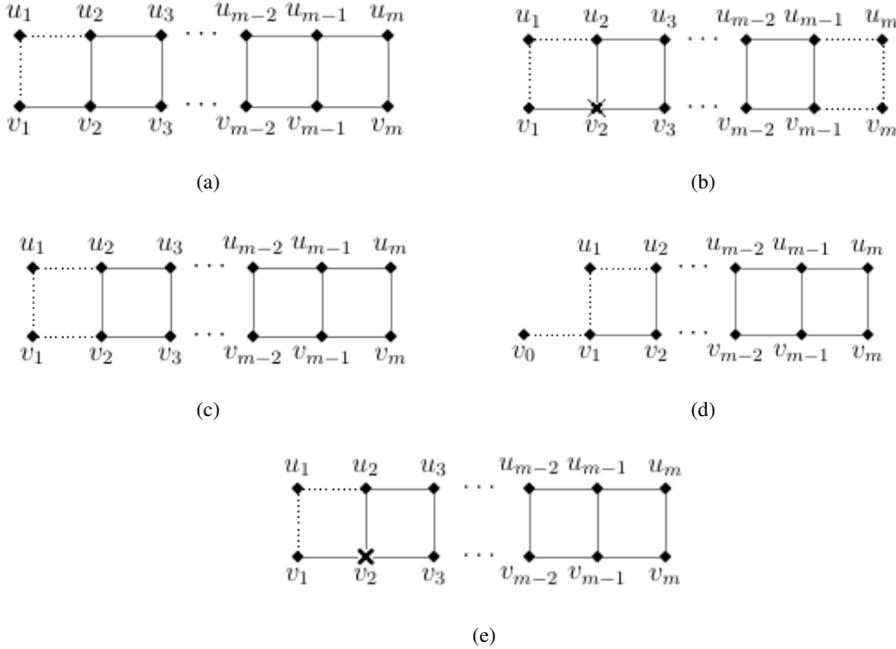


Fig. 1: Sub-figures: (a) \mathfrak{X}_m (b) \mathfrak{Y}_m (c) \mathfrak{Z}_m (d) \mathfrak{W}_m (e) \mathfrak{R}_m . Two incident edges being dotted indicates that the vertex is already dominated by Dominator; the cross indicates that the vertex is claimed by Staller.

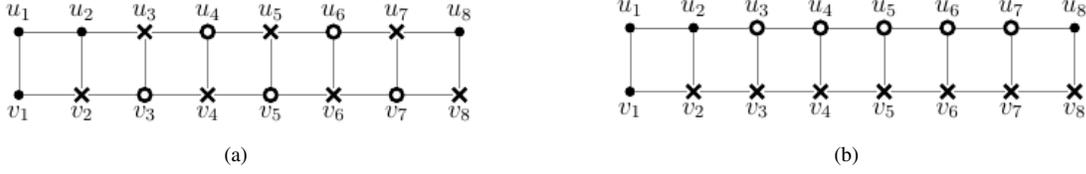


Fig. 2: Illustrations of the two type of traps. (a) Sequence of triangle traps $v_3 - -v_7$; (b) Sequence of line traps $u_3 - -u_7$.

We define the *Maker—Breaker domination number of an MBD graph* — abbreviated as “domination number of an MBD graph”, in the later context of this paper — naturally as follows. Dominator intends to dominate all undominated vertices, namely those vertices with the second coordinate of $\mathcal{I}(v)$ being 0; both players can only claim the unclaimed vertices, namely vertices v with the first coordinate of $\mathcal{I}(v)$ being n . We use the same notations γ_{MB} and γ'_{MB} for the minimum number of moves for Dominator to guarantee his winning in D-game and S-game on the MBD graph, respectively. Actually a normal graph can be viewed as a special case of an MBD graph, namely with all vertices assigned to with value $(n, 0)$ by the function \mathcal{I} .

Now we define two types of traps that Staller can create in the MBD game on $P_2 \square P_n$ for $n \geq 3$, so as to prevent the winning of Dominator. These two strategies of Staller will be used a lot later in the proofs.

Trap 1 - triangle trap. We say that Staller created a *triangle trap* if after her move Dominator is forced to claim the vertex v_i (or u_i) — so as to dominate v_i (or u_i). This is because all neighbouring vertices of v_i (or u_i) are already claimed by Staller and she can isolate v_i (or u_i) by claiming it in her next move, if Dominator did not claim it. We say that Staller created a *sequence of triangle traps* $v_i - -v_j$ (or $v_i - -u_j$), $i < j$, if Dominator is consecutively forced to claim vertices $v_i, u_{i+1}, v_{i+2}, u_{i+3}, \dots, v_j$ (or $v_i, u_{i+1}, v_{i+2}, u_{i+3}, \dots, u_j$) because of the triangle traps one after another set up by Staller.

The sequence of triangle traps $v_3 - -v_7$ is illustrated on Figure 2(a). The first trap shows up when Staller has claimed v_2, u_3 and v_4 , which forces Dominator to claim v_3 ; immediately after, Staller claims u_5 to force Dominator to claim u_4 , which creates the second trap. Notice that after Staller claims v_8 , Dominator has to claim v_7 ; but Staller can then claim u_8 , which will isolate u_8 and further leads to her winning of the game. This strategy will be applied a lot later on in our proofs. The sequences of triangle traps $u_i - -v_j$ and $u_i - -u_j$ are defined analogously.

Trap 2 - line trap. We say that Staller created a *line trap* if after her move Dominator is forced to claim the vertex v_i (or u_i) — so as to dominate u_i (or v_i). This is because all other neighbours of u_i (or v_i) and u_i (or v_i) itself are already claimed by Staller. Staller can isolate u_i (or v_i) by claiming v_i (u_i) in her next move. We say that Staller creates a *sequence of line traps* $v_i - -v_j$ (or $u_i - -u_j$), $i < j$, if Dominator is consecutively forced to claim the vertices $v_i, v_{i+1}, v_{i+2}, v_{i+3}, \dots, v_j$ (or $u_i, u_{i+1}, u_{i+2}, u_{i+3}, \dots, u_j$) because of the line traps one after another set up by Staller.

The sequence of line traps $u_3 - -u_7$ is illustrated on Figure 2(b). The first trap shows up when Staller has claimed v_2, v_3 and v_4 , which forces Dominator to claim u_3 (in order to dominate v_3); directly after, Staller claims v_5 , which forces Dominator to claim u_4 , which creates the second trap. Notice that after Dominator claims u_7 , Staller can then claim u_8 , which makes it impossible for Dominator to dominate v_8 anymore, then win the game. This strategy will be employed a lot later on in our proofs.

Next, let us briefly recall the *pairing strategy*. We follow the content provided in section 2 of [Gledel et al., 2019]. Let G be a graph, $k \geq 1$, and $u_1, \dots, u_k, v_1, \dots, v_k$ pairwise different vertices of G . Then we say that

$$X = \{\{u_1, v_1\}, \dots, \{u_k, v_k\}\}$$

is a *pairing dominating set* if

$$\bigcup_{i=1}^k N[u_i, v_i] = V(G),$$

where $N[u_i, v_i]$ means the set of common neighbouring vertices of u_i and v_i . Note that here the set of “neighbouring vertices” of vertex v includes the vertex v itself. When a graph has a pairing dominating set, then it is not hard to see that Dominator has a winning strategy in both S-game and D-game. He just need to claim the remaining vertex in the pair whenever there is one vertex left in a pair. We call this strategy the **pairing strategy**. Note that in this paper, the pairing strategy is adopted many times; most of the time it is adopted in order to prove the upper bound for the number of steps that Dominator needs in order to win. However, sometimes Dominator cannot conduct the game as suggested since the vertex that he ought to claim will not dominate any new vertices. In this case, we let Dominator skip that move and continue the game. By the “No-Skip Lemma” (Lemma 1.2), we know that Dominator — in a real optimal game — would need less steps to win. Hence the upper bound we tried to prove holds still. We only state this point once here, and omit it later on in the context where the pairing strategy are mentioned.

With this we finished introducing the basic definitions which we will need later on in the proofs of our main results.

3.2 Maker–Breaker domination numbers of MBD graphs

In this section, we introduce some results on the Maker–Breaker domination numbers on MBD graphs, which then will lead to the completion of the proofs of Theorem 1.4 and Theorem 1.5. However, the proofs of Theorem 1.4 and Theorem 1.5 will be given in Section 3. We first list the main results for this section, and then we give the proofs.

3.2.1 Domination numbers of MBD graphs: main results

Proposition 3.1. $\gamma_{MB}(\mathfrak{R}_m) = m$ for $m \geq 0$ — note that \mathfrak{R}_0 is simply the null graph; when $m \geq 2$, Dominator would lose the game if he skipped any moves in the game.

Proposition 3.2. $\gamma_{MB}(\mathfrak{Q}_m) = m - 1$, $m \geq 3$.

Proposition 3.3. $\gamma'_{MB}(\mathfrak{Z}_m) = m - 1$, $m \geq 1$.

Proposition 3.4. $\gamma'_{MB}(\mathfrak{W}_m) = m$ for $1 \leq m \leq 3$, and $\gamma'_{MB}(\mathfrak{W}_m) = m - 1$ for $m \geq 4$.

Theorem 3.5. $\gamma_{MB}(\mathfrak{X}_1) = 1$, $\gamma_{MB}(\mathfrak{X}_m) = m - 1$ for $2 \leq m \leq 5$, and $\gamma_{MB}(\mathfrak{X}_m) = m - 2$ for $m \geq 6$.

3.2.2 Domination numbers of MBD graphs: proofs

Proof of Proposition 3.1: First, the following five statements are not hard to verify:

1. $\gamma_{MB}(\mathfrak{R}_0) = 0$;
2. $\gamma_{MB}(\mathfrak{R}_1) = 1$;

3. $\gamma_{MB}(\mathfrak{R}_2) = 2$;
4. Dominator would lose in the D-games on graph \mathfrak{R}_2 if he skipped any moves;
5. On \mathfrak{R}_1 , Dominator would lose the S-game if he skipped any moves.

Assume that $\gamma_{MB}(\mathfrak{R}_k) \leq k$ holds when $2 \leq k \leq m - 1$, then when $k = m$, let Dominator adopt the strategy for \mathfrak{R}_{m-1} on the (left) subgraph \mathfrak{R}_{m-1} , let him adopt the pairing strategy for the K_2 (complete graph with two vertices) on the right. We then obtain that

$$\gamma_{MB}(\mathfrak{R}_m) \leq \gamma_{MB}(\mathfrak{R}_{m-1}) + 1 \leq m.$$

Hence by induction we obtain that $\gamma_{MB}(\mathfrak{R}_m) \leq m$ for any $m \geq 0$. We only need to show that $\gamma_{MB}(\mathfrak{R}_m) \geq m$ and Dominator would not win if he skipped any moves in the game. Assume that $\gamma_{MB}(\mathfrak{R}_k) \geq k$ and that Dominator would not win if he skipped any moves in D-game on graph \mathfrak{R}_k for any $k < m$. Now we consider the situation when $k = m$; notice that here we can assume $m \geq 3$.

