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On the asymptotic enumeration of accessible automata

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We simplify the known formula for the asymptotic estimate of the number of deterministic and accessible automata with \( n \) states over a \( k \)-letter alphabet. The proof relies on the theory of Lagrange inversion applied in the context of generalized binomial series.

Keywords: asymptotic enumeration, finite automata, Lagrange series, generalized binomial series.

1 Main result

Let \( A_n \) be the set of non-isomorphic accessible (also called initially connected) complete and deterministic automata with \( n \) states over a \( k \)-letter alphabet. By using suitable combinatorial transformations, Bassino and Nicaud (2007) prove that the order of magnitude of the cardinality of \( A_n \) is related to the Stirling number of the second kind \( {kn \choose n} \). Then they reformulate the asymptotic estimate of \(|A_n|\) due to Korshunov (1978) in terms of this Stirling number. Our purpose is to state a simplification of this formula.

First we need to recall a few definitions. For \( m, n \) non-negative integers, the Stirling number of the second kind, denoted by \( \{m \atop n\} \), is the number of ways of partitioning a set of \( n \) elements into \( m \) non-empty subsets. Let \( W_0 \) denote the principal branch of the so-called Lambert \( W \) function (which is the inverse of the function \( x \mapsto xe^x \)). See Graham et al. (1994) and Corless et al. (1996) for more details about these definitions. Also define the constants

\[
v_k = 1 + \frac{W_0(-ke^{-k})}{k},
\]

\[
\zeta_k = k v_k = k + W_0(-ke^{-k}),
\]

\[
\beta_k = \frac{k^k}{e^{k-1}} \left( e^{\zeta_k} - 1 \right) \zeta_k^k.
\]

Except for \( v_k \), these constants are considered in Bassino and Nicaud (2007). Notice that \( v_k \) is the unique root in \((0, 1)\) of the function \( \phi_k(v) = kv + \log(1-v) \) and it satisfies \( v_k > (k-1)/k \). See Figure [1]

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The next result, due to Bassino and Nicaud (2007), reformulates in terms of the Stirling numbers the asymptotic estimate for $|A_n|$ previously obtained by Korshunov (1978).

**Theorem 1 (Bassino and Nicaud (2007))** The number $|A_n|$ of accessible complete and deterministic automata with $n$ states over a $k$-letter alphabet satisfies

$$|A_n| \sim E_k n 2^n \binom{k n}{n},$$

where

$$E_k = \frac{1 + \sum_{r=1}^{\infty} \frac{1}{r} \left( \frac{k r}{r-1} \right) (e^{k-1} \beta_k)^{-r}}{1 + \sum_{r=1}^{\infty} \frac{1}{r} \left( e^{k-1} \beta_k \right)^{-r}}.$$"
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</tbody>
</table>

**Tab. 1:** Values of \(v_k\) and \(E_k\).

2  **Proof of Theorem 2**

The proof relies on the theory of Lagrange series, which constitutes a fundamental tool for the solution of implicit equations and for the reversion of series. We refer to [Henrici (1988)] for an extensive treatment of this subject. To recall the two basic forms of Lagrange series, suppose that the dependence between the variables \(y\) and \(z\) is implicitly given by the equation

\[
y = \frac{z}{\varphi(z)},
\]

where \(\varphi\) is a regular function in a neighborhood of the point \(z = 0\) and \(\varphi(0) \neq 0\). If \(f\) is a regular function in a neighborhood of \(z = 0\), then

\[
f(z) = f(0) + \sum_{r=1}^{\infty} \frac{y^r}{r!} \left[ \frac{d^{r-1} \left( f'(z) \varphi(z) \right)^r}{dz^{r-1}} \right]_{z=0},
\]

where the primes denote derivatives.

To prove Theorem 2, we study a classical application of this theory, namely, the solutions of Problems 212 and 216 in Pólya and Szegő (1998, Chapter 5 of Part III). Let \(a\) and \(b\) be real numbers and consider the equation 

\[
1 - x + y x^b = 0,
\]

that is, 

\[
y = \frac{x - 1}{x - b}.
\]

We define \(z = x - 1\) and apply the formulae in (1) with \(\varphi(z) = (1 + z)^b\) and \(f(z) = (1 + z)^a\). Notice that

\[
\frac{1}{r!} \left[ \frac{d^{r-1} \left( f'(z) \varphi(z) \right)^r}{dz^{r-1}} \right]_{z=0} = \frac{1}{r!} \left[ \frac{d^{r-1} \left( a (1 + z)^{a+br-1} \right)}{dz^{r-1}} \right]_{z=0} = \frac{a}{a + br} \left( \frac{a + br}{r} \right).
\]

