

Partitioning the vertex set of a bipartite graph into complete bipartite subgraphs

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Given a graph and a positive integer k , the biclique vertex-partition problem asks whether the vertex set of the graph can be partitioned into at most k bicliques (connected complete bipartite subgraphs). It is known that this problem is NP-complete for bipartite graphs. In this paper we investigate the computational complexity of this problem in special subclasses of bipartite graphs. We prove that the biclique vertex-partition problem is polynomially solvable for bipartite permutation graphs, bipartite distance-hereditary graphs; and that it remains NP-complete for perfect elimination bipartite graphs and bipartite graphs containing no 4-cycles as induced subgraphs.

Keywords: bicliques, bipartite graph, computational complexity, partitioning problem

1 Introduction

In this paper we consider the computational complexity of the problem of partitioning the vertex set of the given graph into at most k bicliques (BICLIQUE VERTEX-PARTITION) in subclasses of bipartite graphs. The problem arises in information security, data mining, e-commerce, information retrieval, and network management [3, 11]. It is known that the BICLIQUE VERTEX-PARTITION problem is NP-complete for bipartite graphs [8, 11] and polynomially solvable for trees [1]. Recently, Martin et al. [15] showed that this problem is computationally hard even for fixed $k = 2$.

The aim of this paper is to investigate the following question: which subclasses of bipartite graphs does the BICLIQUE VERTEX-PARTITION problem become polynomially solvable for? In Section 3 we establish a relation between the considered problem restricted to (K_3, C_5, C_6) -free graphs and colouration problems of graphs. In Section 4, using this relation, we show that the BICLIQUE VERTEX-PARTITION problem is solvable in polynomial time for bipartite permutation graphs and bipartite distance-hereditary graphs. Finally, we establish the NP-completeness of this problem for perfect elimination bipartite graphs and bipartite C_4 -free graphs.

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2 Preliminaries

We use standard graph-theoretic terminology (see, for example, [2, 4, 7]). In this paper we consider only finite undirected simple graphs $G = (V, E)$ with the vertex set $V = V(G)$ and the edge set $E = E(G)$. A graph with the an empty edge set is referred to as an *empty graph*. We denote the set of all vertices of G adjacent to a vertex $v \in V$ by $N(v)$. For a vertex subset $S \subseteq V$, we denote by $G[S]$ the subgraph of G induced by S . If $v \in V$ and $S \subseteq V \setminus \{v\}$, the notation $v \sim S$ means that every vertex in S is adjacent to v . A vertex subset $I \subseteq V$ of G is *independent* if no two vertices in I are adjacent in G . A vertex subset $D \subseteq V$ of G is a *dominating set* if every vertex in $V \setminus D$ is adjacent to at least one vertex in D . The minimum size among the dominating sets of G is the *domination number* of G and is denoted by $\gamma(G)$. A vertex subset $U \subseteq V$ of G is a *vertex cover* of G if every edge of G is incident with at least one vertex in U . The minimum size of a vertex cover of G is the *vertex covering number* of G and is denoted by $\tau(G)$.

A $u - v$ *path* of G is a path in G connecting vertices u, v of G . The number of edges of a $u - v$ path is its *length*. The *distance* $d_G(u, v)$ between vertices u, v in G is the length of a shortest $u - v$ path. If G does not contain such a path, then we set $d_G(u, v) = \infty$. For a graph $G = (V, E)$ and a positive integer k , the k th *power* of G is the graph G^k with the same vertex set V and two vertices u, v are adjacent in G^k if and only if $d_G(u, v) \leq k$. The *complement* \overline{G} of a graph G is the graph with vertex set $V(G)$ and $\{u, v\} \in E(\overline{G})$ if and only if $\{u, v\} \notin E(G)$.

For a family Γ of graphs, a graph G is Γ -*free* if no induced subgraph of G is isomorphic to a graph in Γ . A graph $G = (V, E)$ is a *bipartite graph* if there is a partition of the vertex set $V = V_1 \cup V_2$ such that V_1, V_2 are independent sets. For such a bipartite graph G we write $G = ((V_1, V_2), E)$. A graph G is a *permutation graph* if there are two permutations P, Q of the vertex set $V(G)$ such that there is an edge between vertices u, v if and only if u precedes v in one of $\{P, Q\}$ and v precedes u in the other permutation [18]. A *bipartite permutation graph* is a permutation graph that is also a bipartite graph. A graph G is *distance-hereditary* if each connected induced subgraph F of G has the property that $d_F(u, v) = d_G(u, v)$ for any two vertices $u, v \in V(F)$ [12]. A graph is a *bipartite distance-hereditary graph* if it is a distance-hereditary graph and a bipartite graph.