In the case if Dominator had skipped the first move — this situation is denoted by $D_1 = \emptyset$ — we propose the following strategy for Staller: Let $S_1 = v_1$, which forces D_2 to be u_1 . Then let Staller make line trap from v_3 to v_n , which forces Dominator to respond on u_2 until on u_{m-1} . Then let Staller claim u_m , Dominator cannot dominate v_m anymore. Hence he would lose the game if Dominator skipped the first move.

So now we consider all possibilities for D_1 . And we will propose correspondingly Staller's strategy.

- Case 1: $D_1 = u_i, i \geq 2$. Let $S_1 = v_1$ which forces $D_2 = u_1$. Then if $i \neq 2$, let Staller make line trap from v_3 to v_i , which forces Dominator to claim from u_2 until u_{i-1} . If $i = m$, then we already see that Dominator won with m moves. If $i \notin \{m, m - 1\}$, then let Staller claim v_{i+2} , on the right we obtain graph \mathfrak{R}_{m-i} . If $m - i = 1$, just let Staller claim any remaining vertex. By induction we know that Dominator would not skip any moves and he needs m moves to win.
- Case 2: $D_1 = v_i, i \geq 3$. Suppose $i > 5$. Let $S_1 = u_2$, which forces $D_2 \in \{v_1, u_1\}$. But no matter which choice D_1 was assigned to, Staller will claim v_3, u_3 in the next two rounds so that either Dominator cannot dominate v_4 or he cannot dominate u_4 . So he would lose the game, therefore this is not an optimal choice for Dominator. Hence $i \in \{3, 4\}$.

Case 2.1: $i = 3$. Staller will then claim u_1 and u_3 , which forces Dominator to respond on v_1 and u_2 , then Staller will claim u_5 , on the right we can use the induction hypothesis.

Case 2.2: $i = 4$. Let Staller claim u_2 , then Dominator needs to claim u_1 or v_1 .

When $D_2 = u_1$, then let Staller claim v_3 which forces $D_3 = v_1$. Then let $S_3 = u_4$, which forces $D_4 = u_3$. Then $S_4 = u_6$, we can use induction hypothesis for the graph on the right.

When $D_2 = v_1$, then let Staller claim u_3 which forces $D_3 = u_1$. Then let $S_3 = u_4$ which forces $D_4 = v_3$. Then $S_4 = u_6$, we can use the induction hypothesis for the graph on the right.

- Case 3: $D_1 = u_1$, let $S_1 = v_3$. In this case if $D_2 = \emptyset$, Staller could then claim u_2 which would force Dominator to claim v_1 . Then let Staller claim u_3 , which gives Dominator a dilemma of not being able to claim u_4 and v_4 at one step: he would fail. Hence Dominator would lose the game if skipped the move D_2 . For the same reason, we know that $D_2 \in \{v_1, u_2, u_3, u_4, v_4\}$. However

if $D_2 \in \{v_1, u_2\}$, then Staller could make line traps by claiming v_4 until v_n , then at the last step claiming u_n , which would make Dominator lose the game. Hence $D_2 \in \{u_3, u_4, v_4\}$.

Case 3.1: $D_2 = u_3$. Let $S_2 = u_2$, forcing $D_3 = v_1$. Then $S_3 = v_5$, we can use the induction hypothesis for the graph on the right.

Case 3.2: $D_2 = u_4$. Let $S_2 = u_2$, forcing $D_3 = v_1$. Let $S_3 = v_4$, forcing $D_4 = u_3$. Then $S_4 = v_6$, we can use the induction hypothesis for the graph on the right.

Case 3.3: $D_2 = v_4$. Let $S_2 = u_2$, forcing $D_3 = v_1$. Let $S_3 = u_4$, forcing $D_4 = u_3$. Then $S_4 = u_6$, we can use induction for the graph on the right.

- Case 4: $D_1 = v_1$, let $S_1 = u_3$. In this case if $D_2 = \emptyset$, Staller could then claim u_2 which would force $D_3 = u_1$. Then let Staller claim v_3 ; we see that Dominator would not be able to dominate both u_3 and v_3 within only one step. So he would lose the game if he skipped D_2 . For the same reason, we know that $D_2 \in \{u_1, u_2, v_3, u_4, v_4\}$. However if $D_2 \in \{u_1, u_2\}$, then Staller could make triangle traps by claiming v_4 until v_n/u_n (depending on whether m is odd or even), then at the last step claiming u_m/v_m , which would make Dominator lose the game. Hence $D_2 \in \{v_3, u_4, v_4\}$.

Case 4.1: $D_2 = v_3$. Let $S_2 = u_2$, forcing $D_3 = u_1$. Then $S_3 = u_5$, we can use the induction hypothesis for the graph on the right.

Case 4.2: $D_2 = u_4$. Let $S_2 = u_2$, forcing $D_3 = u_1$. Let $S_3 = v_4$, forcing $D_4 = v_3$. Then $S_4 = v_6$, we can use the induction hypothesis for the graph on the right.

Case 4.3: $D_2 = v_4$. Let $S_2 = u_2$, forcing $D_3 = u_1$. Let $S_3 = u_4$, forcing $D_4 = v_3$. Then $S_4 = u_6$, we can use the induction hypothesis for the graph on the right.

Notice that in some case it may happen that the graph is not big enough for the strategy of Staller, but it is not hard to verify that in those cases we simply have an \mathfrak{R}_0 or an \mathfrak{R}_1 on the right, still we can use the induction hypothesis.

Hence notice that in all the above cases Dominator in total needs m moves to win and shall not skip any moves during the game. By induction, we conclude that for any $m \geq 2$, in the D-game on graph \mathfrak{R}_m , Dominator needs at least m steps to win and he shall not skip any moves in the game. \square

Remark 3.6. Note that the D-game on graph \mathfrak{R}_m can be considered as the S-game on \mathfrak{X}_m with $S'_1 = v_2$.

Proof of Proposition 3.2: Vertex v_2 is pre-claimed by Staller: we denote this by $S_0 = v_2$. The proof and case analysis can be done analogously to those of Proposition 3.1. \square

Proof of Proposition 3.3: It is not hard to check that $\gamma'_{MB}(\mathfrak{Z}_m) = m - 1$ holds for $m = 1, 2, 3$. Assume that $\gamma'_{MB}(\mathfrak{Z}_k) \leq k - 1$ for all $3 \leq k < m$. When $k = m$, we view graph \mathfrak{Z}_m as two parts — \mathfrak{Z}_{m-1} , and u_m, v_m together with the edge connecting them. Of course these two parts are connected via two edges in \mathfrak{Z}_m , namely $\{u_{m-1}, u_m\}$ and $\{v_{m-1}, v_m\}$ — but we do not consider them for now, since the extra edges will just make it easier for Dominator to win. We propose the following strategy for Dominator. If Staller plays on the \mathfrak{Z}_{m-1} part of the graph, let Dominator respond on \mathfrak{Z}_{m-1} with his strategy on \mathfrak{Z}_{m-1} ; if Staller plays on the other part, Dominator will use the pairing strategy, namely claim the other vertex in this part. In this way, by the induction hypothesis, Dominator can win within

$$\gamma'_{MB}(\mathfrak{Z}_{m-1}) + 1 \leq (m - 2) + 1 = m - 1$$

steps. Hence,

$$\gamma'_{MB}(\mathfrak{Z}_m) \leq m - 1, \quad m \geq 1.$$

For the other direction of the proof, we need to show that $\gamma'_{MB}(\mathfrak{Z}_m) \geq m - 1$ holds for $m \geq 4$. We prove it by proposing a strategy for Staller. Let $S'_1 = u_m$, then it has to be that $D'_1 \in \{u_{m-1}, v_{m-1}, v_m\}$; otherwise $S'_2 = v_m$, then at least one of v_m and u_m can get isolated after S'_3 (Staller's third step) — after D'_2 (Dominator's second step) at least one of u_{m-1} and v_{m-1} will still be free, so Staller can choose this free vertex at her third step. Now we make case distinctions for this three choices of D'_1 .

1. $D'_1 = v_{m-1}$. Let $S'_2 = u_{m-1}$: this forces $D'_2 = v_m$ (in order to dominate u_m). Let $S'_3 = u_{m-3}$. The remaining MBD graph in this game is a \mathfrak{Y}_{m-2} with vertex set

$$V(\mathfrak{Y}_{m-2}) = \{u_1, u_2, \dots, u_{m-2}, v_1, v_2, \dots, v_{m-2}\}.$$

It is Dominator's turn now. Therefore he needs $\gamma_{MB}(\mathfrak{Y}_{m-2}) = m - 3$ (by Proposition 3.2) more moves to win. So he needs in total $m - 1$ steps to win in this case.

2. $D'_1 = u_{m-1}$. Let $S'_2 = v_{m-1}$: this forces $D'_2 = v_m$ (in order to dominate v_m). Let $S'_3 = v_{m-3}$. Remaining part of this case is the same as the last case. Dominator needs in total $m - 1$ steps to win.
3. $D'_1 = v_m$. Let $S'_2 = u_{m-2}$: the remaining MBD graph in this game is a \mathfrak{Y}_{m-1} . Dominator needs $\gamma_{MB}(\mathfrak{Y}_{m-1}) = m - 2$ (by Proposition 3.2) many moves to win on this part and he needs in total $(m - 2) + 1 = m - 1$ many moves to win in this case.

Hence, Staller has a strategy such that Dominator needs at least $m - 1$ steps in order to win in an S-game on the MBD graph \mathfrak{Z}_m . We obtain that $\gamma'_{MB}(\mathfrak{Z}_m) \geq m - 1, m \geq 4$. \square

Proof of Proposition 3.4: One can check that

$$\gamma'_{MB}(\mathfrak{W}_m) = m, \quad 1 \leq m \leq 3.$$

Let $m \geq 4$. Since \mathfrak{W}_m has one more undominated vertex than \mathfrak{Z}_m (namely v_1), Dominator needs to play at least as many moves on \mathfrak{W}_m as on \mathfrak{Z}_m . Therefore we have

$$\gamma'_{MB}(\mathfrak{W}_m) \geq \gamma'_{MB}(\mathfrak{Z}_m) = m - 1,$$

which shows the lower bound. For the upper bound, we start our consideration from $m = 4$. When $m = 4$, we propose strategies for Dominator accordingly to the steps of Staller.

- Case 1: $S'_1 = v_2$. Let $D'_1 = u_3$.
 If $S'_2 = v_1$, then let $D'_2 = v_3$. Dominator only needs one step to dominate v_1 before his winning; so as to do this, he only needs to claim v_0 or u_1 in the next move. So he needs in total 3 moves to win.
 If $S'_2 \neq v_1$, then let $D'_2 = v_1$. Then Dominator just needs one more step to dominate v_4 . In total he needs 3 steps to win.

- Case 2: $S'_1 \neq v_2$.

If $S'_1 = u_4$ (or v_4). Let $D'_1 = u_3$. If $S'_2 = v_3$, then let $D'_2 = v_4$ (or u_4) and let $D'_3 \in \{v_1, v_2\}$. If $S'_2 = v_4$ (or u_4), let $D'_2 = v_3$ and $D'_3 \in \{v_1, v_2\}$. So Dominator needs in total 3 steps to win.

If $S'_1 \notin \{u_4, v_4\}$. Let $D'_1 = v_2$. Dominator needs at most two more moves to dominate the remaining vertices.