Hence, (1) yields

\[
x^a = \sum_{r=0}^{\infty} \frac{a}{a + br} \left( \frac{a + br}{r} \right) y^r \tag{2}
\]

and

\[
\frac{x^{a+1}}{(1 - b)x + b} = \sum_{r=0}^{\infty} \left( \frac{a + br}{r} \right) y^r, \tag{3}
\]
with \( y = (x-1)x^{-b} \). Next, we consider the case \( b > 1 \) and investigate the neighborhood in \( \mathbb{R} \) of the point \( x = 1 \) in which (2) and (3) hold.

For \( a = 1 \), the power series appearing in (2)

\[
B_b(y) := \sum_{r=0}^{\infty} \frac{br+1}{br+1} y^r
\]

is called the generalized binomial series. Thus, \( B_b \) is the solution of the functional equation \( x(y) = 1 + y x(y)^b \) that satisfies \( x(0) = 1 \). Moreover, (2) and (3) can be rewritten as

\[
B_b(y) = \sum_{r=0}^{\infty} \frac{a}{a+br} \binom{a+br}{r} y^r \quad \text{and} \quad \frac{B_b(y)^{a+1}}{(1-b)B_b(y) + b} = \sum_{r=0}^{\infty} \binom{a+br}{r} y^r.
\]

In [Graham et al. (1994), Sections 5.4 and 7.5], \( B_b \) is introduced as a formal power series, these properties are proved and used to derive some combinatorial identities. See also [Heggie and Nickerson (1996)] for a representation of the function \( B_b \) as a contour integral in the complex plane.

Now let \( b > 1 \) be fixed. We are interested in studying the set of values of \( x \) for which the identity \( B_b((x-1)x^{-b}) = x \) holds. To this end, we observe that by the ratio test the series in (4) converges for \( |y| < R \), where

\[
R := \frac{(b-1)^{b-1}}{b}. \]

**Lemma 1** There exists \( \bar{x}_b \in (1/2, 1) \) such that

\[ B_b((x-1)x^{-b}) = x \]

for all \( x \in (\bar{x}_b, b(b-1)^{-1}) \).

**Proof:** It is not difficult to show that the function \( g_b(x) = (x-1)x^{-b} \) is continuous and strictly increasing in the interval \((0, b(b-1)^{-1})\). The result follows since \( g_b(b(b-1)^{-1}) = R \), \( g_b(1) = 0 \) and \( |g_b(1/2)| > 1 \).

Consequently, identities (2) and (3) are valid for every \( x \in (\bar{x}_b, b(b-1)^{-1}) \).

Now define \( \bar{v}_b = (\bar{x}_b)^{-1} \in (1, 2) \). From Lemma 1 we obtain the following result.

**Lemma 2** For any \( v \in ((b-1)/b, \bar{v}_b) \),

\[
\frac{1}{v^a} = \sum_{r=0}^{\infty} \frac{a}{a+br} \binom{a+br}{r} y^r \quad \text{and} \quad \frac{1}{v^a} \frac{1}{b v - (b-1)} = \sum_{r=0}^{\infty} \binom{a+br}{r} y^r,
\]

with \( y = (1-v)v^{b-1} \).

**Proof:** For each \( v \in ((b-1)/b, \bar{v}_b) \), define \( x = 1/v \), so \( x \in (\bar{x}_b, b(b-1)^{-1}) \) and \( x^{-b} = (1-v)v^{b-1} \). The result follows from (2) and (3). \( \square \)
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We are ready to prove our main result.

**Proof of Theorem 2:** Note that

\[
e^{k-1} \beta_k = \frac{k^k (e^\zeta_k - 1)}{\zeta_k} = \frac{1}{(1-v_k)v_k^{k-1}}.
\]

Defining \(y_k = (1 - v_k)v_k^{k-1}\), we conclude from Lemma 2 that the numerator of \(E_k\) equals

\[
1 + \sum_{r=1}^{\infty} \frac{1}{r} \binom{kr}{r-1} y_k^r = \sum_{r=0}^{\infty} \frac{1}{kr+1} \binom{kr+1}{r} y_k^r = \frac{1}{v_k},
\]

whereas the denominator simplifies to

\[
\sum_{r=0}^{\infty} \binom{kr}{r} y_k^r = \frac{1}{kv_k - (k-1)}.
\]

The result follows. \(\square\)

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**References**