A *biclique* of a graph G is a connected complete bipartite (not necessarily induced) subgraph of G . We assume that a subgraph consisting of exactly one vertex is a biclique. A biclique $B = ((U_1, U_2), E)$ is a *star centred at a vertex* v if $U_1 = \{v\}$ or $U_2 = \{v\}$. A set \mathcal{S} of bicliques of G is a *biclique vertex-cover* of G of *size* $|\mathcal{S}|$ if each vertex of G is contained in at least one of the bicliques in \mathcal{S} . A biclique vertex-cover of G consisting of mutually vertex-disjoint bicliques is a *biclique vertex-partition* of G .

Fleischner et al. [8] observed that one can always make the bicliques of a biclique vertex-cover of a graph vertex-disjoint without increasing the size of the cover. Let $\mathcal{S} = \{B_1, B_2, \dots, B_k\}$ be a biclique vertex-cover of a graph G and bicliques B_i, B_j ($i \neq j$) have at least one common vertex. Let $S_i = V(B_i) \setminus V(B_j)$, $S_j = V(B_j) \setminus V(B_i)$ and $S_{ij} = V(B_i) \cap V(B_j)$. If $S_i = \emptyset$ (or $S_j = \emptyset$), then the biclique B_i (respectively B_j) can be eliminated from the cover \mathcal{S} . Thus, we can assume that $S_i \neq \emptyset$ and $S_j \neq \emptyset$. If $B_i - S_{ij}$ is a biclique of G , then replace B_i in the cover by the biclique $B_i - S_{ij}$. Otherwise, there is a vertex $u \in S_{ij}$ such that $u \sim S_i$. If B_j is not a star centred at u , then replace the pair of bicliques B_i, B_j in the cover \mathcal{S} by vertex-disjoint bicliques $B_j - u, ((\{u\}, S_i), E_i)$. In the case when B_j is a star centred at u , we merge the stars B_j and $((\{u\}, S_i), E_i)$ into one a star, which add to the cover \mathcal{S} instead of the bicliques B_i, B_j . It is not hard to see that after replacing the bicliques B_i and B_j by at most two vertex-disjoint bicliques, the cover \mathcal{S} remains biclique vertex-cover of G without increasing of the size. This implies the following remark:

Remark 1 ([8]) For every graph G , there is a biclique vertex-cover of G of size at most k if and only if there is a biclique vertex-partition of G of size at most k .

For a graph G , we denote by $b(G)$ the minimum partition size over all biclique vertex-partitions of G . The decision problem associated with the parameter $b(G)$ is defined as follows:

BICLIQUE VERTEX-PARTITION

Instance: A graph G and a positive integer k .

Question: Does G have a biclique vertex-partition of size at most k ? In other words, is $b(G) \leq k$?

The optimisation version of BICLIQUE VERTEX-PARTITION asks for a biclique vertex-partition of the given graph G with minimum size.

3 Chromatic characterisation of biclique vertex-partitions

In this section we characterise $b(G)$ for a (K_3, C_5, C_6) -free graph G as the chromatic number of a special vertex colouring of the graph that leads to the construction of a biclique vertex-cover of G such that all vertices of the same colour appear in a single biclique of the cover. It allows us to transform BICLIQUE VERTEX-PARTITION for (K_3, C_5, C_6) -free graphs into an equivalent problem involving vertex colouring of graphs.

Theorem 2 Let G be a (K_3, C_5) -free graph. If $S \subseteq V(G)$ is a non-empty vertex subset of G such that the induced subgraph $G[S]$ contains at least one edge of G and $d_G(u, v) \leq 2$ for any two vertices u, v in S , then $G[S]$ is a biclique of G .

Proof: Let G be a (K_3, C_5) -free graph. Assume that $S \subseteq V(G)$ is a non-empty vertex subset of G such that the induced subgraph $G[S]$ contains an edge $e = \{a, b\}$ of G and $d_G(u, v) \leq 2$ for any two vertices u, v in S .