So far, we obtain that $\gamma'_{MB}(\mathfrak{W}_4) \leq 3$. Suppose that

$$\gamma'_{MB}(\mathfrak{W}_{m-1}) \leq m - 2 \text{ for } m \geq 5.$$

Then consider the MBD graph \mathfrak{W}_m ($m \geq 5$) as two parts (of course plus the edges connecting the two parts): one is the rightmost K_2 , namely u_m, v_m and the edge connecting them; the other is on the left an MBD subgraph \mathfrak{W}_{m-1} . We let Dominator react on the left \mathfrak{W}_{m-1} according to his strategy on \mathfrak{W}_{m-1} whenever Staller claims a vertex on this subgraph; otherwise we let Dominator claim the remaining vertex among u_m and v_m . By induction, Dominator needs no more than $(m - 2) + 1 = m - 1$ moves to win. Hence $\gamma'_{MB}(\mathfrak{W}_m) \leq m - 1$. \square

Theorem 3.5 is the essence for proving Theorem 1.5. The case analyses in the proof can be lengthy. We provide a strategy tree (see Figure 3) showing the idea of the case studies of Theorem 3.5, which, although not covering all cases, but can at least express the idea behind. We leave it as an exercise for readers, to complete the whole strategy tree taking into consideration of all cases.

Proof of Theorem 3.5: For $m \in \{1, 2, 3\}$ the situation is not hard to directly see. For $m = 4$ and $m = 5$ simple cases analysis gives the result. Let $m \geq 6$. First, we consider the D-game on \mathfrak{X}_6 . Let $D_1 = v_2$, then we see an MBD subgraph \mathfrak{W}_4 on the rightmost which already contains all the undominated vertices. By Proposition 3.4, we know that $\gamma'_{MB}(\mathfrak{W}_4) = 3$. Hence, $\gamma_{MB}(\mathfrak{X}_6) = 4$.

Suppose that $\gamma_{MB}(\mathfrak{X}_k) \leq k - 2$ for $k < m$. When $k = m$, consider the graph as two parts (plus edges connecting the two parts): one is the K_2 on the rightmost, the other is the MBD subgraph \mathfrak{X}_{m-1} on the left. Let Dominator adopt the strategy for \mathfrak{X}_{m-1} on the left subgraph whenever Staller plays on that subgraph and let him claim the remaining vertex among u_m and v_m when Staller claims one of them. In this way, by the induction hypothesis we know that Dominator needs at most $(m - 3) + 1 = m - 2$ steps to win. Hence $\gamma_{MB}(\mathfrak{X}_m) \leq m - 2$ for $m \geq 6$.

We prove the lower bound by induction and proposing strategies for Staller. Suppose that

$$\gamma_{MB}(\mathfrak{X}_k) \geq k - 2, \text{ for } 4 \leq k < m, m \geq 6.$$

Now consider the case when $k = m, m \geq 6$. We want to show that

$$\gamma_{MB}(\mathfrak{X}_m) \geq m - 2, m \geq 6.$$

According to the first step of Dominator, there are three cases: $D_1 \in \{u_1, v_1, u_2, v_2\}$, $D_1 = u_i$ ($i \geq 3$), or $D_1 = v_i$ ($i \geq 3$).

- **Case 1:** $D_1 \in \{u_1, v_1, u_2, v_2\}$.

If $D_1 \in \{u_1, v_1\}$, then we see that the remaining game is the S-game on the MBD subgraph \mathfrak{W}_{m-1} . By Proposition 3.4, $\gamma'_{MB}(\mathfrak{W}_m) = m - 1$ when $m \geq 4$. Hence Dominator needs $m - 2$ more steps to win; he needs in total $m - 1$ steps to win.

If $D_1 = v_2$, then we see that on the right is an MBD subgraph \mathfrak{W}_{m-2} . By Proposition 3.4, Dominator needs in total $(m-3) + 1 = m-2$ moves to win.

If $D_1 = u_2$, then on the left is the MBD subgraph \mathfrak{W}_{m-2} , while on the right Dominator needs at most one more step, so as to dominate v_1 . Hence Dominator needs at least $(m-3) + 1 = m-2$ steps to win.

- **Case 2:** $D_1 = u_i, i \geq 3$. Let $S_1 = v_2$, then according to the second move of Dominator, we carry out case distinctions.

(2.1) $D_2 = u_1$. Then we further consider two sub-cases: $i = 3$ or $i \geq 4$.

(2.1.1) If $D_1 = u_3$, then we see that the remaining game is the S-game on the MBD subgraph \mathfrak{W}_{m-3} on the left of \mathfrak{X}_m ; where on the right Dominator needs at most one more step, so as to dominate v_2 . Hence he needs in total at least $(m-4) + 2 = m-2$ steps to win.

(2.1.2) If $D_1 = u_i, i \geq 4$, we let $S_2 = v_3$. We need to further distinct cases according to the choice of D_3 .

(2.1.2.1) $D_3 \in \{v_1, u_2\}$. Let Staller take the strategy of creating a sequence of line traps $u_3 - u_{i-1}$ which ends with $S_{i-1} = v_i$ and forces $D_i = u_{i-1}$. Then, if $m-i \geq 2$, let $S_i = v_{i+2}$. We see that the rightmost is the MBD graph \mathfrak{R}_{m-i} . Dominator wins after the D-game on the rightmost \mathfrak{R}_{m-i} . By Proposition 3.1, $\gamma_{MB}(\mathfrak{R}_{m-i}) = m-i$. Hence Dominator needs in total m moves to win. If $m-i = 1$, then let $S_i \in \{u_m, v_m\}$, Dominator needs only one more step to win. If $m-i = 0$, the game finished already. In either case, Dominator wins with m steps.

(2.1.2.2) $D_3 = u_j, j \geq 3$ or $D_3 = v_j, j \geq 4$. Suppose that $\min\{i, j\} \notin \{3, 4\}$. Let $S_3 = u_2$, which forces $D_4 = v_1$. Then let $S_4 = u_3$, Dominator cannot dominate both u_4 and v_4 , hence he would not be able to win. Therefore, $\min\{i, j\} \in \{3, 4\}$.

(2.1.2.2.a) When $D_3 = u_j$, where $j \geq 3$ and $j < i$, we know that $j \in \{3, 4\}$. When $j = 3$, we have $D_3 = u_3$. Recall that $D_1 = u_i, i \geq 4, S_1 = v_2, D_2 = u_1, S_2 = v_3, D_3 = u_3$. Now let $S_3 = u_2$, which forces $D_4 = v_1$. We see on the rightmost the MBD graph \mathfrak{X}_{m-3} , but, with u_i being claimed by Dominator already. Now it is Staller's turn. By the induction hypothesis, we know that Dominator needs at least $\gamma_{MB}(\mathfrak{X}_{m-3}) - 1 \geq m-6$ more steps to win the game. In total, he needs $(m-6) + 4 = m-2$ steps.

(2.1.2.2.b) When $j = 4$, we have $D_3 = u_4$. Let $S_3 = u_2$, which forces $D_4 = v_1$. Then let $S_4 = v_4$, which forces $D_5 = u_3$. Similarly as in the last case, the remaining graph is the MBD graph \mathfrak{X}_{m-4} on the right, but with u_i being claimed by Dominator already. Now it is Staller's turn. By the induction hypothesis, we know that Dominator needs at least $\gamma_{MB}(\mathfrak{X}_{m-4}) - 1 \geq m-7$ more steps to win the game. In total, he needs $(m-7) + 5 = m-2$ steps.

(2.1.2.2.c) When $D_3 = u_j$, where $j > i$, we know that $i \in \{3, 4\}$. Since $i \geq 4$, we know that $D_1 = u_i = u_4$. Let Staller take the same strategy as in the last case, we get that Dominator needs at least $m-2$ steps to win.

(2.1.2.2.d) When $D_3 = v_j$, where $j < i$, since $S_2 = v_3$, we know that $j = 4$, i.e., $D_3 = v_4$. Let $S_3 = u_2$, which forces $D_4 = v_1$. Then let $S_4 = u_4$, which forces $D_5 = u_3$. The remaining graph on the right is the MBD graph \mathfrak{X}_{m-4} , but with u_i claimed by Dominator already, and now it is Staller's turn. By the induction hypothesis, we know that Dominator needs at least $\gamma_{MB}(\mathfrak{X}_{m-4}) - 1 \geq m - 7$ more steps to win the game. In total, he needs $(m - 7) + 5 = m - 2$ steps.

(2.1.2.2.e) When $D_3 = v_j$, where $j = i$, since $S_2 = v_3$, we know that $j = i = 4$, that is, $D_1 = u_4$ and $D_3 = v_4$. Let $S_3 = v_2$, which forces $D_4 = v_1$. The remaining graph on the right is the MBD graph \mathfrak{Z}_{m-4} , and now it is Staller's turn. By Proposition 3.3, Dominator needs in $\gamma'_{MB}(\mathfrak{Z}_{m-4}) = m - 5$ more steps to win. In total, he needs $(m - 5) + 4 = m - 1$ steps to win.

(2.1.2.2.f) When $D_3 = v_j$, where $j > i$, then $i \in \{3, 4\}$, that is, $D_1 \in \{u_3, u_4\}$. The analyses are analogous to items (a) and (b). Dominator needs in total at least $m - 2$ steps to win.

(2.2) $D_2 = u_2$. Recall that we are currently under the case where $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$. Let $S_2 = u_1$, which forces $D_3 = v_1$. The remaining graph is the MBD graph \mathfrak{X}_{m-2} on the right, but with u_i claimed already by Dominator, and now it is Staller's turn. By the induction hypothesis, Dominator needs at least $m - 5$ more steps to win. In total, he needs at least $(m - 5) + 3 = m - 2$ steps to win.

(2.3) $D_2 = u_j$, where $j \geq 3$. Recall that we are currently under the case where $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$. Now let $S_2 = v_1$, which forces $D_3 = u_1$. Let $l := \min\{i, j\}$ and $h := \max\{i, j\}$. Let Staller create a sequence of line traps $u_1 - \dots - u_{l-1}$. The remaining graph is the MBD graph \mathfrak{X}_{m-l} on the right, but with u_h claimed already by Dominator, and now it is Staller's turn. By induction, Dominator needs at least $m - l - 2$ more steps to win. In total, he needs at least $(m - l - 2) + l = m - 2$ steps to win.

(2.4) $D_2 = v_1$.

(2.4.1) $D_1 = u_3$. Recall that $S_1 = v_2$, $D_2 = v_1$. The remaining graph is the MBD graph \mathfrak{W}_{m-3} on the right, and now it is Staller's turn. By Proposition 3.4, we know that Dominator needs $\gamma'_{MB}(\mathfrak{W}_{m-3}) \geq m - 4$ more steps to win. In total, he needs at least $(m - 4) + 2 = m - 2$ steps to win.

(2.4.2) $D_1 = u_i$, where $i \geq 4$. Let $S_2 = u_3$. We need to further distinct cases according to the choice of D_3 .

(2.4.2.1) $D_3 \in \{u_1, u_2\}$.

(2.4.2.1.a) When i is even, let Staller create a sequence of triangle traps $v_3 - \dots - v_{i-1}$, which ends with $S_{i-1} = v_i$, $D_i = v_{i-1}$. Then, let $S_i = v_{i+2}$ if $m - i \geq 2$. The remaining graph is the MBD graph \mathfrak{R}_{m-i} and now it is Dominator's turn. By Proposition 3.1, we know that Dominator needs $\gamma_{MB}(\mathfrak{R}_{m-i}) = m - i$ more steps to win. In total, he needs $(m - i) + i = m$ steps to win.