Let us show that $G[S]$ is a connected graph. Suppose not, and let $w \in S$ be a vertex such that $G[S]$ does not contain any $a - w$ paths and $b - w$ paths. Then the vertices a, b are not adjacent to w in G . As $d_G(a, w) \leq 2$ and $d_G(b, w) \leq 2$, the graph G contains a $a - w$ path and a $b - w$ path that have length 2. There are only two possibilities: *i*) the $a - w$ path and the $b - w$ path have two common vertices and *ii*) the $a - w$ path and the $b - w$ path have only one common vertex w . Both possibilities are illustrated in Figure 1. In case *i*) the graph G contains K_3 as an induced subgraph and in case *ii*) G contains either C_5

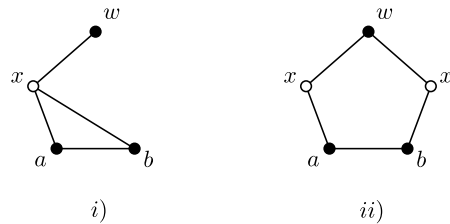


Fig. 1: In the graph G , a $a - w$ path of length 2 and a $b - w$ path of length 2 have either *i*) two common vertices, i.e., w, x or *ii*) only one common vertex w .

(if $\{x, x'\} \notin E(G)$) or K_3 (if $\{x, x'\} \in E(G)$) as an induced subgraph. We obtain a contradiction with

the assumption that the graph G is a (K_3, C_5) -free graph. Therefore, $G[S]$ is indeed a connected graph. In fact, we have shown that each vertex $w \in S \setminus \{a, b\}$ is adjacent to at least one vertex in $\{a, b\}$. As a consequence, the graph $G[S]$ does not contain induced subgraphs isomorphic to \overline{P}_3 .

It is well-known that a connected graph is complete bipartite if and only if no induced subgraph is isomorphic to K_3 or \overline{P}_3 . The connected graph $G[S]$ does not contain induced subgraphs isomorphic to K_3 (as G doesn't have an induced copy of K_3) or \overline{P}_3 (as shown above). Therefore, $G[S]$ is indeed a biclique of G . \square

Theorem 3 *Let G be a (K_3, C_5, C_6) -free graph. If $S \subseteq V(G)$ is a non-empty vertex subset of G such that $d_G(u, v) \leq 2$ for any two vertices u, v in S , then the induced subgraph $G[S]$ is either a biclique of G or an empty graph with the condition that there is a vertex $w \in V(G) \setminus S$ such that $w \sim S$.*

Proof: Let G be a (K_3, C_5, C_6) -free graph and let $S \subseteq V(G)$ be a non-empty vertex subset of G such that $d_G(u, v) \leq 2$ for any two vertices u, v in S .

If $|S| = 1$ then, by definition, $G[S]$ is a biclique of G .

If $|S| \geq 2$ and $G[S]$ contains at least one edge of G then $G[S]$ is a biclique of G . It follows from Theorem 2.

Consider the case when $|S| \geq 2$ and $G[S]$ does not contain any edges of G , i.e., $G[S]$ is an empty graph or, in other words, S is an independent set in G . To show that in this case there is a vertex $w \in V(G) \setminus S$ such that $w \sim S$, we use induction on the size of S .

Let S consist of only two non-adjacent vertices u and v . As $d_G(u, v) \leq 2$, the graph G contains a $u - v$ path of length 2. It implies the existence of a vertex $w \in V(G) \setminus S$ adjacent to both u, v .

Let $k > 2$ be a positive integer number. Assume that for every independent set $S \subseteq V(G)$ of size strictly less than k there exists a vertex $w \in V(G) \setminus S$ such that $w \sim S$ (*induction hypothesis*). Consider an arbitrary independent set $S \subseteq V(G)$ such that $|S| = k$. We denote by v_1, v_2, v_3 three arbitrary vertices in S . By the induction hypothesis, there are vertices w_1, w_2, w_3 in $V(G) \setminus S$ such that $w_i \sim S \setminus \{v_i\}$, $i \in \{1, 2, 3\}$ (as illustrated in Figure 2). We note that the vertices w_1, w_2, w_3 are pairwise non-adjacent

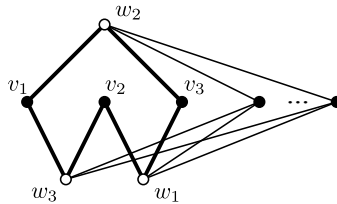


Fig. 2: The set S consists of black vertices. Vertices $w_1, w_2, w_3 \in V(G) \setminus S$ such that $w_i \sim S \setminus \{v_i\}$, $i \in \{1, 2, 3\}$.

(otherwise, the graph G would contain K_3 as an induced subgraph). Since the graph G does not contain C_6 as an induced subgraph, there exists $i \in \{1, 2, 3\}$ such that w_i is adjacent to v_i . Therefore $w_i \sim S$ for some $i \in \{1, 2, 3\}$. It ends the proof of Theorem 3. \square

Let G be a graph. A map $c : V(G) \mapsto \{1, 2, \dots, k\}$ is a *long-distance colouring* of G if $c(u) \neq c(v)$ for any two vertices u, v of G such that $d_G(u, v) > 2$. The smallest positive integer k such that a graph G has a long-distance colouring $c : V(G) \mapsto \{1, 2, \dots, k\}$ is denoted by $\overline{\chi}(G)$. In other words, $\overline{\chi}(G)$ is the chromatic number of a long-distance colouring of G . The following observation is easy to see.

Observation 4 If $B = ((U_1, U_2), E)$ is a biclique of a graph G then $d_G(u, v) \leq 2$ for any two vertices u, v in $U_1 \cup U_2$.

Lemma 5 $\bar{\chi}(G) \leq b(G)$ for every graph G .

Proof: Let $\{B_1, B_2, \dots, B_r\}$ be a biclique vertex-partition of a graph G , where every B_i is a biclique of G . Observation 4 implies that a map $c : V(G) \mapsto \{1, 2, \dots, r\}$ is a long-distance colouring when $c(v)$ is defined as the index i for which B_i contains $v \in V(G)$. \square

Theorem 6 $\bar{\chi}(G) = b(G)$ for every (K_3, C_5, C_6) -free graph G .

Proof: The inequality $\bar{\chi}(G) \leq b(G)$ was proved in Lemma 5. To show the converse inequality, let $c : V(G) \mapsto \{1, 2, \dots, r\}$ be a long-distance colouring of a (K_3, C_5, C_6) -free graph G . Let $V_i = \{v \in V(G) : c(v) = i\}$. Assume without loss of generality that $V_i \neq \emptyset$ for every $i = 1, 2, \dots, r$. By Theorem 3, the induced subgraph $G[V_i]$ is either a biclique of G or an empty graph satisfying the condition that there is a vertex $w_i \in V(G) \setminus V_i$ such that $w_i \sim V_i$.

If $G[V_i]$ is not a biclique of G then the vertices in V_i have a common neighbour w_i and V_i is included in the vertex set of a star $((\{w_i\}, V_i), E_i)$ of G . Thus, for each $i = 1, 2, \dots, r$ we have a biclique of G , say B_i that includes all vertices in V_i . Hence $\{B_1, B_2, \dots, B_r\}$ is a biclique vertex-cover of G of size r . By Remark 1, if G has a biclique vertex-cover of size at most r , then G has a biclique vertex-partition of size at most r . Therefore, $b(G) \leq \bar{\chi}(G)$. \square

As usual, $\chi(G)$ denotes the chromatic number of a graph G . For a graph G , the complement of the square $\overline{G^2}$ has an edge $\{u, v\}$ if and only if $d_G(u, v) > 2$. Therefore, $\bar{\chi}(G) = \chi(\overline{G^2})$. The following statement follows directly from Theorem 6.

Corollary 7 $b(G) = \chi(\overline{G^2})$ for every (K_3, C_5, C_6) -free graph G .

Thus, for every (K_3, C_5, C_6) -free graph G computing $b(G)$ is equivalent to computing $\chi(\overline{G^2})$. It is known that the chromatic number for perfect graphs is computable in polynomial time [10]. Hence, if G is a (K_3, C_5, C_6) -free graph such that $\overline{G^2}$ is perfect, then one can find $b(G)$ in polynomial time, because $\chi(\overline{G^2})$ can be found in polynomial time. It is well-known that a graph is perfect if and only if its complement is perfect [14]. Summarising these ideas, we obtain the following remark.

Remark 8 Let Γ be a graph class. If Γ is a subclass of (K_3, C_5, C_6) -free graphs such that $\overline{G^2}$ is perfect for every $G \in \Gamma$, then BICLIQUE VERTEX-PARTITION is polynomially solvable when restricted to Γ .