(2.4.2.1.b) When i is odd, let Staller create a sequence of triangle traps $v_3 - -v_{i-2}$, which ends with $S_{i-2} = v_{i-1}$, $D_{i-1} = v_{i-2}$. The remaining graph is the MBD graph \mathfrak{W}_{m-i} on the right and now it is Staller's turn. By Proposition 3.4, Dominator needs at least $\gamma'_{MB}(\mathfrak{W}_{m-i}) \geq m - i - 1$ more steps to win. In total, he needs at least $(m - i - 1) + i - 1 = m - 2$ steps to win.

(2.4.2.2) $D_3 = u_j$, $j \geq 4$ or $D_3 = v_j$, $j \geq 3$. A similar reasoning as in case **(2.1.2.2)** shows that $\min\{i, j\} \in \{3, 4\}$ must hold.

(2.4.2.2.a) When $D_3 = u_j$, where $j < i$; since $S_2 = u_3$, we know that $j = 4$, that is, $D_3 = u_4$. Let $S_3 = u_2$, which forces $D_4 = u_1$; then let $S_4 = v_4$, which forces $D_5 = v_3$. The remaining graph is the MBD graph \mathfrak{X}_{m-4} , but with u_j already claimed by Dominator, and now it is Staller's turn. By the induction hypothesis, Dominator needs at least $m - 6 - 1 = m - 7$ more steps to win. In total, he needs $(m - 7) + 5 = m - 2$ steps to win.

(2.4.2.2.b) When $D_3 = u_j$, where $j > i$; since $S_2 = u_3$, we know that $i = 4$, i.e., $D_1 = u_4$. With the same strategy as in item (a), we obtain that Dominator needs at least $m - 2$ steps to win.

(2.4.2.2.c) When $D_3 = v_j$, where $j > i$; since $S_2 = u_3$, we know that $i = 4$, i.e., $D_1 = u_4$. With the same strategy as in item (a), we obtain that Dominator needs at least $m - 2$ steps to win.

(2.4.2.2.d) When $D_3 = v_j$, where $j = i$; since $S_2 = u_3$, we know that $i = j = 4$, i.e., $D_1 = u_4$, $D_3 = v_4$. Let $S_3 = u_2$, which forces $D_4 = u_1$. The remaining graph is the MBD graph \mathfrak{Z}_{m-4} on the right and it is Staller's turn. By Proposition 3.3, Dominator needs at least $\gamma'_{MB}(\mathfrak{Z}_{m-4}) = m - 5$ more steps to win.

(2.4.2.2.e) When $D_3 = v_j$, where $j < i$, then $j \in \{3, 4\}$. When $j = 3$, $D_3 = v_3$. Let $S_3 = u_2$, which forces $D_4 = u_1$. Consider the remaining graph as the MBD graph \mathfrak{X}_{m-3} , but with u_i already claimed by Dominator, and now it is Staller's turn. By the induction hypothesis, Dominator needs at least $k - 6$ more steps to win. In total, he needs at least $k - 2$ steps to win. When $j = 4$, $D_3 = v_4$. Let $S_3 = u_2$, which forces $D_4 = u_1$. Then let $S_4 = u_4$, which forces $D_5 = v_3$. We get the MBD graph \mathfrak{X}_{m-4} with u_i already claimed by Dominator on the right. By induction, Dominator needs at least $k - 7$ steps to win. In total, he needs at least $k - 2$ steps to win.

(2.5) $D_2 = v_j$, where $j > i$.

Recall that we are currently under the case where $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$. Let $S_2 = v_1$, which forces $D_3 = u_1$. Then let Staller create a sequence of line traps $u_2 - -u_{i-1}$. Then the remaining graph is the MBD graph \mathfrak{X}_{m-i} on the right, but with v_j already claimed by Dominator, and now it is Staller's turn. By induction, Dominator needs at least $m - i - 3$ more steps to win. In total, he needs $(m - i - 3) + i + 1 = m - 2$ steps to win.

(2.6) $D_2 = v_j$, where $j = i \geq 3$.

Recall that $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$, $D_2 = v_i$. Let $S_2 = v_1$, which forces $D_3 = u_1$. Then let Staller create a sequence of line traps $u_2 - -u_{i-2}$. If $i = m$, then Dominator already won with m steps; otherwise, the remaining graph is the MBD graph \mathfrak{Z}_{m-i} on the right. By Proposition 3.3,

Dominator needs $\gamma'_{MB}(\mathfrak{Z}_{m-i}) = m - i - 1$ more steps to win. In total, he needs $(m - i - 1) + i = m - 1$ steps to win. In either case, he needs more than $m - 2$ steps to win.

(2.7) $D_2 = v_j$, where $i > j \geq 2$ and j is even.

Recall that $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$, so actually $j \geq 4$. Let $S_2 = u_2$, then we claim that $D_3 \in \{u_1, v_1\}$. Suppose not. After D_3 either v_3 or u_3 should be free. If v_3 is free after D_3 , let $S_3 = v_1$ and $S_4 \in \{u_1, v_3\}$. We see that Dominator cannot dominate both u_1 and v_2 , hence he would not win. If u_3 is free after D_3 , let $S_3 = u_1$ and $S_4 \in \{v_1, u_3\}$. We see that Dominator cannot dominate both v_1 and u_2 , hence he would not win.

(2.7.a) $D_3 = u_1$. Recall that $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$, $D_2 = v_j$, $i > j \geq 4$ is even, $S_2 = u_2$. Let $S_3 = v_3$, which forces $D_4 = v_1$. Then let Staller create a sequence of triangle traps $u_3 - -u_{j-1}$. The remaining graph is the MBD graph \mathfrak{X}_{m-j} on the right, but with u_i claimed already by Dominator, and now it is Staller's turn. By induction, Dominator needs at least $m - j - 2$ more steps to win. In total, he needs $(m - j - 2) + j = m - 2$ steps to win.

(2.7.b) $D_3 = v_1$. Recall that $D_1 = u_i$, $i \geq 3$, $S_1 = v_2$, $D_2 = v_j$, $i > j \geq 4$ is even, $S_2 = u_2$. Let $S_3 = u_3$, which forces $D_4 = u_1$. Actually $j = 4$ must hold, otherwise in the next move Staller can either isolate u_4 or v_4 . Let $S_4 = u_4$, which forces $D_5 = v_3$. The remaining graph is the MBD graph \mathfrak{X}_{m-4} on the right, but with u_i already claimed by Dominator, and now it is Staller's turn. By induction, Dominator needs at least $m - 7$ more steps to win. In total, he needs at least $(m - 7) + 5 = m - 2$ steps to win.

(2.8) $D_2 = v_j$, where $i > j \geq 2$ and j is odd. Let $S_2 = u_1$, which forces $D_3 = v_1$. Then let Staller create a sequence of triangle traps $u_2 - -u_{j-1}$. The remaining graph is the MBD graph \mathfrak{X}_{m-j} on the right, but with u_i already claimed by Dominator, and now it is Staller's turn. By induction, Dominator needs at least $m - j - 3$ more steps to win. In total, he needs at least $(m - j - 3) + j + 1 = m - 2$ steps to win.

- **Case 3:** $D_1 = v_i$, $i \geq 3$. Let $S_1 = v_2$, then according to the second move of Dominator, we carry out case distinctions.

(3.1) $D_2 = u_1$.

(3.1.1) If $i = 3$, i.e., $D_1 = v_3$, then the remaining graph is the MBD graph \mathfrak{W}_{m-3} on the right, and it is Staller's turn now. By Proposition 3.4, Dominator needs at least $(m - 3) - 1 = m - 4$ more steps to win. In total, he needs $(m - 4) + 2 = m - 2$ steps to win.

(3.1.2) If $i \geq 3$, let $S_2 = v_3$. We carry out case distinctions according to the third move of Dominator in the sequel.

(3.1.2.1) $D_3 = v_1$. If $i = 4$, i.e., $D_1 = v_4$, on the right we see the MBD graph \mathfrak{W}_{m-4} , and it is Staller's turn now. Dominator needs at least $(m - 4) - 1 = m - 5$ more steps to dominate those undominated vertices in this graph \mathfrak{W}_{m-4} . Hence he needs in total at least $(m - 5) + 3 = m - 2$ steps to win. Otherwise, if $i > 4$, let $S_3 = v_4$ starting a sequence of line traps

$u_3 - -u_{i-2}$. The remaining graph is the MBD graph \mathfrak{W}_{m-i} on the right, and it is Staller's turn now. By Proposition 3.4, Dominator needs at least $(m-i) - 1$ more steps to win. In total, he needs at least $(m-i-1) + (i-1) = m-2$ steps to win.

(3.1.2.2) $D_3 = v_j, j \geq 4$ or $D_3 = u_j, j \geq 3$. A similar reasoning as in case **(2.1.2.2)** shows that $\min\{i, j\} \in \{3, 4\}$ must hold.

(3.1.2.2.1) $D_3 = v_j, j \geq 4$. Denote by $l := \min\{i, j\}$, $h := \max\{i, j\}$, then we know that $l = 4$ since v_3 is claimed by the second move of Staller. Let $S_3 = u_2$, which forces $D_4 = v_1$. Then let $S_4 = u_4$, which forces $D_5 = u_3$ by making a triangle trap. The remaining graph is the MBD graph \mathfrak{X}_{m-4} on the right, but with u_h claimed by Dominator, and it is Staller's turn now. By induction, Dominator needs at least $(m-4) - 2 = m-6$ more steps to win. In total, he needs at least $(m-6) + 4 = m-2$ steps to win.

(3.1.2.2.2) $D_3 = u_j, j > i$. Since v_3 is already claimed, we know that $i = 4$, that is, $D_1 = v_4$. The reasoning for this case is analogous to the last case (case **(3.1.2.2.1)**).

(3.1.2.2.3) $D_3 = u_j, j = i$. Since v_3 is already claimed, we know that $i = j = 4$. The MBD graph on the right is \mathfrak{Z}_{m-4} , and it is Staller's turn now. By Proposition 3.3, Dominator needs at least $(m-4) - 1 = m-5$ more steps to dominate those undominated vertices in \mathfrak{Z}_{m-4} , and he needs one more step to dominate v_2 . Hence he needs at least $(m-5) + 4 = m-1$ steps to win.

(3.1.2.2.4) $D_3 = u_j, j < i$, hence $j \in \{3, 4\}$.

(3.1.2.2.4.a) If $j = 3$, i.e., $D_3 = u_3$. Let $S_3 = v_1$, which forces $D_4 = u_2$ by making a line trap. The remaining graph on the right is \mathfrak{X}_{m-3} , but with v_i claimed by Dominator already, and it is Staller's turn now. By induction, Dominator needs at least $(m-3) - 2 = m-5$ more steps to win. In total, he needs at least $(m-5) + 3 = m-2$ steps to win.

(3.1.2.2.4.b) If $j = 4$, i.e., $D_3 = u_4$. Let $S_3 = u_2$, which forces $D_4 = v_1$. Then let $S_4 = v_4$, which forces $D_5 = u_3$ by making a line trap. The remaining graph on the right is \mathfrak{X}_{m-4} , but with v_i claimed by Dominator already, and it is Staller's turn now. By induction, Dominator needs at least $(m-4) - 2 = m-6$ more steps to win. In total, he needs at least $(m-6) + 4 = m-2$ steps to win.