4 Complexity results

A hierarchy of commonly studied bipartite graphs is given in Figure 3. Previously, the complexity of BICLIQUE VERTEX-PARTITION was known for forests (polynomial) and general bipartite graphs (NP-complete). In this section, we establish the computational complexity of this problem for bipartite permutation, bipartite distance-hereditary, perfect elimination bipartite, and bipartite C_4 -free graphs.

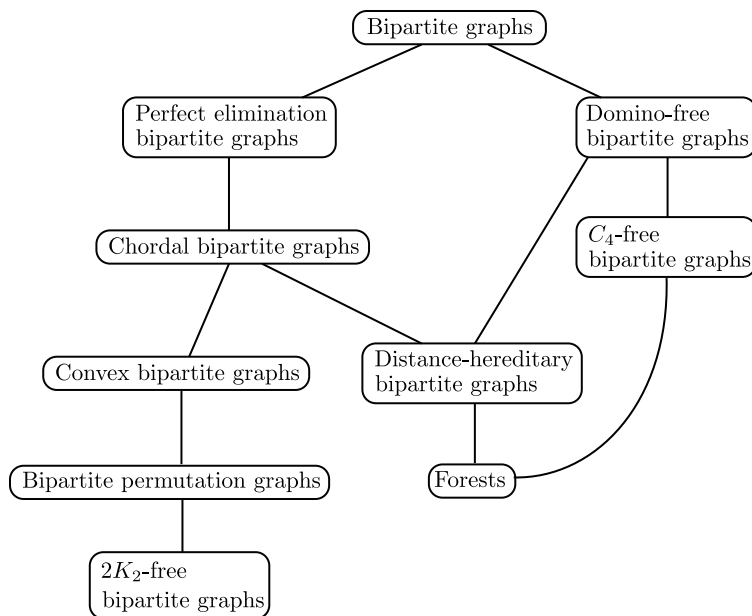


Fig. 3: A hierarchy of bipartite graph classes.

4.1 Polynomial cases

We begin with the polynomial cases. BICLIQUE VERTEX-PARTITION is trivially solvable in linear time when restricted to bipartite P_4 -free graphs. Let G be an arbitrary bipartite P_4 -free graph. As every connected (K_3, P_4) -free graph is complete bipartite, each connected component of G is a biclique of G . Therefore, a biclique vertex-partition of G of the minimum size is the set of its connected components, which can be found in linear time by breadth-first search.

Theorem 9 BICLIQUE VERTEX-PARTITION is polynomially solvable for bipartite permutation graphs and bipartite distance-hereditary graphs.

Proof: A bipartite graph is a *chordal bipartite graph* if each simple cycle of length at least six has a chord, i.e., an edge between two vertices which are not consecutive in the cycle. Thus, every chordal bipartite graph is a (K_3, C_5, C_6) -free graph. As the class of bipartite permutation graphs and the class of bipartite distance-hereditary graphs are subclasses of chordal bipartite graphs, both these graph classes are contained in the class of (K_3, C_5, C_6) -free graphs.

Let us show that $b(G)$ can be found in polynomial time for every bipartite permutation graph G . We begin with a definition. A graph is a *cocomparability graph* if its complement is a comparability graph. Damaschke in [6] showed that the k th power of each cocomparability graph is a cocomparability graph, for all $k \geq 1$. It is known that the class of bipartite permutation graphs is equal to the class of bipartite cocomparability graphs [2, p. 93]. Hence, the square of each bipartite permutation graph is a cocomparability graph. As a consequence of this, G^2 is a comparability graph for every bipartite permutation graph G . As mentioned above, the class of bipartite permutation graphs is a subclass of (K_3, C_5, C_6) -free

graph. By Corollary 7, $b(G) = \chi(\overline{G^2})$ for every bipartite permutation graph G . The chromatic number of comparability graphs can be found in linear time [16]. Therefore, $b(G)$ can be found in polynomial time for every bipartite permutation graph G by finding of the chromatic number of the comparability graph $\overline{G^2}$.

It is known that even powers of distance-hereditary graphs are chordal [2, p. 164]. Thus, if G is a bipartite distance-hereditary graph then G^2 is a chordal graph and, consequently, a perfect graph. We conclude that BICLIQUE VERTEX-PARTITION is polynomially solvable for bipartite distance-hereditary graphs; this follows from Remark 8. \square

We note that the well-studied class of convex bipartite graphs is also a subclass of (K_3, C_5, C_6) -free graphs. In view of Corollary 7, $b(G) = \chi(\overline{G^2})$ for every convex bipartite graph G . However, there is a convex bipartite graph whose square is not perfect (see Figure 4). So, we cannot show the polynomial

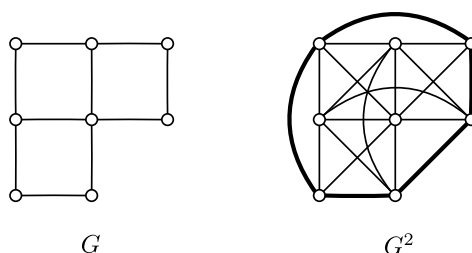


Fig. 4: A convex bipartite graph G whose square is not perfect (G^2 contains C_5 as an induced subgraph).

time solvability of BICLIQUE VERTEX-PARTITION for convex bipartite graphs by using Remark 8.

4.2 NP-completeness

Next, we study the question of whether BICLIQUE VERTEX-PARTITION becomes easier when restricted to bipartite C_4 -free graphs.

Given a dominating set D of a graph G , we can partition the vertex set of G using only stars centred at the vertices of D . Thus, the domination number $\gamma(G)$ of G is an upper bound on $b(G)$. Let G be a graph containing no C_4 as a subgraph. As G has no bicliques other than stars, we have $b(G) = \gamma(G)$. Therefore, computing $b(G)$ is equivalent to computing $\gamma(G)$ for all graphs without C_4 as a subgraph.

Let us observe that every bipartite C_4 -free graph does not contain C_4 as a subgraph. So, we have $b(G) = \gamma(G)$ for every bipartite C_4 -free graph G . Thus, instead of BICLIQUE VERTEX-PARTITION we can consider the decision problem associated with $\gamma(G)$:

DOMINATING SET

Instance: A graph G and a positive integer k .

Question: Does there exist a dominating set of G of size at most k ? In other words, is $\gamma(G) \leq k$?

Lemma 10 For any $t \geq 2$, DOMINATING SET is NP-complete in the class of 2-connected planar $(C_4, C_6, \dots, C_{2t})$ -free bipartite graphs of vertex degree at most 6.

Proof: DOMINATING SET restricted to 2-connected planar $(C_4, C_6, \dots, C_{2t})$ -free bipartite graphs of vertex degree at most 6 is trivially in NP. For the hardness proof, we use a reduction from the VERTEX

COVER problem in the class of 2-connected cubic planar graphs, which is known to be NP-complete [17]. Given a graph and a positive integer k , VERTEX COVER asks whether there is a vertex cover of the graph of size at most k . We closely follow ideas from [5].

Let $t \geq 2$ be an arbitrary integer positive number. Consider an arbitrary *odd* integer positive number $p \geq t$ such that $p \equiv 1 \pmod 3$. Given a 2-connected cubic planar graph $G = (V, E)$ with m edges, form the graph G' by replacing each edge $e = \{u, v\}$ in G by a gadget G_e (a simple cycle with $2p + 4$ vertices) illustrated in Figure 5. It is not hard to see that the obtained graph G' is a 2-connected planar bipartite

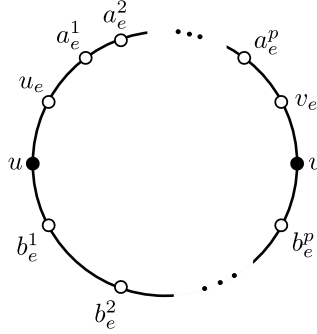


Fig. 5: A gadget G_e .

graph with maximum degree at most 6 and the graph G' does not contain cycles C_4, C_6, \dots, C_{2t} as a subgraphs. We claim that $\gamma(G') = \tau(G) + \frac{m}{3}(2p + 1)$, where $\tau(G)$ is the vertex covering number of the graph G .

Let S be a vertex cover of G . We construct a dominating set D of G' as follows. Initially set $D = S$. We have that at least one vertex of any two vertices u, v adjacent in G is in D . For each edge $e = \{u, v\} \in E$, if $u \in D$ and $v \notin D$ then add to D the following vertices: the vertex v_e , vertices from minimum dominating set of the $a_e^1 - a_e^{p-1}$ path and vertices from minimum dominating set of the $b_e^2 - b_e^p$ path; otherwise $v \in D$ and in this case add to D the following vertices: the vertex u_e , vertices from minimum dominating set of the $a_e^2 - a_e^p$ path and vertices from minimum dominating set of the $b_e^1 - b_e^{p-1}$ path. It is easy to see that D is a dominating set of G' . As $\gamma(P_n) = \lceil \frac{n}{3} \rceil$ for any n [13],

$$|D| = |S| + \sum_{e \in E(G)} \left(1 + \frac{p-1}{3} + \frac{p-1}{3} \right) = |S| + \frac{m}{3}(2p + 1).$$

If S is a vertex cover of G such that $|S| = \tau(G)$ then we obtain that $\gamma(G') \leq \tau(G) + \frac{m}{3}(2p + 1)$.