(3.2) $D_2 = v_1$. Let $S_2 = u_3$, we carry out case distinctions according to the third move of Dominator in the sequel.

(3.2.1) $D_3 \in \{u_1, u_2\}$.

(3.2.1.1) i is even. Let Staller create a sequence of triangle traps $v_3 - -u_{i-2}$. The remaining graph is the MBD graph \mathfrak{W}_{m-i} on the right, and it is Staller's turn now. By Proposition 3.4, Dominator needs at least $(m-i) - 1$ more steps to win. In total, he needs at least $(m-i-1) + (i-1) = m-2$ steps to win.

(3.2.1.2) i is odd. Let Staller create a sequence of triangle traps $v_3 - -u_{i-1}$, which ends with $S_{i-1} = u_i$, $D_i = u_{i-1}$. Let $S_i = u_{i+2}$, which creates the MBD graph \mathfrak{X}_{m-i} on the right, and it is Dominator's turn now. By Proposition 3.1, $\gamma_{MB}(\mathfrak{X}_{m-i}) = m - i$, hence Dominator needs at least $m - i$ steps to win. In total, he needs at least $(m - i) + i = m$ moves to win.

(3.2.2) $D_3 = u_j$, $j \geq 4$. A similar reasoning as in case **(2.1.2.2)** shows that $\min\{i, j\} \in \{3, 4\}$ must hold.

(3.2.2.a) $j < i$. Since u_3 is claimed, $j = 4$, i.e., $D_3 = u_4$. The analysis is analogous to case **(2.4.2.2.a)**.

(3.2.2.b) $j = i$. Since u_3 is claimed, $i = j = 4$, that is, $D_1 = v_4$, $D_3 = u_4$. Let $S_3 = u_2$, which forces $D_4 = u_1$. Then the remaining graph is the MBD graph \mathfrak{Z}_{m-4} on the right, and it is Staller's turn now. By Proposition 3.3, Dominator needs at least $(m - 4) - 1 = m - 5$ more steps to win. In total, he needs at least $(m - 5) + 4 = m - 1$ steps to win.

(3.2.2.c) $j > i$, hence $i \in \{3, 4\}$. If $i = 3$, that is, $D_1 = v_3$, then let $S_3 = u_2$, which forces $D_4 = u_1$. We get the MBD graph \mathfrak{X}_{m-3} on the right, but with v_j claimed already, and it is Staller's turn now. By induction, Dominator needs at least $(m - 3) - 3 = m - 6$ more steps to win. In total, he needs at least $(m - 6) + 4 = m - 2$ steps to win. If $i = 4$, that is, $D_1 = v_4$. Let $S_3 = u_2$, which forces $D_4 = u_1$. Then let $S_4 = u_4$, which forces $D_5 = v_3$ by creating a line trap. Then consider the MBD graph \mathfrak{X}_{m-4} on the right. Again, Dominator needs in total at least $m - 2$ steps to win.

(3.2.3) $D_3 = v_j$, $j \geq 3$. A similar reasoning as in case **(2.1.2.2)** shows that $\min\{i, j\} \in \{3, 4\}$ must hold.

(3.2.3.a) $j < i$. The reasoning is analogous to cases **(3.2.2.c)**.

(3.2.3.b) $j > i$. The reasoning is analogous to cases **(3.2.2.c)**.

(3.3) $D_2 = u_j$, where $j \geq 2$, $j > i$ and i is even. Let $S_2 = u_2$, the remaining reasoning is analogous to case **(2.7)**.

(3.4) $D_2 = u_j$, where $j \geq 2$, $j > i$ and i is odd. Let $S_2 = u_1$, the remaining reasoning is analogous to case **(2.8)**.

(3.5) $D_2 = u_j$, where $j = i \geq 2$. Let $S_2 = v_1$, the remaining reasoning is analogous to case **(2.6)**.

(3.6) $D_2 = u_j$, where $2 \leq j < i$. Let $S_2 = v_1$, the remaining reasoning is analogous to case **(2.5)**.

(3.7) $D_2 = v_j$, $j \geq 2$, such that $\min\{i, j\}$ is odd. Let $S_2 = u_1$, the remaining reasoning is analogous to case **(2.8)**.

(3.8) $D_2 = v_j$, $j \geq 2$, such that $\min\{i, j\}$ is even. Let $S_2 = u_2$, the remaining reasoning is analogous to case **(2.7)**.

Since we have gone through all cases, it turns out that Dominator anyway needs at least $m - 2$ steps to win the game. By induction, $\gamma_{MB}(\mathfrak{X}_m) \geq m - 2$, when $m \geq 6$. So to sum up,

$$\gamma_{MB}(\mathfrak{X}_m) = m - 2 \text{ for } m \geq 6. \quad \square$$

3.3 Maker–Breaker domination number of $P_2 \square P_n$

Proof of Theorem 1.4: To prove the upper bound $\gamma'_{MB}(P_2 \square P_n) \leq n$, let Dominator adopt the pairing strategy: whenever Staller claims u_i (v_i), let Dominator claim v_i (u_i) in the next round. It is not hard to see that Dominator will win within n moves.

To prove the lower bound $\gamma'_{MB}(P_2 \square P_n) \geq n$, we propose the following strategy for Staller. Let Staller claim v_2 as her first move if $n \geq 2$, let Staller claim u_1 when $n = 1$. The remaining game is harder for Dominator than the D-game on \mathfrak{X}_n , since \mathfrak{X}_n has one less undominated vertex than $P_2 \square P_n$ with v_2 claimed by Staller. By Proposition 3.1, we know that Dominator needs at least n more steps to win, and he does not skip any moves. When $n = 1$, it is easy to see that he shall not skip any moves. Hence $\gamma'_{MB}(\mathfrak{X}_n) \geq n$, and Dominator would not win if he skipped any moves. \square

Next, we would like to address a concrete example of $P_2 \square P_n$, namely when $n = 13$. Before which, we need two more auxiliary results.

Claim 3.7. *Consider the S-game on the MBD graph \mathfrak{W}_4 . Let Dominator skip his first move, namely $D'_1 = \emptyset$. In addition, assume that $S'_1 \notin \{u_3, v_3, u_4, v_4\}$. Then Dominator can win within 4 moves.*

Proof: The proof can be done by case distinctions, we do not go into details here. \square

Claim 3.8. *Consider the S-game on the MBD graph \mathfrak{W}_6 . Let Dominator skip his first move, namely $D'_1 = \emptyset$. In addition, assume that $S'_1 = v_2$. Then Dominator can win within 6 moves.*

Proof: The proof can be done by case distinctions, we do not go into details here. \square

Now we can prove the following result about $P_2 \square P_{13}$.

Theorem 3.9. $\gamma_{MB}(P_2 \square P_{13}) = 11$.

Proof: Since $P_2 \square P_{13}$ has one more undominated vertex than \mathfrak{X}_{13} , it follows that

$$\gamma_{MB}(P_2 \square P_{13}) \geq \gamma_{MB}(\mathfrak{X}_{13}).$$

This is because of the Continuation Principle (Remark 1.1). By Theorem 3.5, $\gamma_{MB}(\mathfrak{X}_{13}) = 11$. Hence $\gamma_{MB}(P_2 \square P_{13}) \geq 11$. Denote by $L = (V_L, E_L)$ the subgraph of $P_2 \square P_{13}$ induced by vertex set

$$\{u_1, v_1, u_2, v_2, \dots, u_6, v_6\}.$$

Denote by $R = (V_R, E_R)$ the subgraph of $P_2 \square P_{13}$ induced by vertex set

$$\{u_7, v_7, u_8, v_8, \dots, u_{13}, v_{13}\}.$$

For the upper bound, we propose strategy for Dominator.

Let $D_1 = v_7$. It suffices to consider Staller's response on $V_L \cup \{u_7\}$, by the symmetric property of $P_2 \square P_{13}$. We carry out case distinctions according to the choice of S_1 as follows:

- Case 1: $S_1 = u_7$. In this case, we see an MBD graph \mathfrak{X}_6 on both subgraphs L and R . Let Dominator carry out the next move on the right MBD graph \mathfrak{X}_6 . Then, whenever Staller claims some vertex on R , let Dominator adopt his strategy on the right MBD graph \mathfrak{X}_6 ; whenever Staller claims some vertex on L , let Dominator respond on the left MBD graph \mathfrak{X}_6 , using the pairing strategy. By Theorem 3.5, Dominator needs 4 steps to dominate all vertices on the right. He needs 6 steps to dominate all vertices on the left. In total, he can win within $6 + 1 + 4 = 11$ steps.
- Case 2: $S_1 = u_5$, then let $D_2 = u_9$. Then let Dominator respond on L if Staller claims any vertex of L , and let him respond on R if Staller claims any vertex of R . On the right is the MBD graph \mathfrak{W}_4 , while on the left it is the situation described in Claim 3.8. By Proposition 3.4, Dominator needs $\gamma'_{MB}(\mathfrak{W}_4) = 3$ steps to dominate all vertices of R . Claim 3.8 tells us that Dominator needs within 6 steps to dominate all vertices of L . Hence, he needs in total within $6 + 3 + 2 = 11$ steps to win the game.
- Case 3: $S_1 \in \{u_3, v_3, u_4, v_4, v_5, u_6, v_6\}$, then let $D_2 = u_5$. Then we have the MBD graph \mathfrak{W}_6 on the right. Let Dominator respond on $V_R \cup \{u_7\}$ whenever Staller claims any vertex in $V_R \cup \{u_7\}$, and let him respond on V_L otherwise. If $S_1 \in \{u_6, v_6\}$, on the left would be \mathfrak{W}_4 . In this case, Dominator needs at most

$$\gamma'_{MB}(\mathfrak{W}_4) + \gamma'_{MB}(\mathfrak{W}_6) + 2 = 3 + 5 + 2 = 10$$

steps to win. Otherwise, $S_1 \in \{u_3, v_3, u_4, v_4, v_5\}$, then the graph on the left fulfills the description in Claim 3.7, where he needs within 4 steps to win. In total, he needs at most $4 + 5 + 2 = 11$ steps to win.

- Case 4: $S_1 \in \{u_2, v_2\}$, then let $D_2 = u_3$. Then let Dominator respond on $V_R \cup \{u_7\}$ whenever Staller claims any vertex of this set. Let Dominator adopt the pairing strategy on the left, where the pairing sets of vertices are $\{u_1, v_1\}$, $\{v_4, v_5\}$, $\{u_5, u_6\}$. In addition, he needs at most one more step to dominate v_2 . By Proposition 3.4, $\gamma'_{MB}(\mathfrak{W}_6) = 5$, hence Dominator needs in total at most $3 + 1 + 5 + 2 = 11$ steps to win the game.
- Case 5: $S_1 \in \{u_1, v_1\}$, then let $D_2 = v_2$. Let Dominator adopt strategy for the S-game on \mathfrak{W}_6 on the right, and the pairing strategy with pairing sets $\{u_3, u_4\}$, $\{v_4, v_5\}$, $\{u_5, u_6\}$. In addition, he needs at most one more step to dominate u_1 . In total he needs within $3 + 1 + 5 + 2 = 11$ steps to win the game.