Given a dominating set D of G' , we construct a vertex cover S of G such that $|S| \leq |D| - \frac{m}{3}(2p + 1)$. First construct a dominating set D' of G' such that $|D'| \leq |D|$ as follows. Initially set $D' = D$. Consider an arbitrary edge $e = \{u, v\}$ of G such that $u \notin D', v \notin D'$ and let $D_e^a = D' \cap \{u_e, v_e, a_e^1, a_e^2, \dots, a_e^p\}$, $D_e^b = D' \cap \{b_e^1, b_e^2, \dots, b_e^p\}$. As $u \notin D', v \notin D'$ and $p \equiv 1 \pmod 3$,

$$|D_e^a| \geq \gamma(P_{p+2}) = \frac{p+2}{3}, \quad |D_e^b| \geq \gamma(P_p) = \left\lceil \frac{p}{3} \right\rceil = \frac{p+2}{3}.$$

In this case, replace $D_e^a \cup D_e^b$ in D' by a set R_e consisting of the vertices u, v_e , vertices from minimum dominating set of the $a_e^1 - a_e^{p-1}$ path and vertices from minimum dominating set of the $b_e^2 - b_e^p$ path, i.e.,

we set $D' = (D' \setminus (D_e^a \cup D_e^b)) \cup R_e$. We note that after applying this transformation to D' , it remains a dominating set of G' without increasing its size. After transforming of D' for each $e = \{u, v\} \in E$ such that $u \notin D'$ and $v \notin D'$ one after another, we obtain a dominating set D' of G' such that $|D'| \leq |D|$. It is not hard to see that for each $e = \{u, v\} \in E$ at least one of the vertices u, v is contained in D' and

$$|D' \cap \{u_e, v_e, a_e^1, a_e^2, \dots, a_e^p, b_e^1, b_e^2, \dots, b_e^p\}| \geq \gamma(P_p) + \gamma(P_{p-2}) = \frac{2p+1}{3}.$$

Therefore, the set $S = D' \cap V$ is a vertex cover of G and $|D'| \geq |S| + m \frac{2p+1}{3}$. Hence

$$|D| \geq |D'| \geq |S| + m \frac{2p+1}{3}.$$

If D is a dominating set of G' such that $|D| = \gamma(G')$ then we obtain $\gamma(G') \geq \tau(G) + m \frac{2p+1}{3}$.

Thus, $\gamma(G') = \tau(G) + \frac{m}{3}(2p+1)$ and G has a vertex cover of size at most k if and only if G' has a dominating set of size at most $k + \frac{m}{3}(2p+1)$, where m is the number of edges of G . \square

As for every bipartite C_4 -free graph G computing $b(G)$ is equivalent to computing $\gamma(G)$, the following statement is a direct consequence of Lemma 10.

Corollary 11 *For any $t \geq 2$, BICLIQUE VERTEX-PARTITION is NP-complete for 2-connected planar $(C_4, C_6, \dots, C_{2t})$ -free bipartite graphs of vertex degree at most 6. In particular, BICLIQUE VERTEX-PARTITION remains NP-complete for bipartite C_4 -free graphs.*

Now we study the computational complexity of BICLIQUE VERTEX-PARTITION for perfect elimination bipartite graphs. Let $G = ((X, Y), E)$ be a bipartite graph. An edge $e = \{x, y\} \in E$ is a *bisimplicial edge* if $N(x) \cup N(y)$ induces a biclique of G . Let (e_1, e_2, \dots, e_k) be a sequence of pairwise non-adjacent edges of G (not necessarily all edges of E). We denote by S_i the set of endpoints of the edges e_1, e_2, \dots, e_i and $S_0 = \emptyset$. The sequence (e_1, e_2, \dots, e_k) is a *perfect edge elimination ordering* for G if $G[(X \cup Y) \setminus S_k]$ is an empty graph and each edge e_i is a bisimplicial edge in $G[(X \cup Y) \setminus S_{i-1}]$. The graph G is a *perfect elimination bipartite graph* (see, for example, [2, p. 87], [9]) if G has a perfect edge elimination ordering.