To sum up, $\gamma_{MB}(P_2 \square P_{13}) \leq 11$. Combining it with the lower bound, we obtain that

$$\gamma_{MB}(P_2 \square P_{13}) = 11. \quad \square$$

Proof of item 3 in Theorem 1.5:

By Theorem 3.9, $\gamma_{MB}(P_2 \square P_{13}) = 11$. When $n > 13$, consider the graph as the two subgraphs $A := P_2 \square P_{13}$ and $B \cong P_2 \square P_{n-13}$ connected by edges $\{u_{13}, u_{14}\}$, $\{v_{13}, v_{14}\}$, up to isomorphism. Let

Dominator respond on A whenever Staller claims a vertex of A , and let Dominator respond on B with the pairing strategy whenever Staller claims a vertex of B . In this way, we see that he needs within

$$11 + (n - 13) = n - 2$$

steps in total, in order to win. Therefore,

$$\gamma_{MB}(P_2 \square P_{13}) \leq n - 2, \quad n \geq 13.$$

For the lower bound, by the Continuation Principle (Remark 1.1) we get

$$\gamma_{MB}(P_2 \square P_{13}) \geq \gamma_{MB}(\mathfrak{X}_n),$$

since $P_2 \square P_{13}$ has one more undominated vertex than the MBD graph \mathfrak{X}_n . By Theorem 3.5,

$$\gamma_{MB}(\mathfrak{X}_n) \geq n - 2, \quad n \geq 6.$$

Hence

$$\gamma_{MB}(P_2 \square P_{13}) \geq \gamma_{MB}(\mathfrak{X}_n) \geq n - 2, \quad n \geq 13.$$

To conclude,

$$\gamma_{MB}(P_2 \square P_{13}) = n - 2, \quad n \geq 13. \quad \square$$

3.4 MBD game on $P_2 \square P_{12}$

In this section, we prove the first two items of Theorem 1.5. We re-state them here as one proposition (Proposition 3.10) and one theorem (Theorem 3.11), and prove them one by one.

Proposition 3.10. *If $1 \leq n \leq 4$, then $\gamma_{MB}(P_2 \square P_n) = n$.*

Proof: For $n = 1$ and $n = 2$, one can easily do the verification. For $n \in \{3, 4\}$, let Dominator adopt the pairing strategy — we see that $\gamma_{MB}(P_2 \square P_n) \leq n$. For the other direction, in the $n = 3$ case, because of the symmetry of the graph, it is sufficient to consider two cases, namely when $D_1 \in \{u_1, u_2\}$.

- Case 1: $D_1 = u_1$. Let $S_1 = v_3$. There are still three undominated vertices, it is not hard to see that Dominator needs at least two more moves to win the game.
- Case 2: $D_1 = u_2$. Let $S_1 = v_2$. One observes that for the remaining two undominated vertices, Dominator needs two more moves to win.

Therefore, $\gamma_{MB}(P_2 \square P_3) = 3$.

When $n = 4$, because of the symmetry of the graph, it is sufficient to consider two cases, namely when $D_1 \in \{u_1, u_2\}$.

- Case 1: $D_1 = u_1$. We see that the remaining graph is the MBD graph \mathfrak{W}_3 on the right. By Proposition 3.4, $\gamma'_{MB}(\mathfrak{W}_3) = 3$. Hence Dominator needs in total four steps to win.
- Case 2: $D_1 = u_2$, let $S_1 = v_4$. One can verify that Dominator cannot dominate the remaining four vertices within two steps.

Thence, $\gamma_{MB}(P_2 \square P_4) = 4$. □

Theorem 3.11. *If $5 \leq n \leq 12$, then $\gamma_{MB}(P_2 \square P_n) = n - 1$.*

In order to prove the above result, we need some preparations.

Lemma 3.12. $\gamma_{MB}(P_2 \square P_5) = 4$.

Proof: First, we prove $\gamma_{MB}(P_2 \square P_5) \geq 4$. It is sufficient to consider three cases because of the symmetry of the graph, namely when $D_1 \in \{u_1, u_2, u_3\}$.

- Case 1: $D_1 = u_1$. The remaining graph is the MBD graph \mathfrak{W}_4 . By Proposition 3.4, $\gamma'_{MB}(\mathfrak{W}_4) = 3$. Hence Dominator needs in total four steps to win.
- Case 2: $D_1 = u_2$. Let $S_1 = v_2$. We see that the remaining part on the right is the MBD graph \mathfrak{X}_3 . By Theorem 3.5, we have $\gamma_{MB}(\mathfrak{X}_3) = 2$. Note that Dominator also needs one step on the left part of the graph, so as to dominate v_1 . He needs in total four steps to win.
- Case 3: $D_1 = u_3$. Let $S_1 = v_5$. Then we see that Dominator needs at least two more steps to dominate v_4, v_5, u_5 . Also he needs at least one more step to dominate u_1, v_1, v_2 . Hence he needs in total no less than four steps.

Now we prove the other direction. Let $D_1 = u_1$. The remaining graph is the MBD graph \mathfrak{W}_4 . Hence no matter how Staller plays, Dominator can win within 3 steps, since $\gamma'_{MB}(\mathfrak{W}_4) = 3$ by Proposition 3.4. □

The following several claims are listed here to prepare for the proof of Theorem 3.18.

Claim 3.13. *In the D -game on $P_2 \square P_{12}$, if we are in the situation where $D_1 = v_6$, $S_1 = u_8$, $D_2 = u_6$, $S_2 = u_9$, if $D_3 \neq v_9$, then Dominator would not win the game.*

Proof of Claim 3.13: The analysis of the situation where $D_3 \neq v_9$ contains three cases.

- Case 1: $D_3 \notin \{u_7, v_7, v_9, u_{10}, v_{10}\}$. Then let $S_3 = v_9$ which forces $D_4 = u_{10}$. Next, let $S_4 = v_8$. We see that Dominator is forced to claim three vertices v_{10}, v_7, u_7 all in the next move, which is not possible.
- Case 2: $D_3 = u_{10}$ (or $D_3 = v_{10}$). Then let $S_3 = v_8$ which forces $D_4 = u_7$. Then let $S_4 = v_9$. We see that Dominator must claim two vertices v_7, v_{10} (or v_7, u_{10}) in the next round, which is obviously not allowed by the rule of this game.
- Case 3: $D_3 = u_7$ (or $D_3 = v_7$). Then let $S_3 = v_9$ which forces $D_4 = u_{10}$. Then let $S_4 = v_8$. We see that Dominator must claim two vertices v_7, v_{10} (or u_7, v_{10}) in the next round, which is obviously not allowed by the rule of this game. □

Claim 3.14. *In the D -game on $P_2 \square P_{12}$, suppose that we are in the situation where $D_1 = v_6$, $S_1 = u_8$, $D_2 = u_7$, $S_2 = v_9$. Then if $D_3 \notin \{u_9, u_{10}, v_{10}\}$, he would not be able to win the game.*

Proof: Suppose that $D_3 \notin \{u_9, u_{10}, v_{10}\}$. We do the analysis in two cases.

- Case 1: $D_3 = v_8$. Let Staller make a sequence of triangle traps which means: $S_3 = u_{10} \rightarrow D_4 = u_9$, $S_4 = v_{11} \rightarrow D_5 = v_{10}$. Next let $S_5 = u_{12}$, Dominator could not anymore dominate both u_{11} and v_{12} .
- Case 2: $D_3 \notin \{v_8, u_9, u_{10}, v_{10}\}$. Let $S_3 = u_9$ which forces $D_4 = u_{10}$. Next, let $S_4 = v_8$. Dominator could not anymore dominate both v_8 and v_9 . \square

Claim 3.15. *In the D-game on $P_2 \square P_{12}$, suppose that we are in the situation where $D_1 = v_6$, $S_1 = u_8$, $D_2 = v_7$, $S_2 = u_9$. Then if $D_3 \notin \{v_9, u_{10}, v_{10}\}$, he would not win the game.*

Proof: Suppose that $D_3 \notin \{u_9, u_{10}, v_{10}\}$. We do the analysis in two cases.

- Case 1: $D_3 = v_8$. Let Staller make a sequence of line traps which means: $S_3 = u_{10} \rightarrow D_4 = v_9$, $S_4 = u_{11} \rightarrow D_5 = v_{10}$. Next let $S_5 = u_{12}$, Dominator would not manage to dominate both v_{11} and v_{12} anymore.
- Case 2: $D_3 \notin \{v_8, u_9, u_{10}, v_{10}\}$. Let $S_3 = v_8$ which forces $D_4 = u_7$. Next, let $S_4 = v_9$. Dominator would not be able to dominate both v_9 and u_9 anymore. \square

Claim 3.16. *In the D-game on $P_2 \square P_{12}$, suppose that we are in the situation where $D_1 = v_6$, $S_1 = u_8$, $D_2 \in \{u_i, v_i \mid 10 \leq i \leq 12\}$, $S_2 = v_8$. Then if $D_3 \notin \{u_7, v_7, u_9, v_9\}$, Dominator would not be able to win the game.*

Proof: Suppose that $D_3 \notin \{u_7, v_7, u_9, v_9\}$. We do the analysis in two cases.

- Case 1: $D_3 = u_6$. Then

$$S_3 = v_9 \rightarrow D_4 = v_7, S_4 = u_9 \rightarrow D_5 = u_7.$$

Now at least one of u_{10} and v_{10} is still free. When u_{10} (or v_{10}) is free. Then let $S_5 = u_{10}$ (or $S_5 = v_{10}$), we see that u_9 (or v_9) would be isolated.

- Case 2: $D_3 \neq u_6$. Then $S_3 = u_7 \rightarrow D_4 = u_9$. Next let $S_4 = v_7$. Then at least one of u_7 and v_8 would not anymore be dominated in the game. \square

Claim 3.17. *In the D-game on $P_2 \square P_{12}$, suppose that we are in the situation where $D_1 = v_3$, $S_1 = u_5$, $D_2 = u_6$, $S_2 = v_8$, $D_3 = v_6$, $S_3 = u_8$. Then if $D_4 \notin \{u_9, v_9\}$, Dominator would not win the game.*

Proof: The proof is left as an exercise. \square

The following result, as the last preparation needed for the proof of Theorem 3.11, has a rather lengthy proof. In order to assist our readers to attain a better understanding, we provide the strategy tree (see Figure 4) for the analysis of Case 1 of the proof of Theorem 3.18.

Theorem 3.18. $\gamma_{MB}(P_2 \square P_{12}) \geq 11$.

Proof: Because of the symmetry of the graph, it is sufficient to consider the cases when

$$D_1 \in \{v_1, v_2, v_3, v_4, v_5, v_6\}.$$

- Case 1: $D_1 = v_6$. Let $S_1 = u_8$.

Case 1.1: $D_2 \in \{u_i, v_i \mid 1 \leq i \leq 5\}$. In this case, let $S_2 = u_8$. We see that on the right part of the graph is the MBD graph \mathfrak{R}_6 . We let Staller plays on this part until Dominator dominates all vertices on this \mathfrak{R}_6 . By Proposition 3.1, we know that he needs 6 steps to accomplish this task and he shall not skip any moves in the middle of this process. After Dominator's last move on this \mathfrak{R}_6 , let Staller continue her response towards D_1 on the left subgraph which is the MBD graph \mathfrak{X}_5 . To view it in another way, it is an MBD game on \mathfrak{X}_5 . By Theorem 3.5, $\gamma_{MB}(M_5) = 4$. So, in total, Dominator needs $1 + 6 + 5 = 12$ moves.

Case 1.2: $D_2 = u_6$. Then let $S_2 = u_9$. In this case, Dominator is forced to claim v_9 — for the reasoning see Claim 3.13. Then let $S_3 = u_7$ which forces $D_4 = v_8$. Next, we let $S_4 = u_{11}$. On the left is the MBD graph \mathfrak{R}_3 , on the right is the MBD graph \mathfrak{Z}_5 . We let Staller play on \mathfrak{R}_3 until the game finishes on this part, and then let her starts playing on \mathfrak{Z}_5 . By Proposition 3.1 and 3.3, $\gamma_{MB}(\mathfrak{R}_3) = 3$ and Dominator would not skip any moves, $\gamma'_{MB}(\mathfrak{Z}_5) = 4$. Hence, Dominator needs in total $4 + 3 + 4 = 11$ moves.

Case 1.3: $D_2 = u_7$. Let $S_2 = v_9$. Now by Claim 3.14, we know that $D_3 \in \{u_9, u_{10}, v_{10}\}$.

Case 1.3.a: $D_3 = u_9$. Let $S_3 = v_8$ which forces $D_4 = v_7$. Next, let $S_4 = u_{11}$. We see \mathfrak{W}_5 on the left and \mathfrak{R}_3 on the right. Let Staller plays on \mathfrak{R}_3 first. Since $\gamma'_{MB}(\mathfrak{W}_5) = 4$ and $\gamma_{MB}(\mathfrak{R}_3) = 3$, Dominator needs $4 + 3 + 4 = 11$ moves to complete the game.

Case 1.3.b: $D_3 = u_{10}$. Let $S_3 = v_8$ which forces $D_4 = v_7$. Then, let $S_4 = v_{10}$ which forces $D_5 = u_9$. Next, let $S_5 = v_{12}$. We see that the remaining part on the right is the MBD graph \mathfrak{R}_2 , while that on the left is the MBD graph \mathfrak{W}_5 . We let Staller first play on \mathfrak{R}_2 until the game finishes on this part, and then let her starts playing on \mathfrak{W}_5 . By Proposition 3.1 and 3.4, we know that $\gamma_{MB}(\mathfrak{R}_2) = 2$ and Dominator shall not skip any moves during the process, $\gamma'_{MB}(\mathfrak{W}_5) = 4$. Hence, Dominator needs in total $5 + 2 + 4 = 11$ moves to accomplish the victory.

Case 1.3.c: $D_3 = v_{10}$. Let $S_3 = v_8$ which forces $D_4 = v_7$. Then, let $S_4 = u_{10}$ which forces $D_5 = u_9$. Next, let $S_5 = u_{12}$. We see that the remaining part on the right is the MBD graph \mathfrak{R}_2 , while that on the left is the MBD graph \mathfrak{W}_5 . We let Staller first play on \mathfrak{R}_2 until the game finishes on this part, and then let her starts playing on \mathfrak{W}_5 . By Proposition 3.1 and 3.4, we know that $\gamma_{MB}(\mathfrak{R}_2) = 2$ and Dominator shall not skip any moves during the process, $\gamma'_{MB}(\mathfrak{W}_5) = 4$. Hence, Dominator needs in total $5 + 2 + 4 = 11$ moves to accomplish the victory.

Case 1.4: $D_2 = v_7$. Let $S_2 = u_9$. This case is analogous to Case 1.3. For the completeness of the contents, we still illustrate the detailed argument. By Claim 3.15 we know that $D_3 \in \{v_9, u_{10}, v_{10}\}$.

Case 1.4.a: $D_3 = v_9$. Let $S_3 = v_8$ which forces $D_4 = u_7$. Then let $S_4 = u_{11}$ which then creates an \mathfrak{R}_3 on the right and let Staller play on this \mathfrak{R}_3 first and then play on the leftmost \mathfrak{W}_5 . Dominator needs $4 + 3 + 4 = 11$ moves to win.

Case 1.4.b: $D_3 = u_{10}$. Let $S_3 = v_8$ which forces $D_4 = u_7$. Then let $S_4 = v_{10}$ which forces $D_5 = v_9$. Then let $S_5 = v_{12}$. Analogous to the cases before, it is \mathfrak{R}_2 on the right and \mathfrak{W}_5 on the left. Let Staller play on \mathfrak{R}_2 first and then on \mathfrak{W}_5 . Dominator needs $5 + 2 + 4 = 11$ moves to win.

Case 1.4.c: $D_3 = v_{10}$. Adopt the following strategy, where \rightarrow refers to “forces” (and this

convention applies also to the later context of this paper):

$$S_3 = u_{10} \rightarrow D_4 = v_9, S_4 = u_7 \rightarrow D_5 = v_8.$$

Let $S_5 = u_{12}$ and let Staller play on the rightmost \mathfrak{R}_2 until the game is locally finished then play on the rightmost \mathfrak{W}_5 . Dominator needs $5 + 2 + 4 = 11$ moves to win.

Case 1.5: $D_2 = v_8$. Let $S_2 = u_{10}$. It is an \mathfrak{R}_4 on the rightmost. By Proposition 3.1, $\gamma_{MB}(\mathfrak{R}_4) = 4$ and Dominator shall not skip any moves in the process. After the game is finished on this \mathfrak{R}_4 , let Staller then claim u_6 . We get an \mathfrak{X}_5 on the leftmost. Since Dominator needs one more step to dominate u_7 and $\gamma_{MB}(\mathfrak{X}_5) = 4$ (by Theorem 3.5), Dominator needs in total $2 + 4 + 4 + 1 = 11$ steps to win.

Case 1.6: $D_2 = u_9$. Let $S_2 = v_7$. Consider the following cases for Dominator's third move.

Case 1.6.a: $D_3 \in \{u_i, v_i \mid 1 \leq i \leq 5\}$. Then

$$S_3 = u_6 \rightarrow D_4 = u_7, S_4 = v_9 \rightarrow D_5 = v_8.$$

Then let $S_5 = v_{11}$ which creates an \mathfrak{R}_3 on the rightmost. Let Staller plays on this \mathfrak{R}_3 first and then starts playing on the left \mathfrak{X}_5 — when forgetting the move D_3 , it is an \mathfrak{X}_5 . Note that D_3 may not be the optimal move, hence Dominator needs at least $\gamma_{MB}(\mathfrak{X}_5) = 4$ (by Theorem 3.5) to finish the local game on the leftmost. And on the leftmost \mathfrak{R}_3 , he needs 3 moves and shall not skip any moves by Proposition 3.1. Hence he needs in total $4 + 3 + 4 = 11$ moves to win the whole game.

Case 1.6.b: $D_3 = u_6$. Let $S_3 = v_9 \rightarrow D_4 = v_8$. Then let $S_4 = v_{11}$ which then creates an \mathfrak{R}_3 on the rightmost, while on the leftmost is an \mathfrak{Z}_5 . Similarly as always, let us just omit this repetitive narrative from now on. Dominator needs in total $4 + 3 + 4 = 11$ steps to win (since $\gamma'_{MB}(\mathfrak{Z}_5) = 4$).

Case 1.6.c: $D_3 = u_7$. Let $S_3 = v_9 \rightarrow D_4 = v_8$. Then $S_4 = v_{11}$. Rightmost: \mathfrak{R}_3 ; leftmost: \mathfrak{W}_5 . Steps to win: $4 + 3 + 4 = 11$.

Case 1.6.d: $D_3 = v_8$. Then $S_3 = u_6 \rightarrow D_4 = u_7$, then $S_4 = u_4$. Leftmost: \mathfrak{R}_5 ; rightmost: \mathfrak{W}_3 . Steps to win: $4 + 5 + 3 = 12$.

Case 1.6.e: $D_3 = v_9$. Then $S_3 = u_6 \rightarrow D_4 = u_7$, then $S_4 = u_4$. Leftmost: \mathfrak{R}_5 ; rightmost: \mathfrak{Z}_3 . Steps to win: $4 + 5 + 2 = 11$.

Case 1.6.f: $D_3 \in \{u_i, v_i \mid 10 \leq i \leq 12\}$. Then

$$S_3 = u_6 \rightarrow D_4 = u_7, S_4 = v_9 \rightarrow D_5 = v_8.$$

Then let $S_5 = u_4$ which creates an \mathfrak{R}_5 on the leftmost. Let Staller plays on this \mathfrak{R}_5 first and then starts playing on the rightmost \mathfrak{X}_3 — when forgetting the move D_3 , it is an \mathfrak{X}_3 . Note that D_3 may not be the optimal move, hence Dominator needs at least $\gamma_{MB}(\mathfrak{X}_3) = 2$ (by Theorem 3.5) to finish the local game on the leftmost. And on the leftmost \mathfrak{R}_5 , he needs 5 moves and shall not skip any moves by Proposition 3.1. Hence he needs in total $4 + 5 + 2 = 11$ moves to win the whole game.

Case 1.7: $D_2 = v_9$. Let $S_2 = u_7$. The proof is analogous to Case 1.6. We omit it here.

Case 1.8: $D_2 \in \{u_i, v_i \mid 10 \leq i \leq 12\}$. Let $S_2 = v_8$. Now, by Claim 3.16,

$$D_3 \in \{u_7, v_7, u_9, v_9\}.$$

Case 1.8.a: $D_3 = u_7$. Then $S_3 = v_9 \rightarrow D_4 = v_7$. Now we look back on the choice of D_2 : If $D_2 \notin \{u_{10}, v_{10}\}$, then he cannot win — because otherwise we let $S_4 = u_9$ which forces Dominator in the next step to claim both u_{10} and v_{10} (in order to not isolate u_9 and v_9) which is not allowed by the rules of this game. Therefore we only need to consider the situation where Dominator actually has an opportunity to win.

- Situation 1: $D_2 = u_{10}$. Then $S_4 = v_{10} \rightarrow D_5 = u_9$. Next let $S_5 = v_{12}$. Now it is back to our familiar pattern. Leftmost: \mathfrak{R}_2 ; rightmost: \mathfrak{W}_5 . Since $\gamma_{MB}(\mathfrak{R}_2) = 2$ and Dominator shall not skip any moves in the game, $\gamma'_{MB}(\mathfrak{W}_5) = 4$, he needs $5 + 2 + 4 = 11$ moves to win.
- Situation 2: $D_2 = v_{10}$. Then $S_4 = u_{10} \rightarrow D_5 = u_9$. Next let $S_5 = u_{12}$. Now it is back to our familiar pattern. Leftmost: \mathfrak{R}_2 ; rightmost: \mathfrak{W}_5 . Since $\gamma_{MB}(\mathfrak{R}_2) = 2$ and Dominator shall not skip any moves in the game, $\gamma'_{MB}(\mathfrak{W}_5) = 4$, he needs $5 + 2 + 4 = 11$ moves to win.

Case 1.8.b: $D_3 = v_7$. The reasoning is analogous to Case 1.8.a — we omit it here.