Theorem 12 *BICLIQUE VERTEX-PARTITION is NP-complete for perfect elimination bipartite graphs.*

Proof: We make a simple reduction from BICLIQUE VERTEX-PARTITION, which is known to be NP-complete for bipartite graphs. Suppose we are given a bipartite graph $G = ((X, Y), E)$ with $X = \{x_1, x_2, \dots, x_p\}$, $Y = \{y_1, y_2, \dots, y_q\}$ and $p \geq 1$. We transform G into a bipartite graph $G' = ((X', Y'), E')$, where $X' = X \cup \{b_1, b_2, \dots, b_p\}$, $Y' = Y \cup \{a_1, a_2, \dots, a_p\} \cup \{c_1, c_2, \dots, c_p\}$ and

$$E' = E \cup \bigcup_{i=1}^p \{\{x_i, a_i\}, \{a_i, b_i\}, \{b_i, c_i\}\}.$$

Thus, G' is obtained from G by appending a path $\{x_i, a_i, b_i, c_i\}$ to x_i for every $i = 1, 2, \dots, p$ (see Figure 6 for an example). The graph G' is a perfect elimination bipartite graph. The sequence

$$\sigma = (\{b_1, c_1\}, \{b_2, c_2\}, \dots, \{b_p, c_p\}, \{x_1, a_1\}, \{x_2, a_2\}, \dots, \{x_p, a_p\})$$

is a perfect edge elimination ordering for G' . We claim that $b(G') = b(G) + p$.

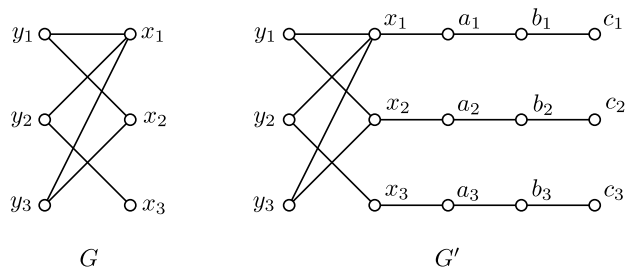


Fig. 6: A bipartite graph G and the graph G' obtained from G .

For each $i = 1, 2, \dots, p$, we denote by B_i the star of G' centred at the vertex b_i that contains the vertices b_i and c_i , i.e., $B_i = ((\{b_i\}, \{a_i, c_i\}), \{\{b_i, a_i\}, \{b_i, c_i\}\})$. Let \mathcal{S} be a biclique vertex-partition of G . Form \mathcal{S}' by adding the stars B_1, B_2, \dots, B_p to \mathcal{S} . It is easy to see that \mathcal{S}' is a biclique vertex-partition of G' of size $|\mathcal{S}| + p$. Taking \mathcal{S} with the size $b(G)$, we obtain $b(G') \leq b(G) + p$. To show the converse inequality, we use the following simple observation.

Observation 13 *The graph G' has a biclique vertex-partition \mathcal{S}' of size $b(G')$ that contain the stars B_1, B_2, \dots, B_p .*

Let \mathcal{S}' be a biclique vertex-partition of G' from Observation 13. Construct \mathcal{S} from \mathcal{S}' by eliminating the stars B_1, B_2, \dots, B_p . It is easy to see that \mathcal{S} is a biclique vertex-partition of G of size $b(G') - p$. Hence $b(G) \leq b(G') - p$.

Thus, we have that $b(G') = b(G) + p$ and, as a consequence, the bipartite graph G has a biclique vertex-partition of size at most k if and only if the perfect elimination bipartite graph G' has a biclique vertex-partition of size at most $k + p$. \square

5 Conclusion and further work

In this paper, we have shown that the BICLIQUE VERTEX-PARTITION problem is NP-complete for perfect elimination bipartite graphs and have provided two subclasses of perfect elimination bipartite graphs for which this problem is solvable in polynomial time. To refine the borderline between polynomial time solvability and NP-completeness, it would be interesting to establish the computational complexity of the given graph problem for convex bipartite graphs and for chordal bipartite graphs.

We have also shown that the BICLIQUE VERTEX-PARTITION problem remains NP-complete even for bipartite C_4 -free graphs. It would be interesting to find subclasses of bipartite C_4 -free graphs for which the considered problem is polynomially solvable.

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