Case 1.8.c: $D_3 = u_9$. Then

$$S_3 = v_7 \rightarrow D_4 = v_9, S_4 = u_6 \rightarrow D_5 = u_7.$$

Then let $S_5 = u_4$ which creates an \mathfrak{R}_5 on the leftmost, and on the rightmost is an \mathfrak{Z}_3 if we forget the move of D_2 . Since D_2 may not be the optimal choice, Dominator needs at least $\gamma'_{MB}(\mathfrak{Z}_3) = 2$ steps to finish the game on the rightmost subgraph \mathfrak{Z}_3 . In total, at least $4 + 5 + 2 = 11$ moves are needed to win.

Case 1.8.d: $D_3 = v_9$. Then

$$S_3 = u_7 \rightarrow D_4 = u_9, S_4 = u_6 \rightarrow D_5 = v_7.$$

Next let $S_5 = u_4$. The rest of the reasoning is the same as in Case 1.8.c.

- Case 2: $D_1 = v_5$. In this case, the analysis is analogous to in Case 1 — the only change is that when Dominator (or Staller) claims vertex u_i (or v_i) in Case 1, they instead now in this case claims vertex u_{i-1} (or v_{i-1}), where $i \geq 1$. For example, we already see that then we have $S_1 = u_7$ (it was u_8 in Case 1). Note that this is feasible since all the vertices claimed by the two players in Case 1 has subscript bigger than 3. We have verified all the cases, and we omit the detailed reasoning (or rather repetition) here.
- Case 3: $D_1 = v_4$. In this case, the analysis is analogous to in Case 1 — the only change is that when Dominator (or Staller) claims vertex u_i (or v_i) in Case 1, they instead now in this case claims vertex u_{i-2} (or v_{i-2}), where $i \geq 1$. For example, we already see that then we have $S_1 = u_6$ (it was u_8 in Case 1). Note that this is feasible since all the vertices claimed by the two players in Case 1 has subscript bigger than 3. We have verified all the cases, and we omit the detailed reasoning here.

- Case 4: $D_1 = v_3$. Let $S_1 = u_5$. In this case, most of the the analysis is analogous to in Case 1 — the only change is that when Dominator (or Staller) claims vertex u_i (or v_i) in Case 1, they instead now in this case claims vertex u_{i-3} (or v_{i-3}), where $i \geq 1$. Note that this is feasible since all the vertices claimed by the two players in Case 1 has subscript bigger than 3. But there are two cases we need to consider separately, namely when $D_2 \in \{u_6, v_6\}$ — because, if in these two cases we let Staller adopt the analogous strategy as in Case 1, Dominator would be able to win within 10 moves. We have verified all other cases. Now, let us talk about these two cases.

Case 4.1: $D_2 = u_6$. Let $S_2 = v_8$. Now, Staller’s strategy depends on Dominator’s third move.

Case 4.1.1: $D_3 \in \{u_i, v_i \mid 1 \leq i \leq 5\}$. Let $S_3 = v_6$ which then creates an \mathfrak{R}_6 on the rightmost. Let Staller play on this \mathfrak{R}_6 now. By Proposition 3.1, Dominator needs 6 steps to win and he shall not skip any moves during the process. Next, Staller’s strategy depends on what was D_3 :

- Case 4.1.1.a: $D_3 \in \{u_1, v_1, u_2, v_2\}$. Let $S_9 = v_4$ which forces $D_{10} = v_5$. Then $S_{10} = u_3 \rightarrow D_{11} = u_4$. Dominator needs at least $4 + 6 + 1 = 11$ steps to win.
- Case 4.1.1.b: $D_3 = u_3$. Then $S_9 = v_4 \rightarrow D_{10} = v_5$. One more step is needed on the leftmost to dominate u_1 and v_1 . He needs in total $4 + 6 + 1 = 11$ moves to win.
- Case 4.1.1.c: $D_3 = u_4$. Then $S_9 = v_5 \rightarrow D_{10} = v_4$. On the left is an \mathfrak{M}_2 . Steps to win: $4 + 6 + 2 = 12$.
- Case 4.1.1.d: $D_3 = v_4$. Then $S_9 = u_1$ which forces Dominator to spend two more steps to finish the game. Steps to win: $3 + 6 + 2 = 11$.
- Case 4.1.1.e: $D_3 = v_5$. Then it is not hard to see that he needs in the best case two steps to win, so he needs at least $3 + 6 + 2 = 11$ steps to finish the game.

Case 4.1.2: $D_3 = v_9$. Let $S_3 = u_8$. By Claim 3.17, $D_4 \in \{u_9, v_9\}$. We leave the rest of the argument for readers as an exercise.

From now on, we only list the cases, since we believe after all the case distinctions above, our readers already got the taste of how this can be done. We have verified all cases, but we omit the detailed long-winded here.

Case 4.1.3: $D_3 = u_7$. Let $S_3 = v_9$. Then it is forced that $D_4 \in \{u_9, u_{10}, v_{10}\}$.

Case 4.1.4: $D_3 = v_7$. Let $S_3 = u_9$. Then it is forced that $D_4 \in \{u_9, u_{10}, v_{10}\}$.

Case 4.1.5: $D_3 = u_8$. Let $S_3 = v_{10}$.

Case 4.1.6: $D_3 = u_9$. Let $S_3 = v_7$.

Case 4.1.7: $D_3 = v_9$. Then $S_3 = u_7$.

Case 4.1.8: $D_3 \in \{u_{10}, u_{11}, u_{12}, v_{10}, v_{11}, v_{12}\}$. Let $S_3 = u_8$. Then it has to be that $D_4 \in \{u_9, v_9\}$.

Case 4.2: $D_2 = v_6$. The reasoning is analogous to Case 4.1.

- Case 5: $D_1 = v_2$. Then $S_1 = u_4$. In this case, the analysis is analogous to in Case 4 — the only change is that when Dominator (or Staller) claims vertex u_i (or v_i) in Case 4, they instead now in this case claims vertex u_{i-1} (or v_{i-1}), where $i \geq 2$. Note that in the case when the subscript of the claimed vertex is 1, we keep it unchanged. We have verified all the cases, and we omit the detailed reasoning (or rather repetition) here.

- Case 6: $D_1 = v_1$. The rest is the MBD graph \mathfrak{W}_{11} . By Proposition 3.4, $\gamma'_{MB}(\mathfrak{W}_{11}) = 10$. So Dominator needs in total $1 + 10 = 11$ steps to win. \square

Proof of Theorem 3.11: Let $n \geq 5$, we consider the D-game on $P_2 \square P_n$. Consider the graph as two subgraphs joined by two edges: $P_2 \square P_5$ on the leftmost and the rightmost is isomorphic to $P_2 \square P_{n-5}$; obviously they are joined by two edges $\{u_5, u_6\}$, and $\{v_5, v_6\}$. Let Dominator start the game on the leftmost $P_2 \square P_5$. Then, whenever Staller claims a vertex on this part, he responds using his strategy for $P_2 \square P_5$. Otherwise, let him adopt the pairing strategy to respond to Staller's move. By Lemma 3.12, $\gamma_{MB}(P_2 \square P_5) = 4$. Dominator can win the game within $4 + (n - 5) = n - 1$ steps. Therefore,

$$\gamma_{MB}(P_2 \square P_n) \leq n - 1 \text{ when } n \geq 5.$$

Let $5 < m < 12$. Suppose that $\gamma_{MB}(P_2 \square P_m) < m - 1$. Then consider the MBD game on $P_2 \square P_{12}$. We consider the graph as two subgraphs joined by two edges: $P_2 \square P_m$ on the leftmost and the rightmost is isomorphic to $P_2 \square P_{12-m}$. Let Dominator starts playing on the leftmost subgraph. Then, whenever Staller claims a vertex on this part, he responds using his strategy for $P_2 \square P_m$. Otherwise, let him adopt the pairing strategy to respond to Staller's move. We see that Dominator will win the game with less than $(m - 1) + (12 - m) = 11$ steps, which contradicts the result of Theorem 3.18. Therefore,

$$\gamma_{MB}(P_2 \square P_n) \geq n - 1, \text{ for } 5 \leq n \leq 12.$$

In the end, we get that $\gamma_{MB}(P_2 \square P_n) = n - 1$, for $5 \leq n \leq 12$. \square

3.5 Union of grids

In this subsection, we prove the two results on the disjoint union of $P_2 \square P_n$ s, using Theorem 1.4 and Theorem 1.5.

Proof of Theorem 1.6: By the pairing strategy where the pairing sets are $\{u_i, v_i\}$ s in each copy, Dominator can win within $k \cdot n$ steps. However, by Theorem 1.4, we know that no matter in which copy Staller claims a vertex, Dominator has to respond on the same copy, otherwise he would lose the game. Hence he needs at least $k \cdot \gamma'_{MB}(P_2 \square P_n) = k \cdot n$ steps to win. So we get

$$\gamma'_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = k \cdot n, \text{ for any } k, n \geq 1. \quad \square$$

Proof of Theorem 1.7: First we prove the upper bounds. Let Dominator adopt the following strategy: let him respond on the same copy where Staller played her move in the last round. By Theorem 1.4, we know that Dominator does not skip any moves in a copy where Staller starts the game — namely, he has to respond on the same copy whenever Staller plays a move on that copy — and needs n steps to dominate all vertices of this copy. For the copy where he starts the game, the situation is just a D-game on $P_2 \square P_n$. Hence we get that

$$\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) \leq (k - 1) \cdot n + \gamma_{MB}(P_2 \square P_n).$$

Now we prove the lower bounds. After Dominator's first move, let Staller continue the game on any other copy than the one where Dominator played his first move, and let Staller continue play on this copy using the strategy for $P_2 \square P_n$ until the game is finished on this copy. By Theorem 1.4, we know that

Dominator does not skip any moves during this process and needs n steps to dominate all vertices of this copy. Then, let Staller repeat this strategy on all the remaining copies except for the one where Dominator starts the game. After this process, let Staller continue by responding on that remaining copy, then it should be considered as a D-game on $P_2 \square P_n$. Hence we get that

$$\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) \geq (k-1) \cdot n + \gamma_{MB}(P_2 \square P_n).$$

Combining the above results, we get that

$$\gamma_{MB}(\dot{\cup}_{i=1}^k (P_2 \square P_n)_i) = (k-1) \cdot n + \gamma_{MB}(P_2 \square P_n).$$

By Theorem 1.5, we get the results in the statement of this theorem. \square

4 Concluding remarks

We proved that Dominator needs exactly n moves to win in the S-game on $P_2 \square P_n$ for every $n \geq 1$, while in the D-game he needs exactly n , $n-1$, $n-2$ moves for $1 \leq n \leq 4$, $5 \leq n \leq 12$, and $n \geq 13$, respectively. We showed that for k copies of $P_2 \square P_n$, Dominator needs exactly $k \cdot n$ steps to win in the S-game, and $k \cdot n$, $k \cdot n - 1$, $k \cdot n - 2$ steps to win in the D-game for $1 \leq n$, $5 \leq n \leq 12$, $n \geq 13$, respectively. The exact result for general Cartesian $P_m \square P_n$ does not seem easy, so it would be interesting to consider the situation for $P_3 \square P_n$, as a starting point.

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⁽ⁱ⁾ At the time of the publication of this paper, Dongsheng Wu was affiliated with Yanqi Lake Beijing Institute of Mathematical Sciences and Applications (BIMSA), and with Tsinghua University, Beijing, China.

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